CS344: Introduction to Artificial Intelligence

Pushpak Bhattacharyya
CSE Dept.,
IIT Bombay

Lecture 8 and 9– Predicate Calculus; Interpretation; Inferencing

Predicate Calculus: well known examples

- Man is mortal : rule

  $\forall x[man(x) \rightarrow mortal(x)]$

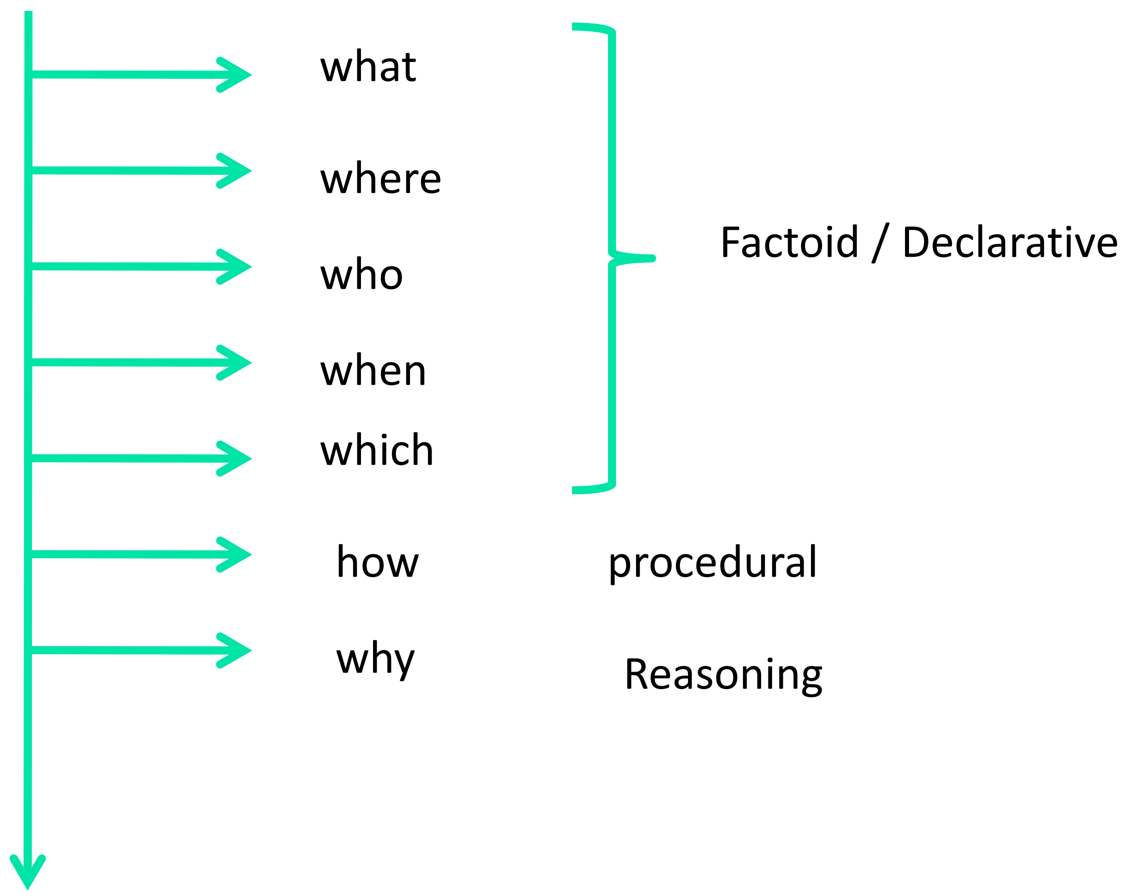
- shakespeare is a man

  man(shakespeare)

- To infer shakespeare is mortal

  mortal(shakespeare)

Wh-Questions and Knowledge



Fixing Predicates

- Natural Sentences

<Subject> <verb> <object>

Verb(subject,object)

predicate(subject)

Examples

- Ram is a boy
  - Boy(Ram)?
  - Is_a(Ram,boy)?

