

# *The Design and Implementation of Gnu Compiler Collection*

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*Part 1*

## *Outline*

# Outline

- Open Source Vs. Free Software
- An Introduction to Compilation
- An overview of Gnu Compiler Collection
- Our work in GCC
- Conclusions

*Part 2*

## *Open Source Vs. Free Software*

# Open Source Vs. Free Software

- The open source initiative: (<http://www.opensource.org/>)  
Emphasis on development methodology
- The Free Software Foundation: (<http://www.fsf.org/>)  
Emphasis on freedom of the user
- In some cases, open source software has restricted user freedom



# Open Source Vs. Free Software

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- The Free Software Foundation: (<http://www.fsf.org/>)  
Emphasis on freedom of the user
- In some cases, open source software has restricted user freedom
  - ▶ May cooperate with restrictive systems
  - ▶ eg. DRM (Digital Rights Management)



# Open Source and Free Software Development Model

- The Cathedral and the Bazaar  
Eric S Raymond, 1997.
- Cathedral: Total Centralized Control  
Design, implement, test, release
- Bazaar: Total Decentralization  
Release early, release often, let users fix bugs

## The Bazaar Approach

Release early, release often, let users fix bugs

- Brooks' law (The Mythical Man Month, 1975)

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OR  
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- “Given enough eyeballs, all bugs are shallow”.  
Code errors, logical errors, and architectural errors.



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OR  
12 persons working for 1 month?
- Bazaar approach believes that the two somewhat equivalent in internet-based distributed development.
- “Given enough eyeballs, all bugs are shallow”.  
Code errors, logical errors, and architectural errors.

