



CS620, IIT BOMBAY

An Introduction to Hybrid Systems Modeling

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Dynamical Systems

Dynamical System: A system whose **state** evolves with **time** governed by a fixed set of **rules** or **dynamics**.

- **state**: valuation to variables (discrete or continuous) of the system
- **time**: discrete or continuous
- **dynamics**: discrete, continuous, or hybrid

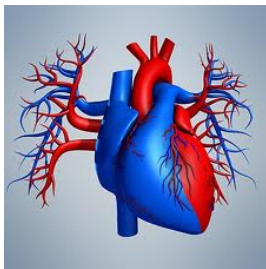
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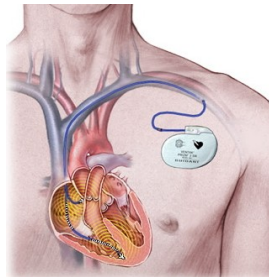
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Discrete System



Continuous System



Hybrid Systems.

Discrete Dynamical Systems

Continuous Dynamical Systems

Hybrid Dynamical Systems

Abstract Model for Dynamical Systems

Definition (State Transition Systems)

A state transition system is a tuple $\mathcal{T} = (S, S_0, \Sigma, \Delta)$ where:

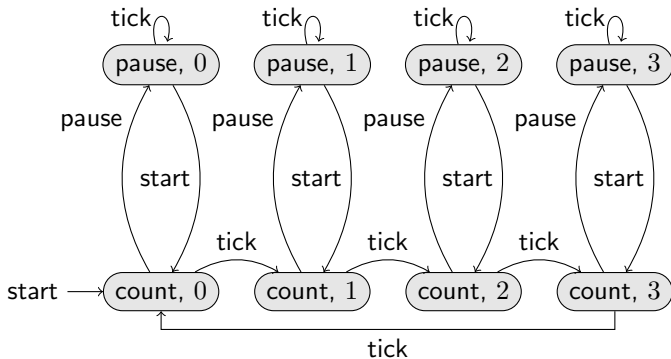
- S is a (potentially infinite) set of **states**;
- $S_0 \subseteq S$ is the set of **initial states**;
- Σ is a (potentially infinite) set of **actions**; and
- $\Delta \subseteq S \times \Sigma \times S$ is the **transition relation**;

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- Finite and countable state transition systems
- A **finite run** is a sequence

$$\langle s_0, a_1, s_1, s_2, s_2, \dots, s_n \rangle$$

such that $s_0 \in S_0$ and for all $0 \leq i < n$ we have that $(s_i, a_{i+1}, s_{i+1}) \in \Delta$.

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- **Reachability** and **Safe-Schedulability** problems

We need efficient computer-readable representations of infinite systems!

Extended Finite State Machines

- Let X be the set of **variables** (real-valued) of the system
- let $|X| = N$.
- A **valuation** ν of X is a function $\nu : X \rightarrow \mathbb{R}$.
- We consider a valuation as a point in \mathbb{R}^N equipped with **Euclidean Norm**.
- A **predicate** is defined simply as a subset of \mathbb{R}^N represented (non-linear) algebraic equations involving X .
 - Non-linear predicates, e.g. $x + 9.8 \sin(z) = 0$
 - Polyhedral predicates:

$$a_1x_1 + a_2x_2 + \dots + a_nx_n \sim k$$

where $a_i \in \mathbb{R}$, $x_i \in X$, and $\sim = \{<, \leq, =, \geq, >\}$.

- Octagonal predicates

$$x_i - x_j \sim k \text{ or } x_i \sim k$$

where $x_i, x_j \in X$, and $\sim = \{<, \leq, =, \geq, >\}$.

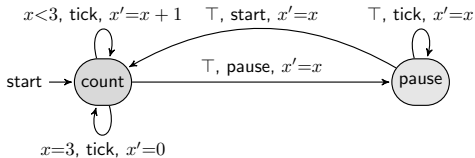
- Rectangular predicates

$$x_i \sim k$$

where $x_i \in X$, and $\sim = \{<, \leq, =, \geq, >\}$.

- Singular Predicates $x_i = c$.