- Ram Playes Football
  - Plays(Ram,football)?
  - Plays_football(Ram)?

Knowledge Representation of Complex Sentence

- *"In every city there is a thief who is beaten by every policeman in the city"*

$$\forall x[\text{city}(x) \rightarrow \{\exists y((\text{thief}(y) \wedge \text{lives\_in}(y, x)) \wedge \forall z(\text{poleceman}(z, x) \rightarrow \text{beaten\_by}(z, y)))\}]$$

Interpretation in Logic

- Logical expressions or formulae are "FORMS" (placeholders) for whom contents are created through interpretation.
- Example:

$$\exists F\left[\{F(a)=b\}\wedge\forall x\{P(x)\rightarrow\left(F(x)=g\left(x,F\left(h(x)\right)\right)\right)\}\right]$$

- This is a Second Order Predicate Calculus formula.
- Quantification on 'F' which is a function.

Examples

- Interpretation:1

  $D=N$ (natural numbers)

  $a = 0$ and $b = 1$

  $x \in N$

  $P(x)$ stands for $x > 0$

  $g(m,n)$ stands for $(m \times n)$

  $h(x)$ stands for $(x - 1)$
- Above interpretation defines **Factorial**

Examples (contd.)

- Interpretation:2

  D={strings}

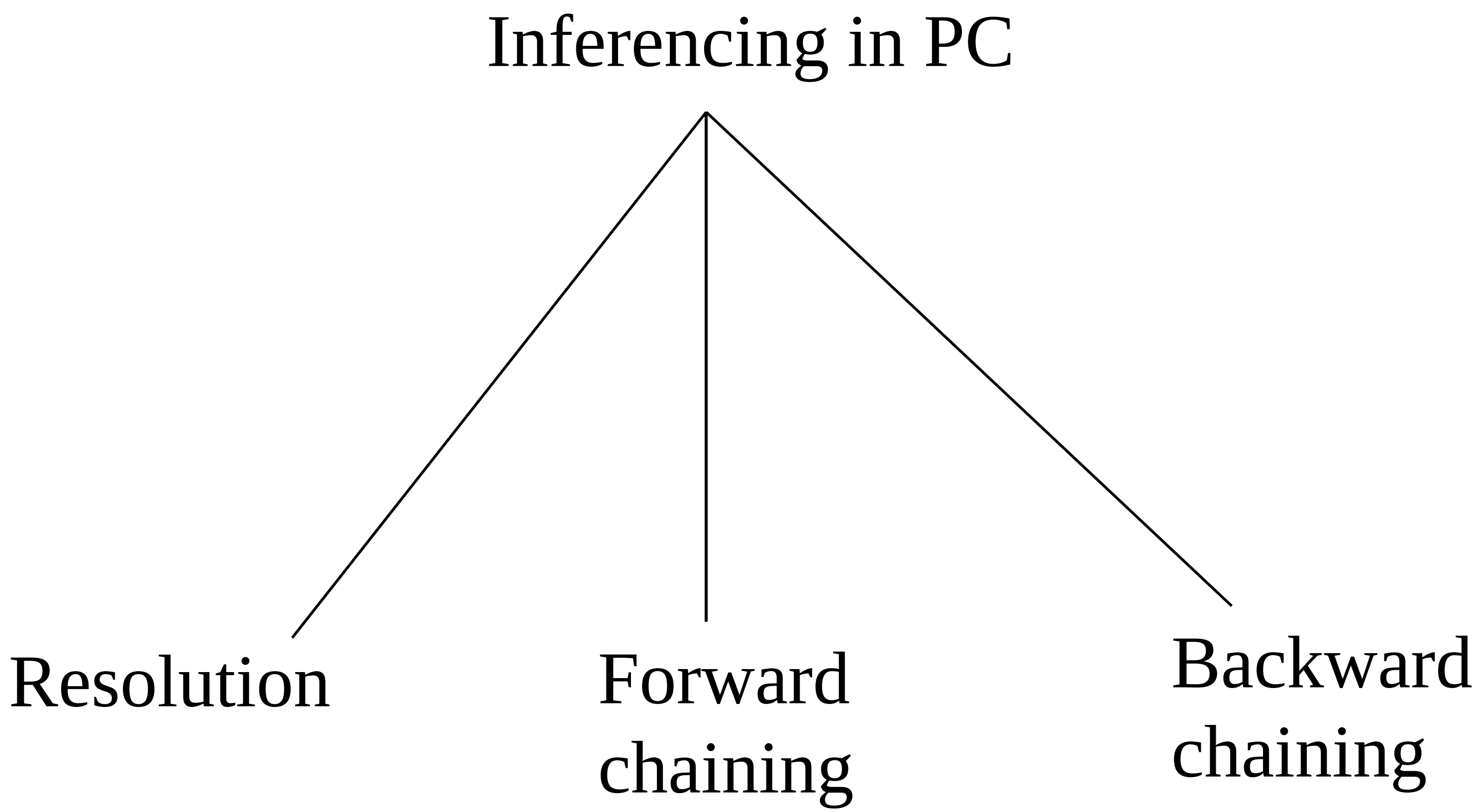
  $a = b = \lambda$

  $P(x)$ stands for “x is a non empty string”

  $g(m, n)$ stands for “append head of m to n”

  $h(x)$ stands for $tail(x)$

- Above interpretation defines “reversing a string”