*A combination of the two seems more sensible*



*Part 3*

## *Introduction to Compilation*

# Implementation Mechanisms

Source Program



Translator

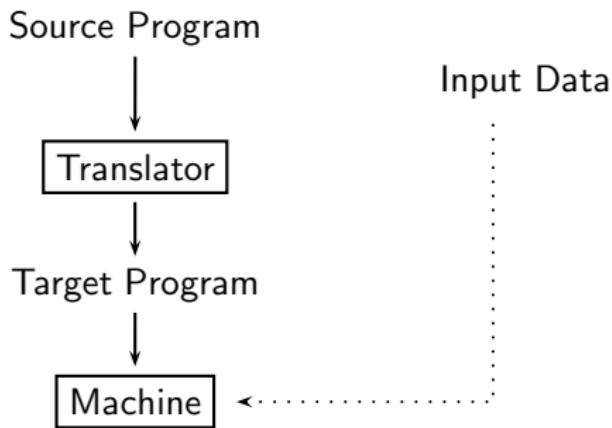


Target Program

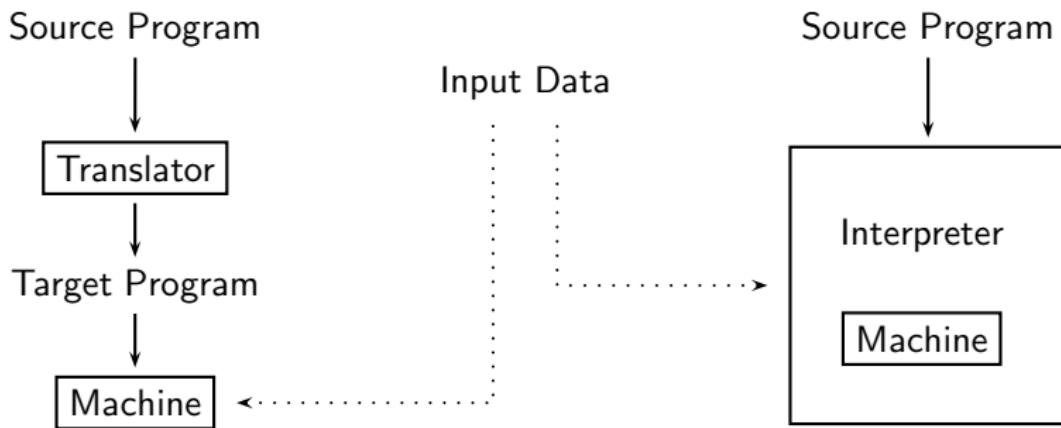


Machine

# Implementation Mechanisms



# Implementation Mechanisms



# Implementation Mechanisms as “Bridges”

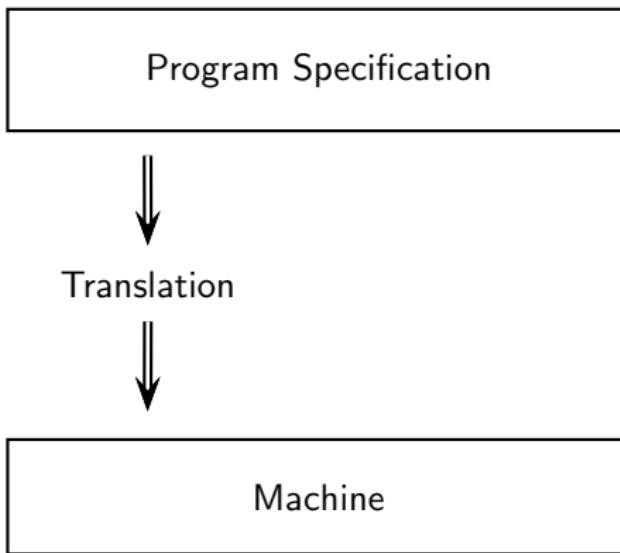
- “Gap” between the “levels” of program specification and execution

Program Specification

Machine

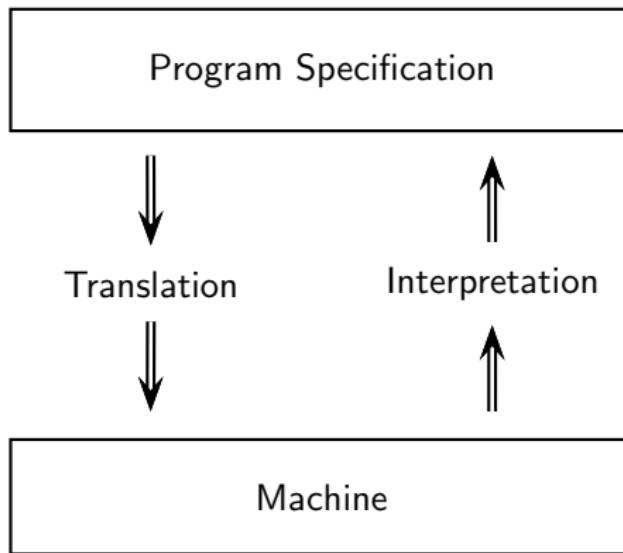
## Implementation Mechanisms as “Bridges”

- “Gap” between the “levels” of program specification and execution



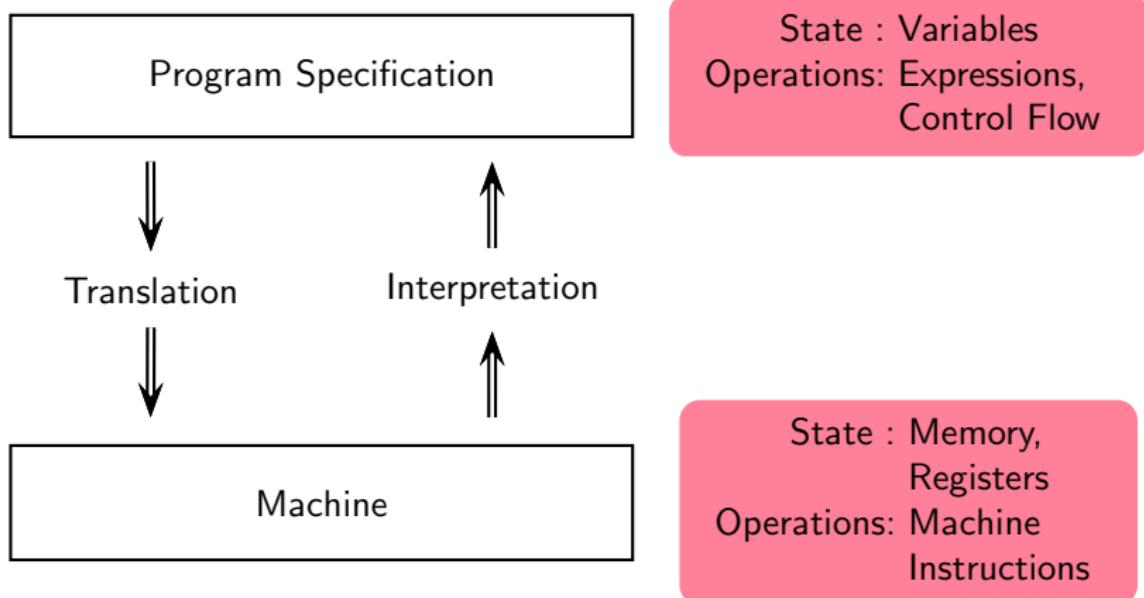
## Implementation Mechanisms as “Bridges”