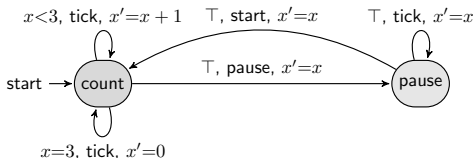
Extended Finite State Machines



Extended Finite State Machines (EFSMs):

- Finite state-transition systems coupled with a **finite set of variables**
- The valuation remains unchanged while system stays in a mode (state)
- The valuation changes during a transition when it **jumps** to the valuation governed by a predicate over $X \cup X'$ specified in the transition relation.
- Transitions are **guarded** by predicates over X
- **Mode invariants**
- **Initial state** and **valuation**

Extended Finite State Machines

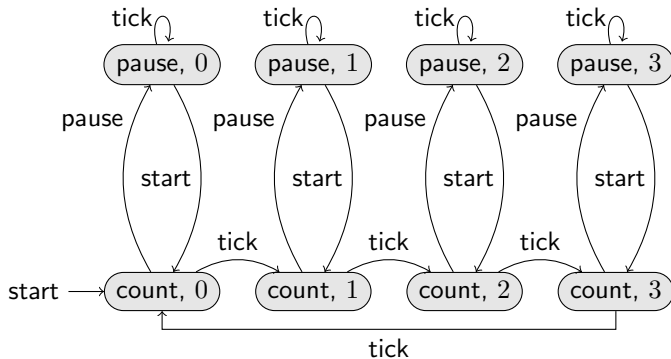
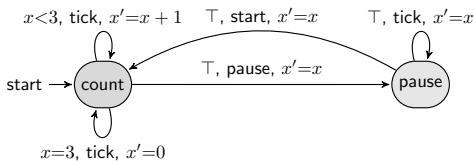


Definition (EFSM: Syntax)

An **extended finite state machine** is a tuple $\mathcal{M} = (M, M_0, \Sigma, X, \Delta, I, V_0)$ such that:

- M is a finite set of control **modes** including a distinguished initial set of control modes $M_0 \subseteq M$,
- Σ is a finite set of **actions**,
- X is a finite set of real-valued **variable**,
- $\Delta \subseteq M \times \text{pred}(X) \times \Sigma \times \text{pred}(X \cup X') \times M$ is the **transition relation**,
- $I : M \rightarrow \text{pred}(X)$ is the mode-invariant function, and
- $V_0 \in \text{pred}(X)$ is the set of initial valuations.

EFSM: Semantics



EFSM: Semantics

The **semantics** of an EFSM $\mathcal{M} = (M, M_0, \Sigma, X, \Delta, I, V_0)$ is given as a **state transition graph** $T^{\mathcal{M}} = (S^{\mathcal{M}}, S_0^{\mathcal{M}}, \Sigma^{\mathcal{M}}, \Delta^{\mathcal{M}})$ where

- $S^{\mathcal{M}} \subseteq (M \times \mathbb{R}^{|X|})$ is the set of **configurations** of \mathcal{M} such that for all $(m, \nu) \in S^{\mathcal{M}}$ we have that $\nu \in \llbracket I(m) \rrbracket$;
- $S_0^{\mathcal{M}} \subseteq S^{\mathcal{M}}$ is the set of **initial configurations** such that $(m, \nu) \in S^{\mathcal{M}}$ if $m \in M_0$ and $\nu \in V_0$;
- $\Sigma^{\mathcal{M}} = \Sigma$ is the set of **labels**;
- $\Delta^{\mathcal{M}} \subseteq S^{\mathcal{M}} \times \Sigma^{\mathcal{M}} \times S^{\mathcal{M}}$ is the set of **transitions** such that $((m, \nu), a, (m', \nu')) \in \Delta^{\mathcal{M}}$ if there exists a transition $\delta = (m, g, a, j, m') \in \Delta$ such that
 - **current valuation satisfies the guard** $\nu \in \llbracket g \rrbracket$;
 - **current and next valuations satisfy the jump constraint** $(\nu, \nu') \in \llbracket j \rrbracket$; and
 - **next valuation satisfies the invariant of the target mode** $\nu' \in \llbracket I(m') \rrbracket$.

Discrete Dynamical Systems

Continuous Dynamical Systems

Hybrid Dynamical Systems

Continuous Dynamical Systems

- A finite set of continuous variables,
- a set of ordinary differential equations (ODE) characterizing the flow of these variables as a function of time
 - $F : \dot{X} \rightarrow \text{pred}(X)$ where \dot{x} is the first derivative of x .
 - Higher-order derivatives can be written using first derivatives by introducing auxiliary variables, e.g. write $\ddot{\theta} + (g/\ell) \sin(\theta) = 0$ can be written as
$$\dot{\theta} = y \text{ and } \dot{y} = -(g/\ell) \sin(\theta).$$
- an initial valuation to the variables.

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Definition (Continuous Dynamical System)

A **continuous dynamical system** is a tuple $\mathcal{M} = (X, F, \nu_0)$ such that

- X is a **finite** set of real-valued **variable**,
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- A **run** or a **trajectory** of $\mathcal{M} = (X, F, \nu_0)$ is given as a solution to the differential equations $\dot{X} = F(X)$ with initial valuation ν_0 .
- Let a differentiable function $f : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^{|X|}$ be a solution to $\dot{X} = F(X)$ that provides the valuations of the variables as a function of time:

$$\begin{aligned}f(0) &= \nu_0 \\ \dot{f}(t) &= F(f(t)) \text{ for every } t \in \mathbb{R}_{\geq 0},\end{aligned}$$

where $\dot{f} : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^{|X|}$ is the time derivative of the function f .