Inferencing in PC
Resolution
Forward chaining
Backward chaining

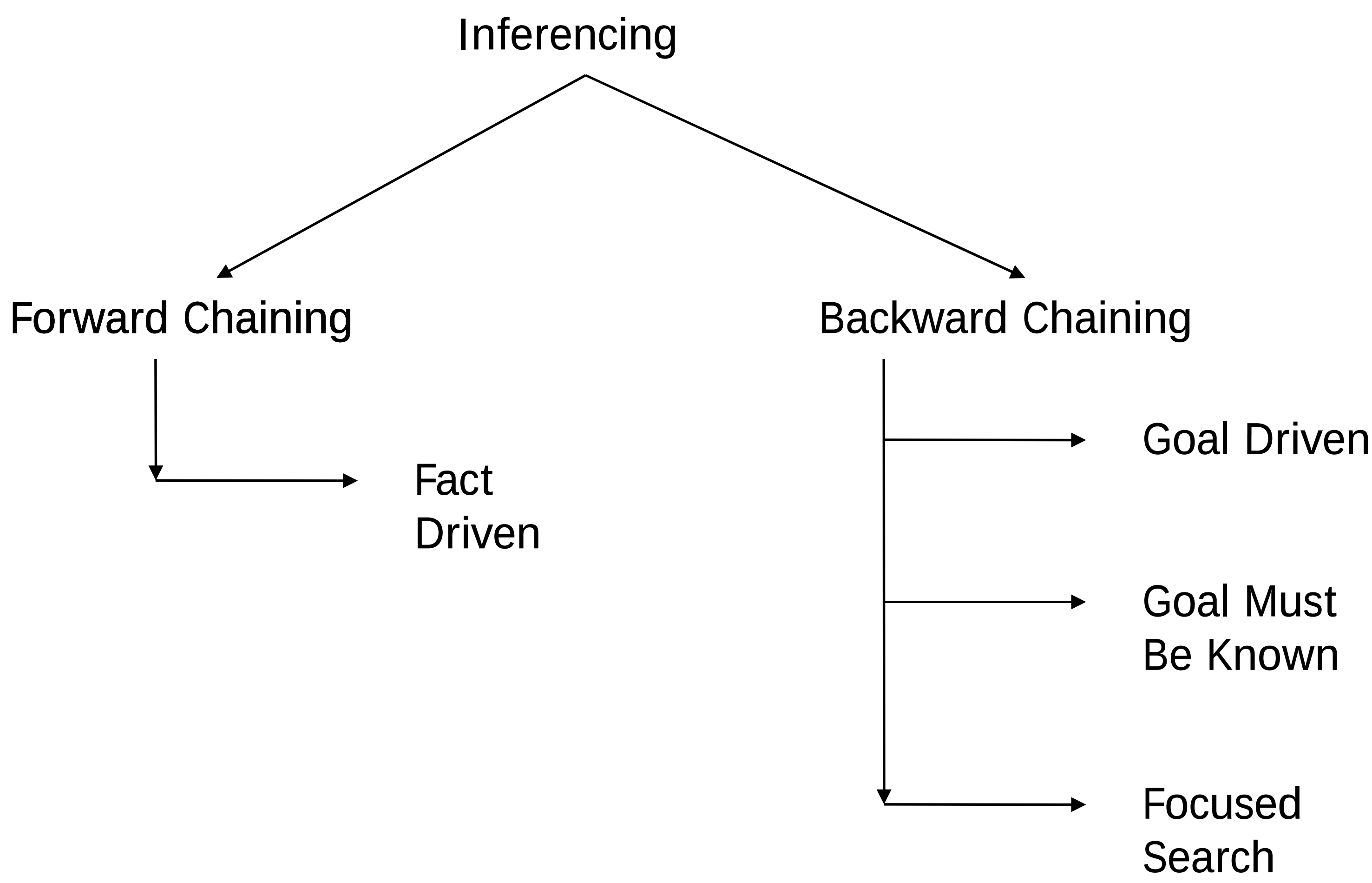
Forward Chaining/ Inferencing

- $man(x) \rightarrow mortal(x)$
  - *Dropping the quantifier, implicitly Universal quantification assumed*
  - *man(shakespeare)*
- Goal mortal(shakespeare)