- “Gap” between the “levels” of program specification and execution



# Implementation Mechanisms as “Bridges”

- “Gap” between the “levels” of program specification and execution



## High and Low Level Abstractions

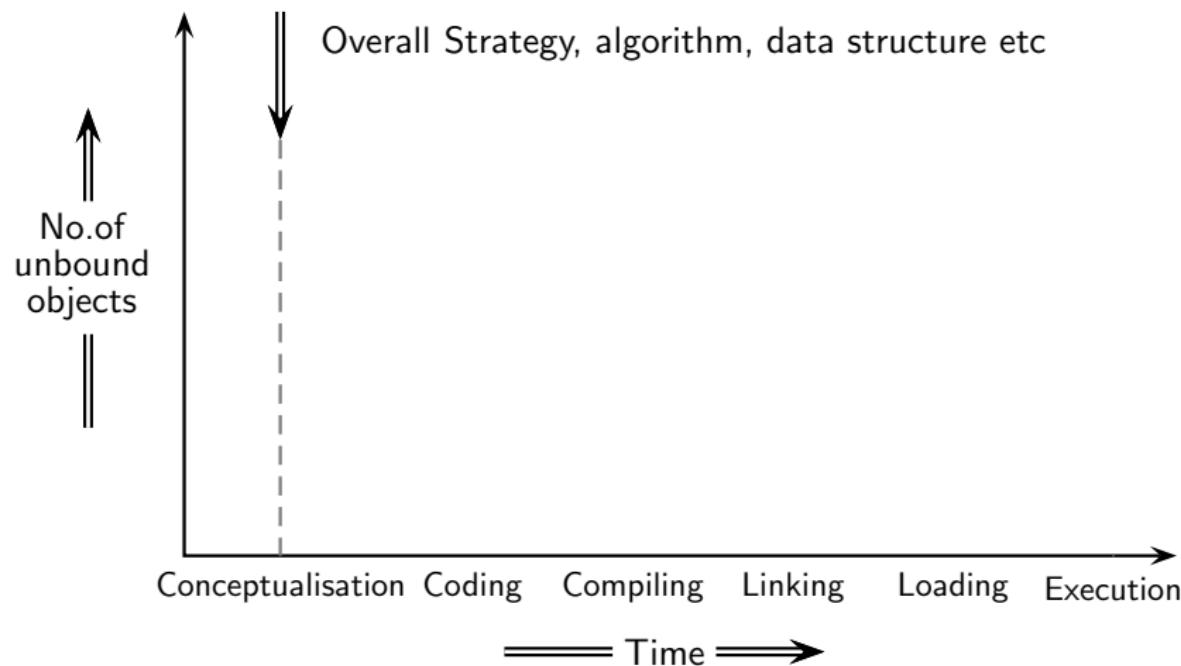
- A source statement

```
a = b < 10 ? b : c;
```