- a run of a continuous dynamical system **may not exist or may not be unique!**

Existence and Uniqueness

Definition (Lipschitz-continuous Function)

We say that a function $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is Lipschitz-continuous if there exists a constant $K > 0$, called the Lipschitz constant, such that for all $x, y \in \mathbb{R}^n$ we have that $\|F(x) - F(y)\| < K\|x - y\|$.

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Theorem (Picard-Lindelöf Theorem)

If a function $F : \mathbb{R}^{|X|} \rightarrow \mathbb{R}^{|X|}$ is Lipschitz-continuous then the differential equation $\dot{X} = F(X)$ with initial valuation $\nu_0 \in \mathbb{R}^{|X|}$ has a unique solution $f : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^{|X|}$ for all $\nu_0 \in \mathbb{R}^{|X|}$.

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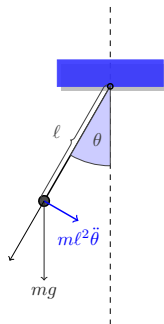
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Theorem (Stability wrt initial valuation)

Let F be a Lipschitz-continuous function with constant $K > 0$ and let $f : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^{|X|}$ and $f' : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^{|X|}$ be solutions to the differential equation $\dot{X} = F(X)$ with initial valuation $\nu_0 \in \mathbb{R}^{|X|}$ and $\nu'_0 \in \mathbb{R}^{|X|}$, respectively. Then, for all $t \in \mathbb{R}_{\geq 0}$ we have that $\|f(t) - f'(t)\| \leq \|\nu - \nu_0\| e^{Kt}$.

Example: Simple Pendulum



- Variables y and θ
- flow equations: $m\ell^2\ddot{\theta} = -mg\ell \sin(\theta)$, or

$$\dot{\theta} = y,$$

$$\dot{y} = -(g/\ell) \sin(\theta),$$

- initial valuations $(\theta, y) = (\theta_0, 0)$.

Simple Pendulum

- To analytically solve these equations, assume that initial angular displacement θ is small.
- hence $\sin(\theta) \approx \theta$.
- Now the equations simplify to

$$\dot{\theta} = y \text{ and } \dot{y} = -(g/\ell)\theta.$$

- The solution for these differential equations is

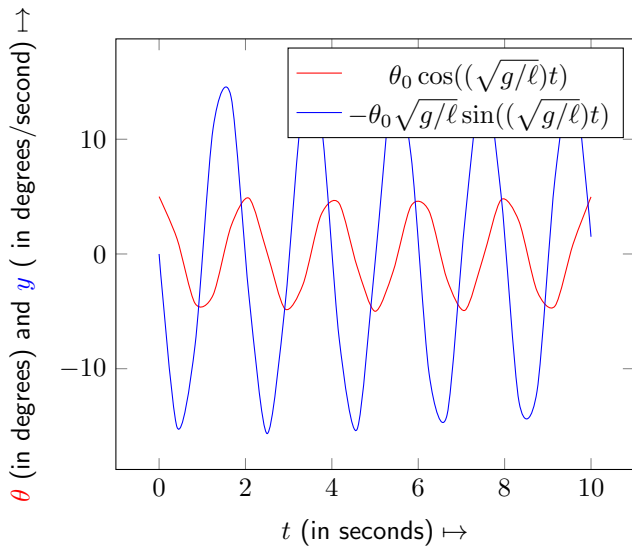
$$\begin{aligned}\theta(t) &= A \cos(Kt) + B \sin(Kt) \\ y(t) &= -AK \sin(Kt) + BK \cos(Kt),\end{aligned}$$

where $K = \sqrt{g/\ell}$.

- Substituting $\theta(0) = \theta_0$ and $y(0) = 0$ from the initial valuation, we get that $A = \theta_0$ and $B = 0$.
- The unique run of the pendulum system can be given as the function $f : \mathbb{R}_{\geq 0} \rightarrow \{\theta, y\}$ as

$$t \mapsto (\theta_0 \cos(Kt), -\theta_0 K \sin(Kt)).$$

Pendulum Motion



(b)

Discrete Dynamical Systems

Continuous Dynamical Systems

Hybrid Dynamical Systems

Another example: Bouncing Ball

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- $x'_1 = x_1$ and $x'_2 = -cx_2$ where c is **Restitution coefficient**.