  - Found in one step
  - x = shakespeare, unification

Backward Chaining/ Inferencing

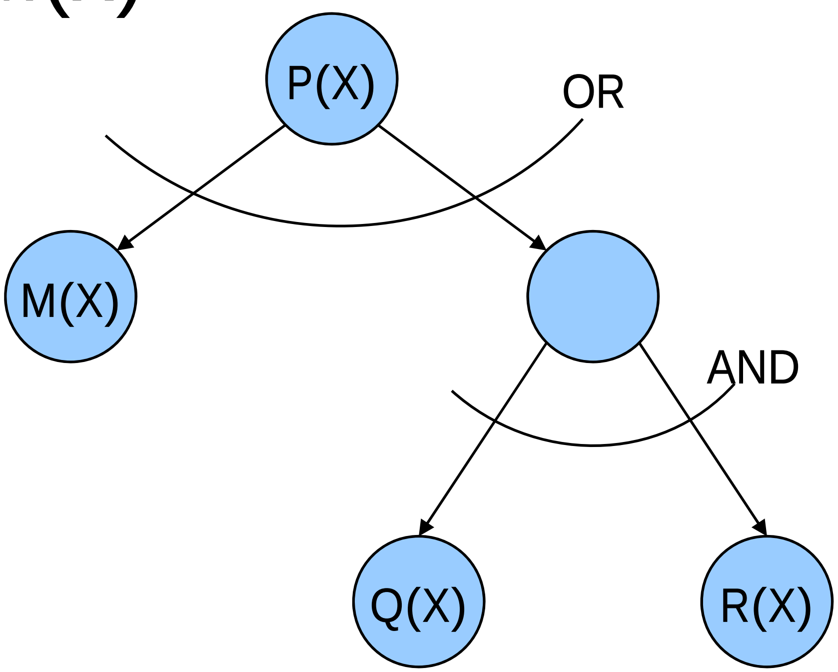
- $man(x) \rightarrow mortal(x)$
- Goal mortal(shakespeare)
  - x = shakespeare
  - Travel back over and hit the fact asserted
  - man(shakespeare)