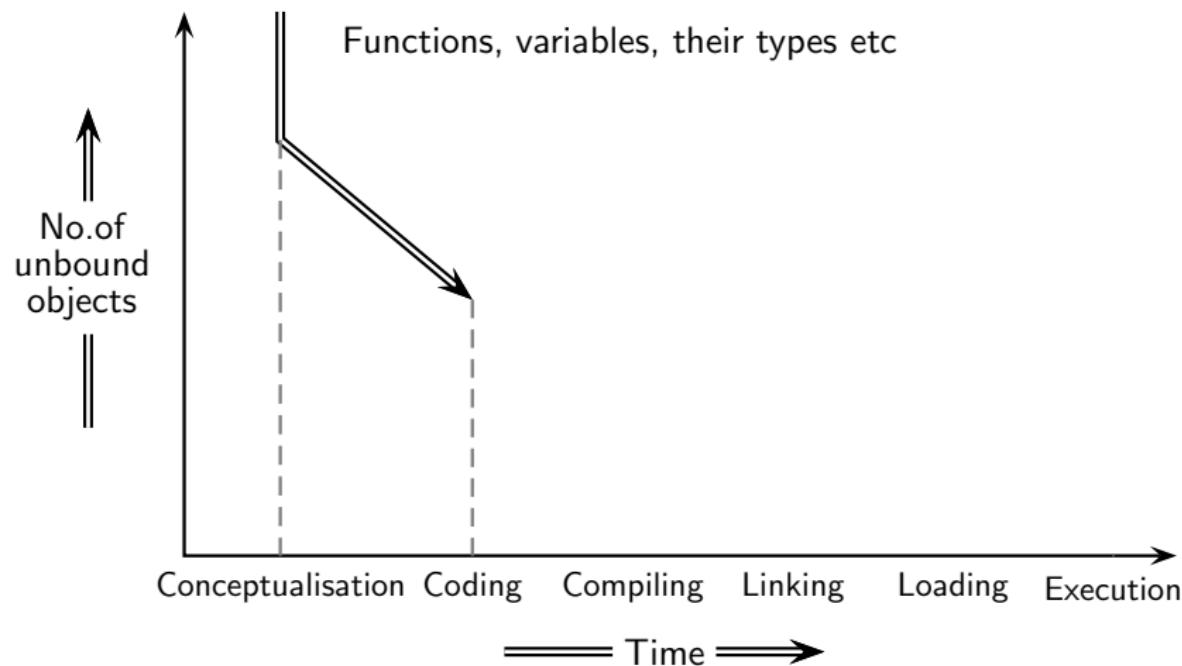
- Spim assembly equivalent

```
lw    $t0, 4($fp)      #      $t0 <- b
slti $t0, $t0, 10      #      $t0 <- $t0 < b
not  $t0, $t0          #      $t0 <- ! $t0
bgtz $t0, L0:         #      if $t0 >= 0 goto L0:
lw    $t0, 4($fp)      #      $t0 <- b
b    L1:                #      goto L1:
L0:  lw    $t0, 8($fp)  # L0: $t0 <- c
L1:  sw    0($fp), $t0  # L1: a <- $t0
```