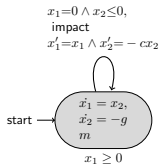
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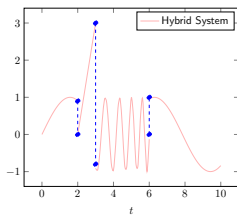
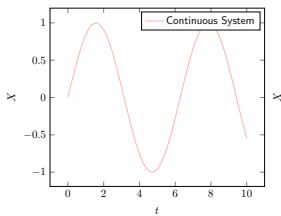
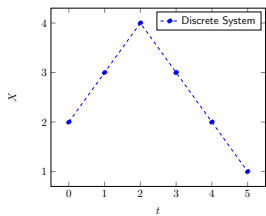
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Discrete, Continuous, and Hybrid Systems

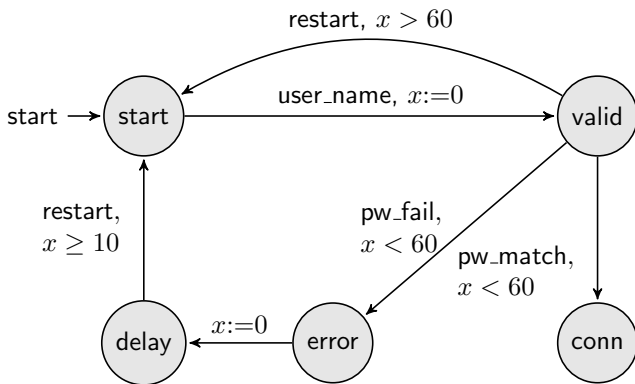


Hybrid Automata: Syntax

Some examples:

- Two leaking-water tanks systems
- Water-level monitor with delayed switch
- A leaking gas-burner
- Green scheduling with lower dwell-time requirements
- Light-bulb with three modes- dim, bright, and off.
- Job-shop scheduling problem

Login Protocol



Hybrid Automata: Syntax

Definition (HA: Syntax)

A hybrid automaton is a tuple $\mathcal{H} = (M, M_0, \Sigma, X, \Delta, I, F, V_0)$ where:

- M is a finite set of control **modes** including a distinguished initial set of control modes $M_0 \subseteq M$,
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- X is a finite set of real-valued **variable**,
- $\Delta \subseteq M \times \text{pred}(X) \times \Sigma \times \text{pred}(X \cup X') \times M$ is the **transition relation**,
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- A **configuration** (m, ν) and a **timed action** (t, a)
- A **transition** $((m, \nu), (t, a), (m', \nu'))$
 - solve flow ODE of mode m with ν as the starting state $\nu \oplus_{F(m)} t$.
 - invariant, guard, and jump conditions.
- A **run** or **execution** is a sequence of transitions

$$(m_0, \nu_0), (t_1, a_1), (m_1, \nu_1), (t_2, a_2) \dots$$

Hybrid Automata: Semantics

Definition (HA: Semantics)

The semantics of a HA $\mathcal{H} = (M, M_0, \Sigma, X, \Delta, I, F, V_0)$ is given as a state transition graph $T^{\mathcal{H}} = (S^{\mathcal{H}}, S_0^{\mathcal{H}}, \Sigma^{\mathcal{H}}, \Delta^{\mathcal{H}})$ where

- $S^{\mathcal{H}} \subseteq (M \times \mathbb{R}^{|X|})$ is the set of configurations of \mathcal{H} such that for all $(m, \nu) \in S^{\mathcal{H}}$ we have that $\nu \in \llbracket I(m) \rrbracket$;
- $S_0^{\mathcal{H}} \subseteq S^{\mathcal{H}}$ s.t. $(m, \nu) \in S_0^{\mathcal{H}}$ if $m \in M_0$ and $\nu \in V_0$;
- $\Sigma^{\mathcal{H}} = \mathbb{R}_{\geq 0} \times \Sigma$ is the set of labels;
- $\Delta^{\mathcal{H}} \subseteq S^{\mathcal{H}} \times \Sigma^{\mathcal{H}} \times S^{\mathcal{H}}$ is the set of transitions such that $((m, \nu), (t, a), (m', \nu')) \in \Delta^{\mathcal{H}}$ if there exists a transition $\delta = (m, g, a, j, m') \in \Delta$ such that
 - $(\nu \oplus_{F(m)} t) \in \llbracket g \rrbracket$;
 - $(\nu \oplus_{F(m)} \tau) \in \llbracket I(m) \rrbracket$ for all $\tau \in [0, t]$;
 - $\nu' \in (\nu \oplus_{F(m)} t)[j]$; and
 - $\nu' \in \llbracket I(m') \rrbracket$.