Inferencing



AND-OR Graphs

- $P(x) \rightarrow Q(x)$ AND $R(x)$
- $P(x) \rightarrow M(x)$



AND-OR Graphs

- Whole rule base can be represented as an AND-OR graph
- AO* Search: A* like search on AND-OR graphs

Deciding the Strategy

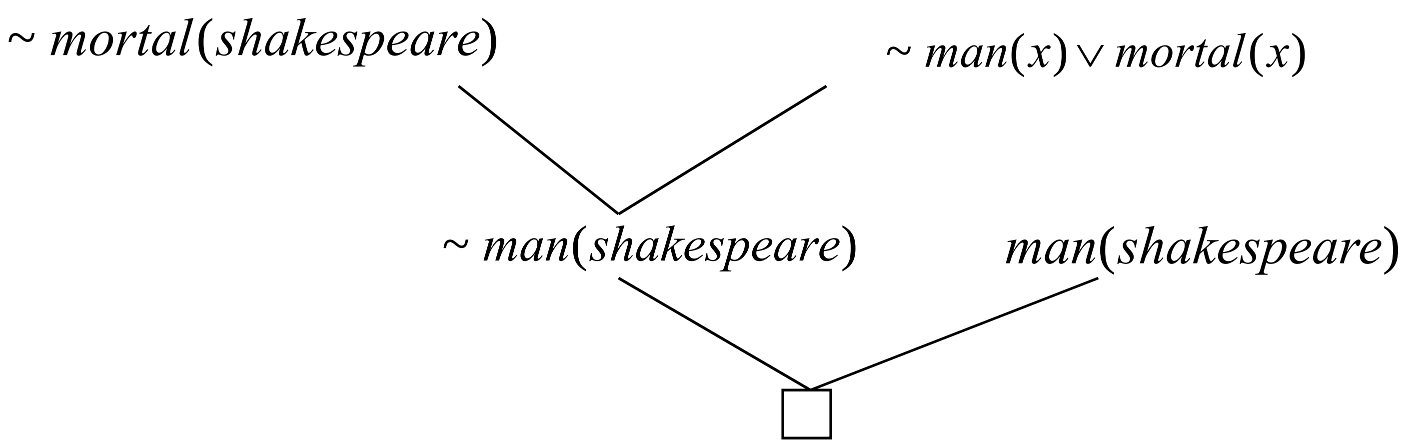
- Forward Chaining/Backward Chaining is decided from:
  - AO Graph and,
  - OR Fan-Out and,
  - Fan-In of Goal Node

Hilbert's Logical Axioms

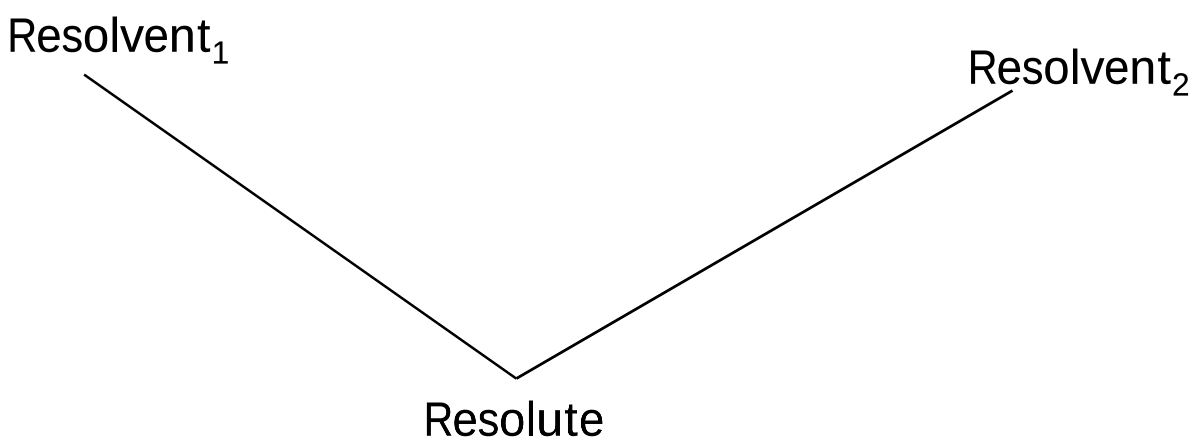
- $P(x) \rightarrow (Q(x) \rightarrow P(x))$
- $[P(x) \rightarrow (Q(x) \rightarrow R(x))] \rightarrow [(P(x) \rightarrow Q(x)) \rightarrow (P(x) \rightarrow Q(x))]$
- $\sim(\sim P(x)) \rightarrow P(x)$

