# CS101 Computer Programming and Utilization 

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(1) General Information
(2) Machines, Mechanisms and Programs
(3) Lights: A mechanism
(4) The Calculator
(5) CAL-Programs

## Course Organization

## Lectures and Laboratory

- Two classes per week, each of 1 hour.
- 2 hours of laboratory work.


## Resources

- Linux Handbook.
- gcc handbook.
- Tutorial and Worksheets.


## Machines

Machines are devices which execute or implement a given operation. Machines have instructions and states.

## Example

Stapler:

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| Ready | Hit | Ready | staples the sheets |

## Example

Vending Machine:

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| Ready | Tea | Ready | Vend Tea |
| Ready | Coffee | Ready | Vend Coffee |

## A more interesting example...

The state is a 2-tuple, such as [off,on], saying that electricity is off while water is on.

## Water Heater

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| [off,off] | SwitchOn | [on,off] | current! |
| [on,off] | TapOn | [on,on] | water! |
| $\vdots$ |  |  |  |

What is the point:

- While switching on, first turn the tap on and then the electricty.
- While switching off, first turn off the electricty and then the tap.


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## Water Heater

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| :---: | :---: | :---: | :---: |
| [off,off] | SwitchOn | [on,off] | current! |
| [on,off] | TapOn | [on,on] | water! |
| $\vdots$ |  |  |  |

What is the point:

- While switching on, first turn the tap on and then the electricty.
- While switching off, first turn off the electricty and then the tap.
- The state [on,off] is UNSTABLE.

Well, actually...

## Example

Stapler:

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| Ready | Hit | Ready | staples the sheets |
| Ready | Hit | Jam | messed it up |
| Jam | Hit | Jam | more mess |

## Example

Vending Machine:

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| Ready | Tea | Ready | Vend Tea |
| Ready | Tea | Over | no more |
| Over | Tea | Over | why??? |
| Ready | Coffee | Ready | Vend Coffee |

## Example

Tape Recorder: States: \{ FF, BB, play, idle, empty\}
Instructions \{ FF, BB, Play, Stop, OpenDoor

| Current State | Instruction | Next State | Action |
| :---: | :---: | :---: | :---: |
| FF | Stop | Idle | stops FF |
| Idle | Play | play | playing tape |
| play | Eject | error | block this |

In summary:

- A machine has a set of states and instructions.
- At any moment, a machine is in one of the states.
- An instruction causes some action and a change of state.
- Not all instructions are executable for any state.


## Mechanisms and Machines

## Mechanism

A mechanism is a device of supplying instrcutions to a machine.


- Thus the mechanism runs the machine by issuing instructions.
- The machine executes this instructions and updates its state.
- The mechanism can access this state and issues the next instruction.

Actually, most modern machines internally are mechanisms. For example, the vending machine above, when you ask for Tea does the following:

| now | Instruction | next | Remarks |
| :---: | :---: | :---: | :---: |
| 1 | Open Tea Powder Nozzle | 2 | allows powder into <br> mixing container |
| 2 | Shut Tea Powder Nozzle | 3 | enough tea powder |
| 3 | Open Hot Water Nozzle | 4 | into mixing container |
| 4 | Shut Hot Water Nozzle | 5 |  |
| 5 | Open Vending Nozzle | 6 |  |
| 6 | Shut Vending Nozzle | 1 |  |

Of course, if the machine state returns Jammed at anytime, then the execution is suspended.

## Fancy Lighting

Lets look at a more interesting mechanism.

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## Fancy Lighting

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## Fancy Lighting:How does it work?

## Brush

- Bulbs are grouped together into (what we call) circuits.
- Each circuit has a brush.
- There is a rotating drum with portrusions.
- Portrusions on the drum establish electrical contact.
- The portruding parts on the circle decide when the lights of that circuit are on.



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- The portruding parts on the circle decide when the lights of that circuit are on.


How do we describe this mechanism


| 122 | create a 1-by-22 grid |
| :--- | :--- | :--- |
| 1234561234561234561234 | the circuits |
| 1 magenta |  |
| 1111111111111111111111 | all the same colour |

This describes the hardware of the system. There is an array of $1 \times 22$ bulbs with fixed bulb colours, and so on.

## How do we describe this mechanism

| 1 | 1 | 1 | 1 | 0 | 0 | light up ckts $1,2,3,4$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 1 | 1 | 1 | 0 |  |
| 0 | 0 | 1 | 1 | 1 | 1 |  |
| 1 | 0 | 0 | 1 | 1 | 1 |  |
| 1 | 1 | 0 | 0 | 1 | 1 |  |
| 1 | 1 | 1 | 0 | 0 | 1 | the cycle ends |
| 0.5 |  |  |  |  |  |  |
| 0 |  |  |  |  |  |  |

This is the program, which resides on the cylinder of the mechanism. As the cylinder rotates, the desired pattern emerges.

## Program: A new concept

The sequence of instructions which are stored on the mechanism which drives the machine.

