

# Ideas from Communicating Sequential Processes

Prof. RKJ  
CS 451 Lecture  
2003



# Background of the times

- ◆ Basic Structuring methods of computer programs

1. Repetitive construct (while, ...)
2. Alternative command (if, ...)
3. Normal sequential program composition (;)

- ◆ Other important program structures

subroutines

Coroutines

Classes

Processes

monitors

Actors

...

# Hoare's contribution

- ◆ Components on multiprocessors must be able to communicate with each other
- ◆ What was existing: shared storage, semaphores, monitors, conditional critical regions..
- ◆ Hoare suggested that  
input and output are basic primitives of programming

Parallel composition of sequential processes is a fundamental program structuring method

→ Development of CSP (Communicating Sequential Processes Paradigm)

# Contributions..

- ◆ Express parallel processes  $p1::P1 \parallel p2::P2$
- ◆  $P1$  and  $P2$  can be composed together through input and output
- ◆ Inputs can be used in guards
- ◆ Guarded commands can be clubbed in an alternative statement through nondeterminism
- ◆ Repeat alternative command

# CSP Basic primitives

- ◆ Specifying parallel processes
- ◆ Communication between processes through input and output commands
- ◆ Guarded commands for obtaining nondeterminism
- ◆ Repetitive command
- ◆ Pattern matching feature for determining input messages

# Specifying parallel processes

- ◆ Specifying parallel processes

```
[room::ROOM ||  
fork(i:0..4)::FORK ||  
phil(0..4)::PHIL]
```

ROOM, FORK, PHIL provide the process specifications or command lists, and room, fork, and phil are the processes created

# Input and output commands

◆  $X?i$

Input a value from  $X$  and assign it to  $i$

◆  $Y!10$

Output value 10 to  $Y$

◆  $Sem!P()$

Output signal  $P()$  to process  $Sem$

◆  $X(i)?V()$

Input a signal  $V()$  from the given process; refuse any other signal

◆  $X?(x,y)$

Input a pair of values and assign them to  $x$  and  $y$  respectively

# Guarded commands and Alternative Command

Guarded commands  $G \rightarrow CL$   
that is, if guard  $G$  is true, execute command list  $CL$

◆  $G_1 \rightarrow CL_1 \mid G_2 \rightarrow CL_2 \mid \dots \mid G_n \rightarrow CL_n$

execute any one of the command list whose guards are found to be true: Nondeterministic choice

IF all guards fail, alternative command fails



# Repetitive command

◆ AS many iterations of its constituent alternative command as possible

◆  $*[ G_1 \rightarrow CL_1 \mid G_2 \rightarrow CL_2 \mid \dots \mid G_n \rightarrow CL_n ]$

execute any one of the command list whose guards are found to be true: Nondeterministic choice and repeat

IF all source processes fail, alternative command fails, else processes wait for corresponding outputs

Try Bounded Buffer



# Bounded Buffer

```
[  
  buffer:: int item; int buff[MAX]; ....  
*]
```

```
(Prod?item)and(buffernotfull)→do insertion
```

```
(consumer?consume())and(buffernotempty)→remove an item;  
  consumer.removeditem
```

```
]
```

```
|| prod::PROD || consumer::CONSUMER  
]
```

Try dining philosophers



# Dining Philosophers

```
[room::ROOM || fork[i:0..4]::FORK ||  
  phil[i:0..4]::PHIL]
```

1. ROOM

```
  [int i, count;
```

```
    *[(i:0..4)phil(i)?enter()andcount<4→
```

```
      count=count+1
```

```
    [(i:0..4)phil(i)?exit()→count=count-1
```

```
  ]
```

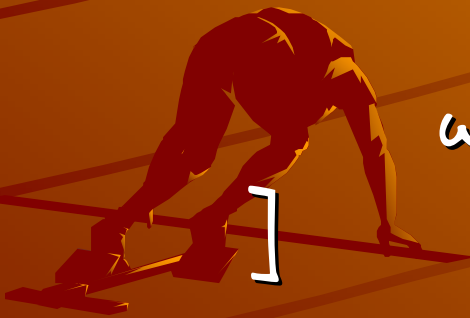


# Example 1

X :: \* [

c:char;

west?c → east!c



# Example 2

X :: \* [

c: char;

west?c → [

c != '\*' → east!c

| c == '\*' → west?c;

[c != '\*' → east!'\*';

east!c;

| c == '\*' → east!'^';

]

]

]

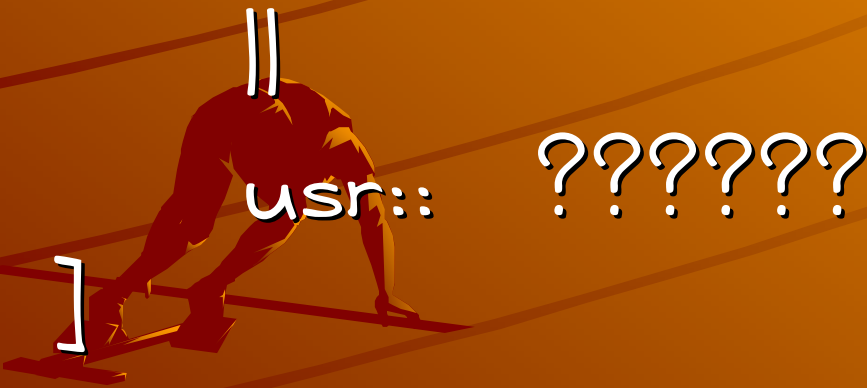
Assume that final character is not '\*'

How do you deal with odd no. of '\*'s at the end?



# A subroutine

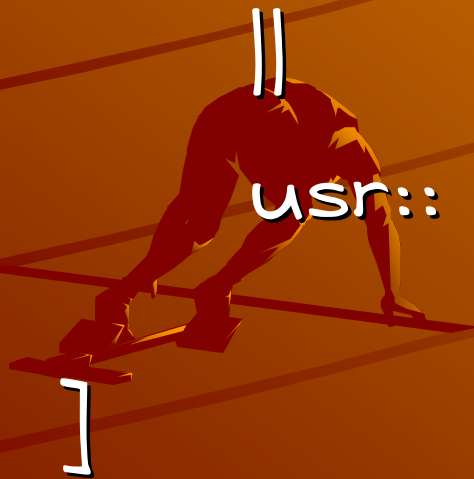
```
[  subr:: usr?(input parameters) →  
    ...;  
    X!(results)
```





# A subroutine

```
[  subr:: usr?(input parameters) →  
    ...;  
    usr!(results)
```




```
usr:: subr!(input params);  
subr?(results);  
]
```

# Example 3

```
[xyz (i: 1..max) ::  
  * [ n:int;  
      xyz(i-1)?n →
```

```
  [ n==0 → xyz(i-1)!1  
    n>0 → xyz(i+1)!n-1  
    r:int;  
    xyz(i+1)?r;  
    xyz(i-1)!(n*r)
```



```
xyz(0)::USER  
]
```

# Readings

- CAR Hoare, Communicating Sequential Processes, CACM, August 1978