Resolution – Refutation contd

- *Negate the goal*
  - *~mortal(shakespeare)*
- Get a pair of resolvents



Resolution Tree



Search in resolution

- Heuristics for Resolution Search
  - Goal Supported Strategy
    - Always start with the negated goal
  - Set of support strategy
    - Always one of the resolvents is the most recently produced resolute

Search in resolution

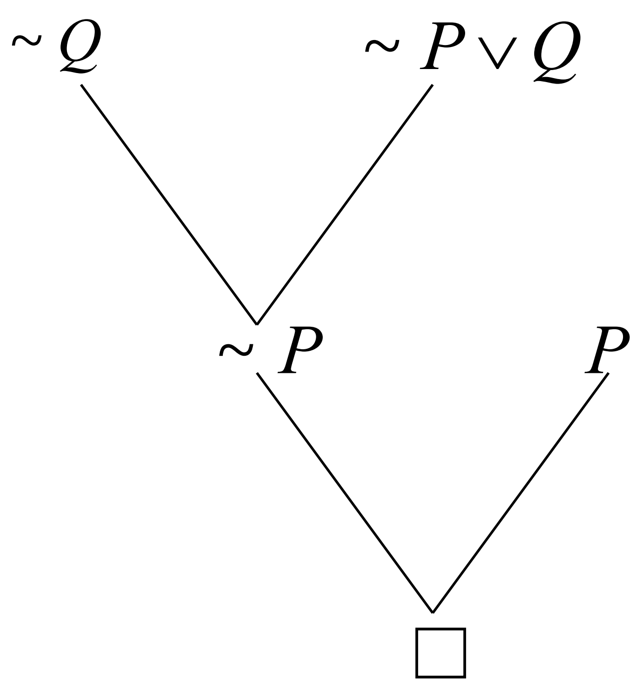
- Heuristics for Resolution Search
  - Goal Supported Strategy
    - Always start with the negated goal
  - Set of support strategy
    - Always one of the resolvents is the most recently produced resolute

Inferencing in Predicate Calculus

- Forward chaining
  - Given P, $P \rightarrow Q$, to infer Q
  - P, match $L.H.S$ of
  - Assert Q from $R.H.S$
- Backward chaining
  - Q, Match $R.H.S$ of $P \rightarrow Q$
  - assert P
  - Check if P exists
- Resolution – Refutation
  - Negate goal
  - Convert all pieces of knowledge into clausal form (disjunction of literals)
  - See if contradiction indicated by null clause $\square$ can be derived

1. P
2. $P \rightarrow Q$ converted to $\sim P \vee Q$
3. $\sim Q$

Draw the resolution tree (actually an inverted tree). Every node is a clausal form and branches are intermediate inference steps.



Terminology

- Pair of clauses being resolved is called the Resolvents. The resulting clause is called the Resolute.
- Choosing the correct pair of resolvents is a matter of search.