## Another Example

The Hardware

| 34 | create a 3-by-4 grid |
| :---: | :---: |
| 1111 | the circuits |
| 2222 |  |
| 3333 |  |
| 1 magenta | specifying colours |
| 2 green |  |
| 3 blue |  |
| 4 red |  |
| 1234 | the colour scheme |
| 2341 |  |
| 3412 |  |



The Program

| 1 | 0 | 0 | a 3-cycle |
| :--- | :--- | :--- | :--- |
| 0 | 1 | 0 |  |
| 0 | 0 | 1 |  |
| 0.5 |  | in seconds |  |

## Assignment 1

- Download the instructions, hardware and program files and execute the lighting program.
- Modify the above files to your taste and observe various combinations.
- Prepare a system which displays your roll-number, one digit after another.


## The Calculator as a Machine

| $\checkmark$ Calculator |  |  |  | - $\quad \mathrm{x}$ |
| :---: | :---: | :---: | :---: | :---: |
| Calculator |  | Edit Help | Help |  |
|  |  |  |  | 171 |
| 1/x | $\mathrm{x}^{2}$ | SQRT | CE/C | AC |
| INV | sin | cos | $\tan$ | DEG |
| e | EE | $\log$ | In | $\mathrm{x}^{\text {y }}$ |
| PI | x ! | ( | ) | 1 |
| STO | 7 | 8 | 9 | * |
| RCL | 4 | 5 | 6 | - |
| Sum | 1 | 2 | 3 | $+$ |
| EXC | 0 | . | +/- | $=$ |

## The basic calculator

Types of keys:

- Unary operators such as $\cos , \pm$,.
- Binary Operators such as,$+ *$.
- Digits
- Memory related: STO, RCL
- Special: = and AC.

Then there is the display.

| $\checkmark$ Calculator |  |  |  | - $\square$ |
| :---: | :---: | :---: | :---: | :---: |
| Calculator | Edit He |  | Help |  |
| 171 |  |  |  |  |
| 1/x | $\mathrm{x}^{2}$ | SQRT | CE/C | AC |
| INV | $\sin$ | $\cos$ | tan | DEG |
| e | EE | $\log$ | In | $\mathrm{x}^{\mathrm{y}}$ |
| PI | x ! | ( | ) | 1 |
| STO | 7 | 8 | 9 | * |
| RCL | 4 | 5 | 6 | - |
| SUM | 1 | 2 | 3 | + |
| EXC | 0 | . | +/- | $=$ |


|  | action | display |
| :---: | :---: | :---: |
| 1 | 3 | 3 |
| 2 | + | 3 |
| 3 | 2 | 2 |
| 4 | $=$ | 5 |

## Observation 1

The number 3 is stored internally and used during the binary operation. Let us call this as the internal register IR.

|  | action | display | IR |
| :---: | :---: | :---: | :---: |
| 1 | 3 | 3 | x |
| 2 | + | 3 | 3 |
| 3 | 2 | 2 | 3 |
| 4 | $=$ | 5 | x |
| x will stand for undefined. |  |  |  |


|  | action | display |
| :---: | :---: | :---: |
| 1 | 3 | 3 |
| 2 | + | 3 |
| 3 | 2 | 2 |
| 4 | $=$ | 5 |

## Observation 1

The number 3 is stored internally and used during the binary operation. Let us call this as the internal register IR.

|  | action | display | IR |
| :---: | :---: | :---: | :---: |
| 1 | 3 | 3 | x |
| 2 | STO | 3 | x |
| 3 | 2 | 2 | x |
| 4 | + | 2 | 2 |
| 5 | RCL | 3 | 2 |
| 6 | $=$ | 5 | x |

## Observation 2

The contents of the memory need to be stored as well. We call this as the M1 register.

|  | action | display | IR |
| :---: | :---: | :---: | :---: |
| 1 | 3 | 3 | x |
| 2 | + | 3 | 3 |
| 3 | 2 | 2 | 3 |
| 4 | $=$ | 5 | x |
| x will stand for undefined. |  |  |  |


| action | display | IR | M 1 |
| :---: | :---: | :---: | :---: |
| 3 | 3 | x | x |
| STO | 3 | x | 3 |
| 2 | 2 | x | 3 |
| + | 2 | 2 | 3 |
| RCL | 3 | 2 | 3 |
| $=$ | 5 | x | 3 |

Lets look at another set of computations.

| action | display | IR | M1 |
| :---: | :---: | :---: | :---: |
| 4 | 4 | x | x |
| SQRT | 2 | x | x |

and this:

| action | display | IR | M1 |
| :---: | :---: | :---: | :---: |
| 3 | 3 | x | x |
| + | 3 | 3 | x |
| 4 | 4 | 3 | x |
| SQRT | 2 | 3 | x |
| $=$ | 5 | x | x |

## Observation 3

A binary operator is evaluated as soon as its operands are specified. Is this really true?