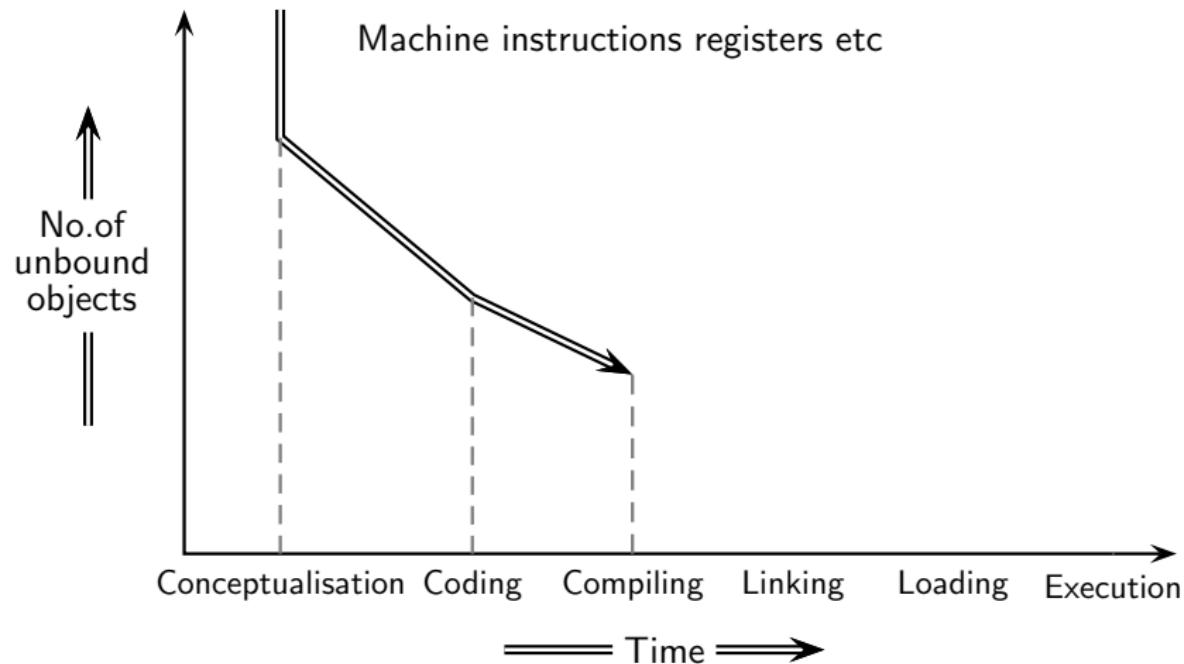
# Binding



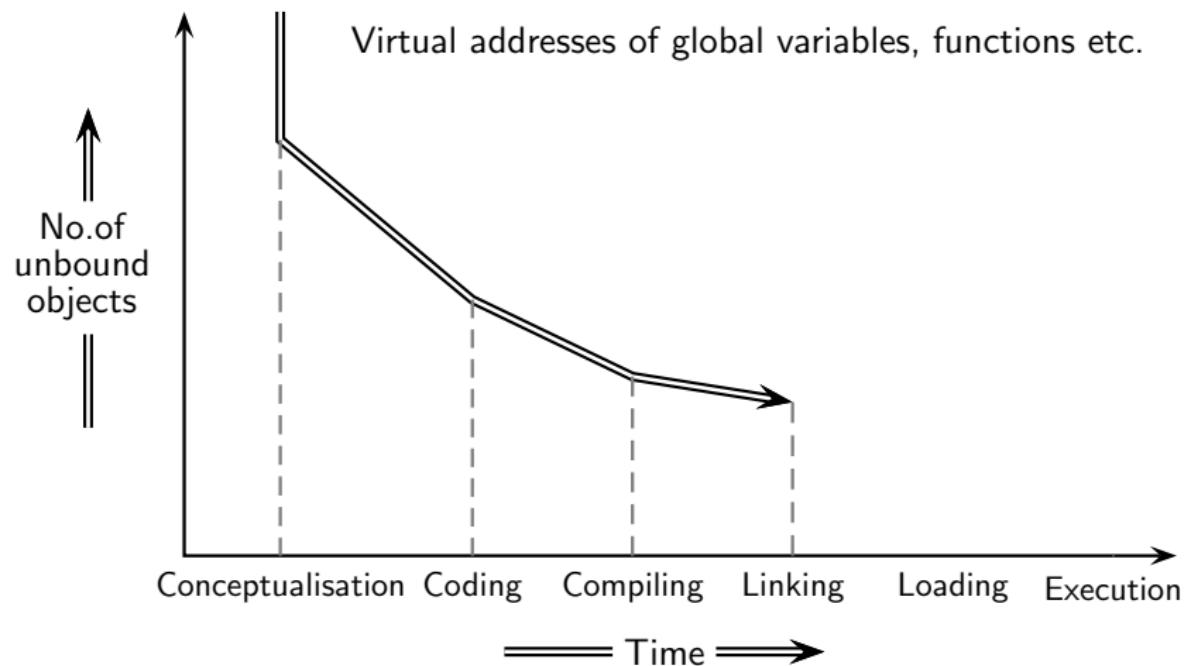
# Binding



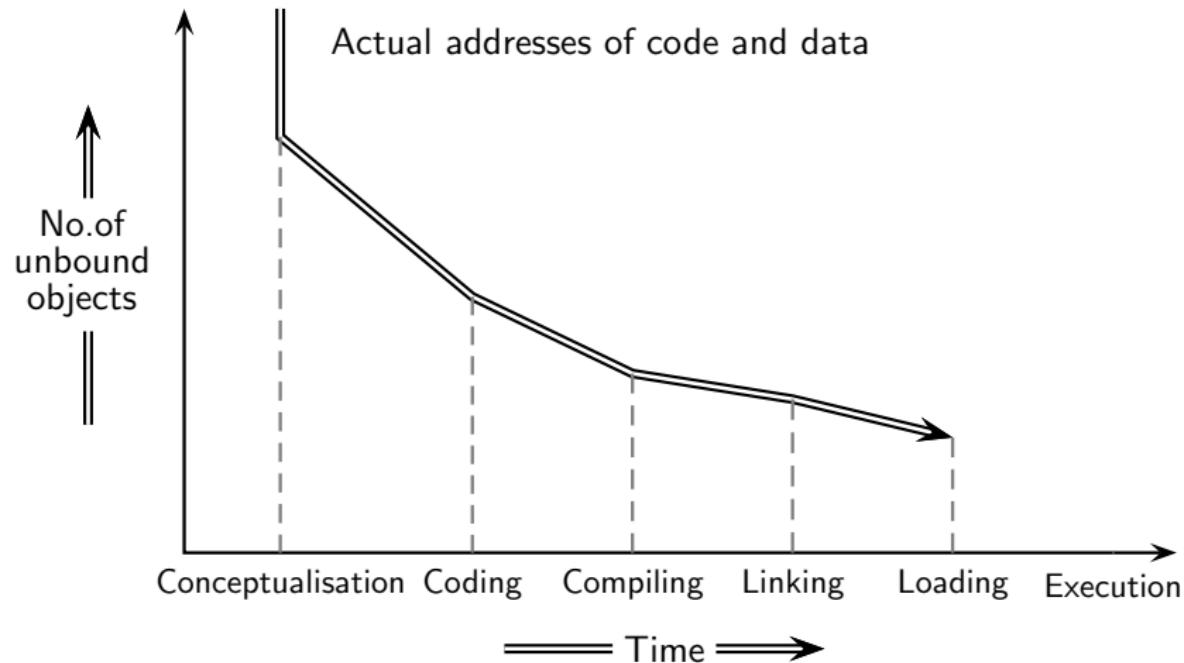
# Binding



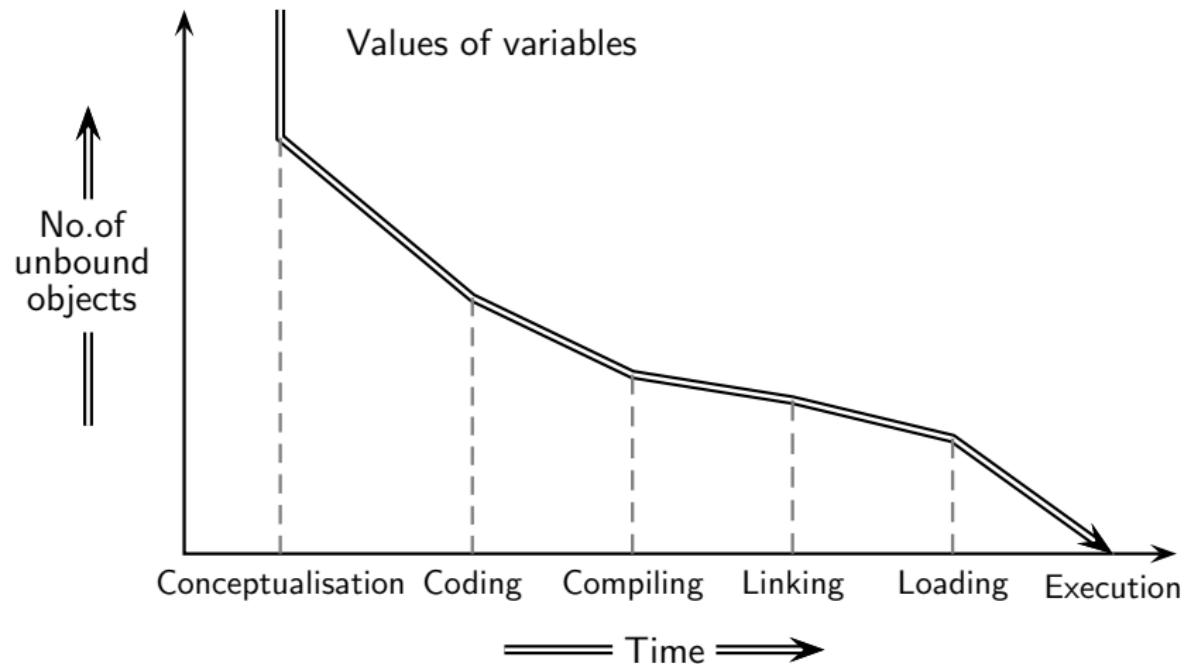
# Binding



# Binding



# Binding



# Implementation Mechanisms

- Translation = Analysis + Synthesis
- Interpretation = Analysis + Execution

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- Translation      Instructions       $\Rightarrow$       Equivalent Instructions