Himalayan Club example

- Introduction through an example *(Zohar Manna, 1974):*
  - Problem: A, B and C belong to the Himalayan club. Every member in the club is either a mountain climber or a skier or both. A likes whatever B dislikes and dislikes whatever B likes. A likes rain and snow. No mountain climber likes rain. Every skier likes snow. *Is there a member who is a mountain climber and not a skier?*
- Given knowledge has:
  - Facts
  - Rules

Example contd.

- Let *mc* denote mountain climber and *sk* denotes skier. Knowledge representation in the given problem is as follows:
  1. *member(A)*
  2. *member(B)*
  3. *member(C)*
  4. $\forall x[member(x) \rightarrow (mc(x) \vee sk(x))]$
  5. $\forall x[mc(x) \rightarrow \sim like(x,rain)]$
  6. $\forall x[sk(x) \rightarrow like(x, snow)]$
  7. $\forall x[like(B, x) \rightarrow \sim like(A, x)]$
  8. $\forall x[\sim like(B, x) \rightarrow like(A, x)]$
  9. *like(A, rain)*
  10. *like(A, snow)*
  11. *Question:* $\exists x[member(x) \wedge mc(x) \wedge \sim sk(x)]$
- We have to infer the 11th expression from the given 10.
- Done through Resolution Refutation.

Club example: Inferencing

1. *member(A)*
2. *member(B)*
3. *member(C)*
4. $\forall x[member(x) \rightarrow (mc(x) \vee sk(x))]$
   - Can be written as
     $$[member(x) \rightarrow (mc(x) \vee sk(x))]$$
   - $$\sim member(x) \vee mc(x) \vee sk(x)$$
5. $\forall x[sk(x) \rightarrow lk(x, snow)]$
   - $$\sim sk(x) \vee lk(x, snow)$$
6. $\forall x[mc(x) \rightarrow \sim lk(x, rain)]$
   - $$\sim mc(x) \vee \sim lk(x, rain)$$
7. $\forall x[like(A, x) \rightarrow \sim lk(B, x)]$
   - $$\sim like(A, x) \vee \sim lk(B, x)$$