## Observation 3

A unary operator is evaluated immediately. A pending binary operator is remembered.

| action | display | IR | M1 |
| :---: | :---: | :---: | :---: |
| 4 | 4 | x | x |
| $*$ | 4 | 4 | x |
| 3 | 3 | 4 | x |
| + | 12 | 12 | x |
| 2 | 2 | 12 | x |
| $=$ | 14 | x | x |


| action | display | IR | M1 |
| :---: | :---: | :---: | :---: |
| 4 | 4 | x | x |
| + | 4 | 4 | x |
| 3 | 3 | 4 | x |
| $*$ | 3 | $?$ | x |
| 2 | 2 | $?$ | x |
| $=$ | 10 | x | x |

## What is happening

The machine encounters a * which it decides must be evaluated before the + . So it has remembered 4 and + , and 3 when it encounters the $*$ and 2.

However, we will not worry about the internals of a calculator right now. We all know how to use it. Let us learn how to make it even more useful.

## Summary-The States of a Calculator

It is clear that the calculator remembers some operands and some operators. Let us assume that the calculator has the following four internal registers and 7 possible inputs:

| D | The Display register |
| :---: | :---: |
| I | The Invisible register |
| M | The memory regsiter |
| O | The current operator |


| op1 | a unary operator |
| :---: | :---: |
| op2 | a binary operator |
| $=$ | is equal to |
| num | a number |
| STO | memory store |
| RCL | memory recall |
| AC | All Cancel |

## The Calculator and the Program

Let us re-look at the machine-mechanism model and understand how we use a calculator.


We issue instructions and we observe the states.
\(\left.\begin{array}{|l|l|l|l|l|}\hline instruction <br>
\hline 23 <br>
+ <br>
10 <br>
= <br>
STO <br>

\vdots\end{array}\right]\)| $D$ | $I$ | $M$ |
| :--- | :--- | :--- |
| 23 | x | x |
| 23 | 23 | x |
| 10 | 23 | x |
| 33 | + |  |
| 33 | x | x |
| x |  |  |
| $\vdots$ | x |  |
| $\vdots$ |  |  |

## The Calculator and the Program

## Program

A program is a sequence of instructions.
Problem: Write a program to convert centigrades to farenheit.

| instruction | display |
| :--- | :--- |
| 40 | 40 |
| $*$ | 40 |
| 8 | 8 |
| DIV | 320 |
| 5 | 5 |
| + | 64 |
| 32 | 32 |
| $=$ | 96 |

## The Calculator and the Program

## Program

A program is a sequence of instructions.
Problem: Write a program to convert centigrades to farenheit.

| instruction | display |
| :--- | :--- |
| 0 | 0 |
| $*$ | 0 |
| 8 | 8 |
| DIV | 0 |
| 5 | 5 |
| + | 0 |
| 32 | 32 |
| $=$ | 32 |

## The Calculator and the Program

## Program

A program is a sequence of instructions.
Problem: Write a program to convert centigrades to farenheit.

| instruction | display |
| :--- | :--- |
| 10 | 10 |
| $*$ | 10 |
| 8 | 8 |
| DIV | 80 |
| 5 | 5 |
| + | 16 |
| 32 | 32 |
| $=$ | 48 |

## The Calculator and the Program

## Program

A program is a saved sequence of instructions.

```
10 % substitute centigrades here
*
8
DIV
5
+
32
= % observe the answer in the display
```

The above may be saved on a piece of paper or written on a computer file and re-used when necessary. It can be shared and transmitted. It can be written in Bangalore but executed in New York.

## Another Program

## Problem

Write a CAL-Program to compute the polynomial $p(t)=3 t^{2}+2 t+1$.

$$
\text { We use } 3 t^{2}+2 t+1=(3 * t+2) * t+1
$$

1 \% substitute $t$ here
STO \% stored in memory
3
*
RCL
$+$
2
$=\quad \% 3 t+2$ is done
*
RCL
$+$
1
= \% read answer in display

## Sample Executions

| instruction | Display | Memory |
| :--- | :--- | :--- |
| 2 | 2 | x |
| STO | 2 | 2 |
| 3 | 3 | 2 |
| $*$ | 3 | 2 |
| RCL | 2 | 2 |
| + | 6 | 2 |
| 2 | 2 | 2 |
| $=$ | 8 | 2 |
| $*$ | 8 | 2 |
| RCL | 2 | 2 |
| + | 16 | 2 |
| 1 | 1 | 2 |
| $=$ | 17 | 2 |

Indeed $p(2)=17$.

## Sample Executions

| instruction | Display | Memory |
| :--- | :--- | :--- |
| 0 | 0 | x |
| STO | 0 | 0 |
| 3 | 3 | 0 |
| $*$ | 3 | 0 |
| RCL | 0 | 0 |
| + | 0 | 0 |
| 2 | 2 | 0 |
| $=$ | 2 | 0 |
| $*$ | 2 | 0 |
| RCL | 0 | 0 |
| + | 2 | 0 |
| 1 | 1 | 0 |
| $=$ | 1 | 0 |

Indeed $p(0)=1$.

## In Summary

## The Programmer

- Writes a CAL-program by assuming a typical input.
- Writes where typical inputs are to be replaced by user inputs.
- Stores/writes and transmits.

10 \% substitute input here *
8
DIV
5
$+$
32
= \% see output in display

## The Bum

- Receives the program.
- Substitutes his inputs.
- Runs the program on his calculator line-by-line in that order.