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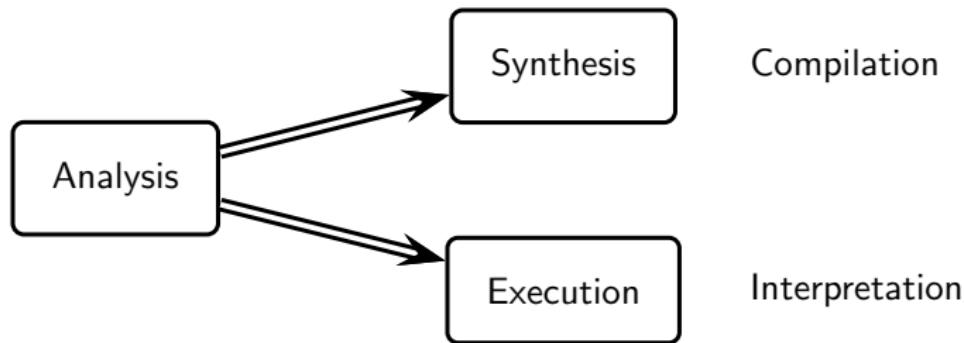
- Translation = Analysis + Synthesis

- Interpretation = Analysis + Execution

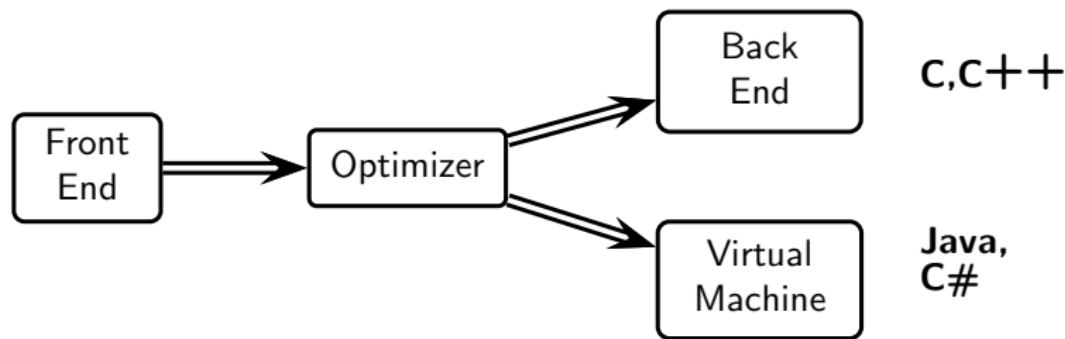
- Translation      Instructions       $\Rightarrow$       Equivalent Instructions

- Interpretation      Instructions       $\Rightarrow$       Actions Implied by Instructions