8. $\forall x[\sim lk(A,x) \rightarrow lk(B,x)]$

– $lk(A,x) \vee lk(B,x)$

9. $lk(A,rain)$

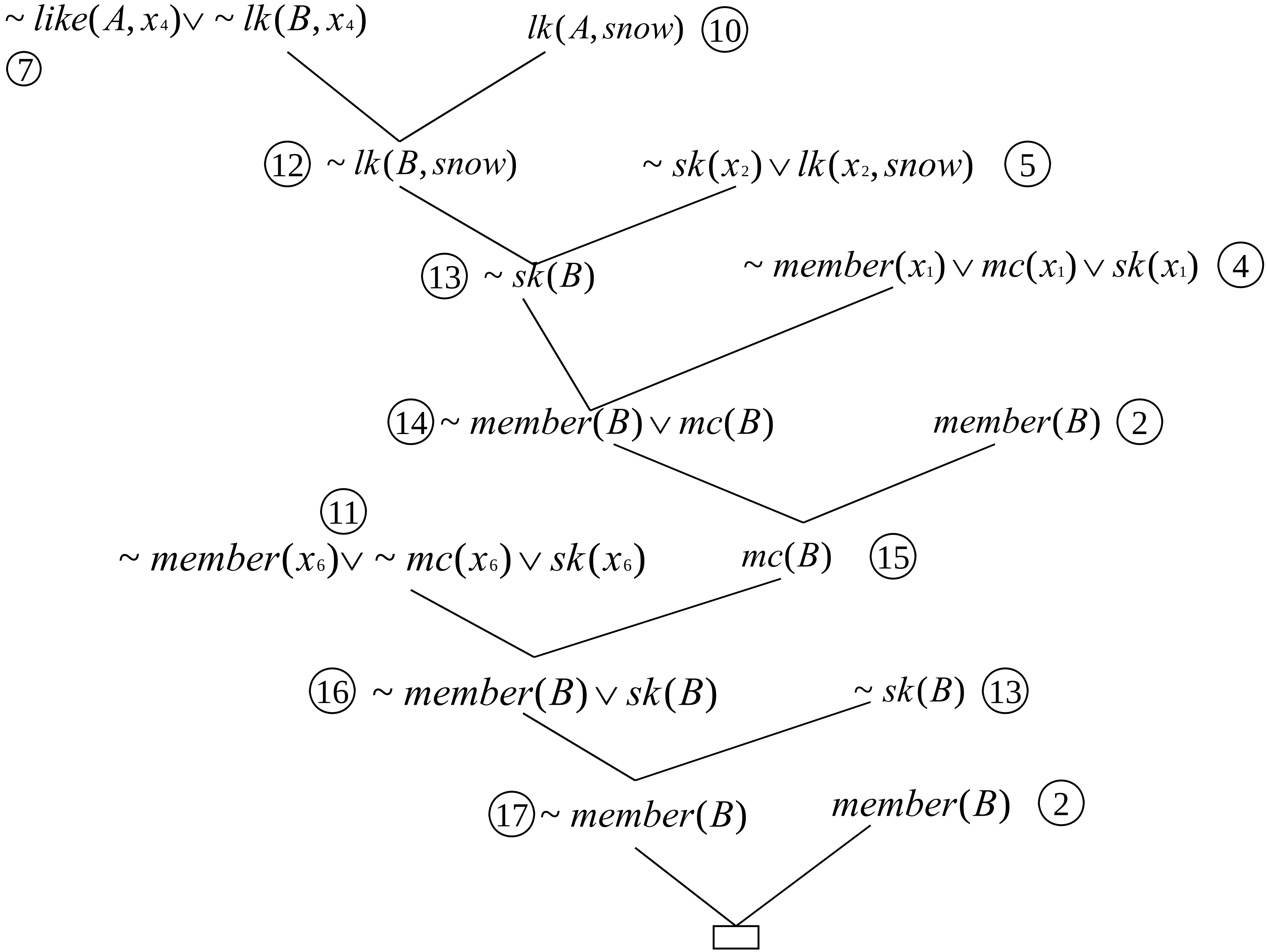
10. $lk(A,snow)$

11. $\exists x[member(x) \wedge mc(x) \wedge \sim sk(x)]$

– Negate– $\forall x[\sim member(x) \vee \sim mc(x) \vee sk(x)]$

- Now standardize the variables apart which results in the following

1. $member(A)$
2. $member(B)$
3. $member(C)$
4. $\sim member(x_1) \vee mc(x_1) \vee sk(x_1)$
5. $\sim sk(x_2) \vee lk(x_2, snow)$
6. $\sim mc(x_3) \vee \sim lk(x_3, rain)$
7. $\sim like(A, x_4) \vee \sim lk(B, x_4)$
8. $lk(A, x_5) \vee lk(B, x_5)$
9. $lk(A, rain)$
10. $lk(A, snow)$
11. $\sim member(x_6) \vee \sim mc(x_6) \vee sk(x_6)$


~ like(A, x4) ∨ ~ lk(B, x4) ⑦
lk(A, snow) ⑩
⑫ ~ lk(B, snow)
~ sk(x2) ∨ lk(x2, snow) ⑤
⑬ ~ sk(B)
~ member(x1) ∨ mc(x1) ∨ sk(x1) ④
⑭ ~ member(B) ∨ mc(B)
member(B) ②
⑪
~ member(x6) ∨ ~ mc(x6) ∨ sk(x6)
mc(B) ⑮
⑯ ~ member(B) ∨ sk(B)
~ sk(B) ⑬
⑰ ~ member(B)
member(B) ②

Assignment

- Prove the inferencing in the Himalayan club example with different starting points, producing different resolution trees.
- Think of a Prolog implementation of the problem
- Prolog Reference (Prolog by Chockshin & Melish)