# Language Implementation Models



# Language Processor Models

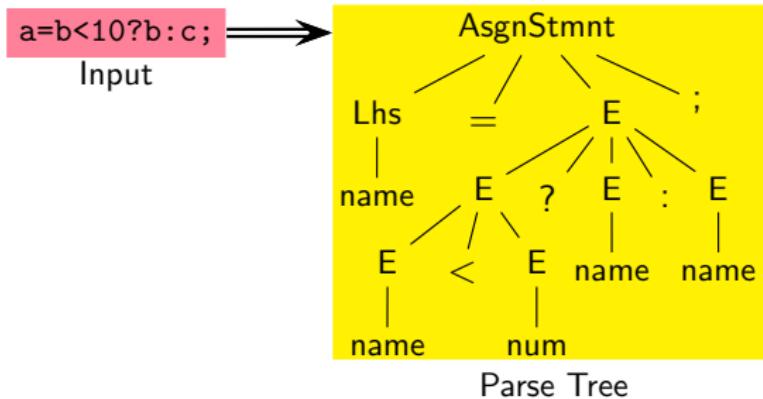


## Translation Sequence in Our Example: Parsing

```
a=b<10?b:c;
```

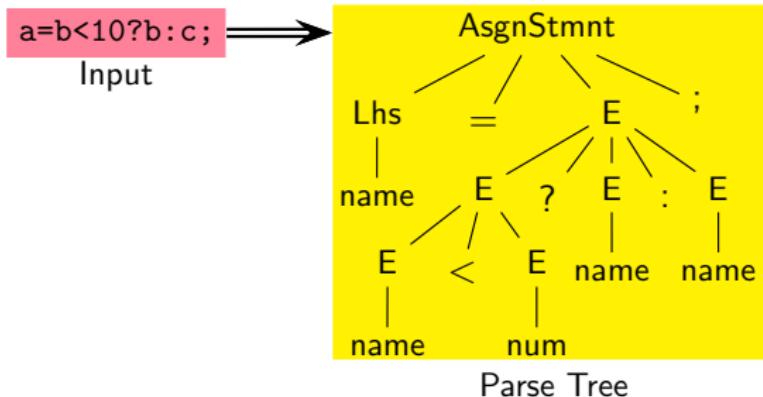
Input

## Translation Sequence in Our Example: Parsing



Issues:

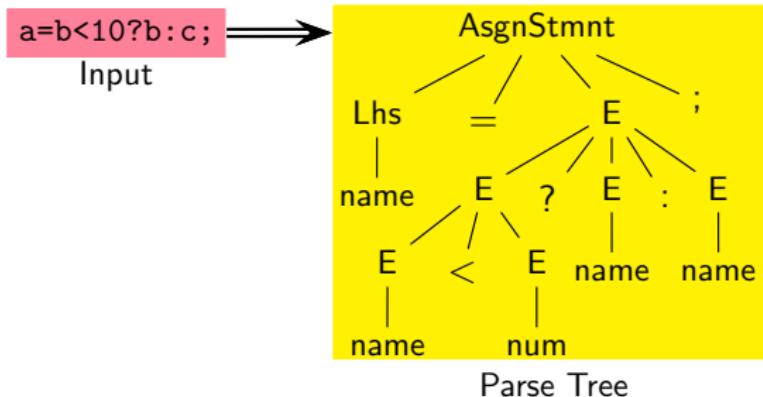
## Translation Sequence in Our Example: Parsing



### Issues:

- Grammar rules, terminals, non-terminals

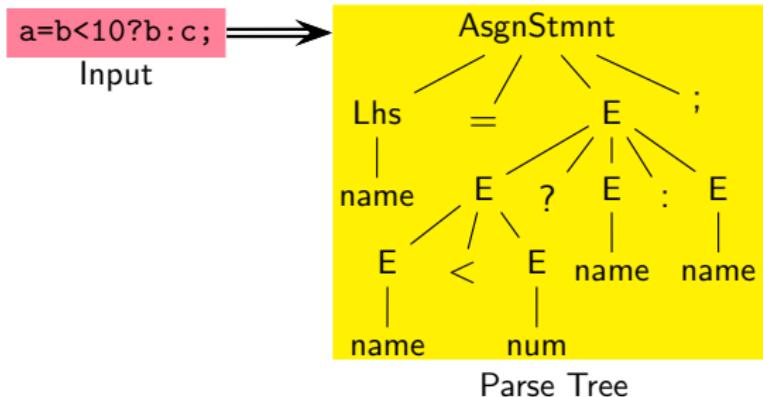
## Translation Sequence in Our Example: Parsing



### Issues:

- Grammar rules, terminals, non-terminals
- Order of application of grammar rules  
eg. is it  $(a = b < 10?)$  followed by  $(b:c)$ ?

## Translation Sequence in Our Example: Parsing



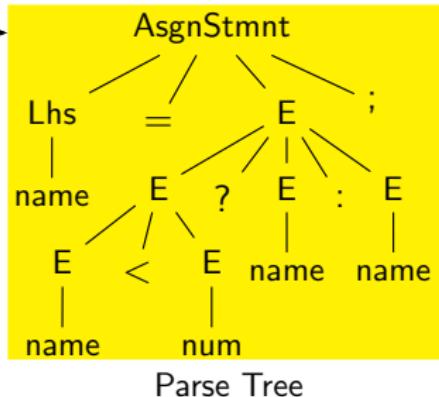
### Issues:

- Grammar rules, terminals, non-terminals
- Order of application of grammar rules
  - eg. is it ( $a = b < 10?$ ) followed by ( $b:c$ )?
- Values of terminal symbols
  - eg. string “10” vs. integer number 10.

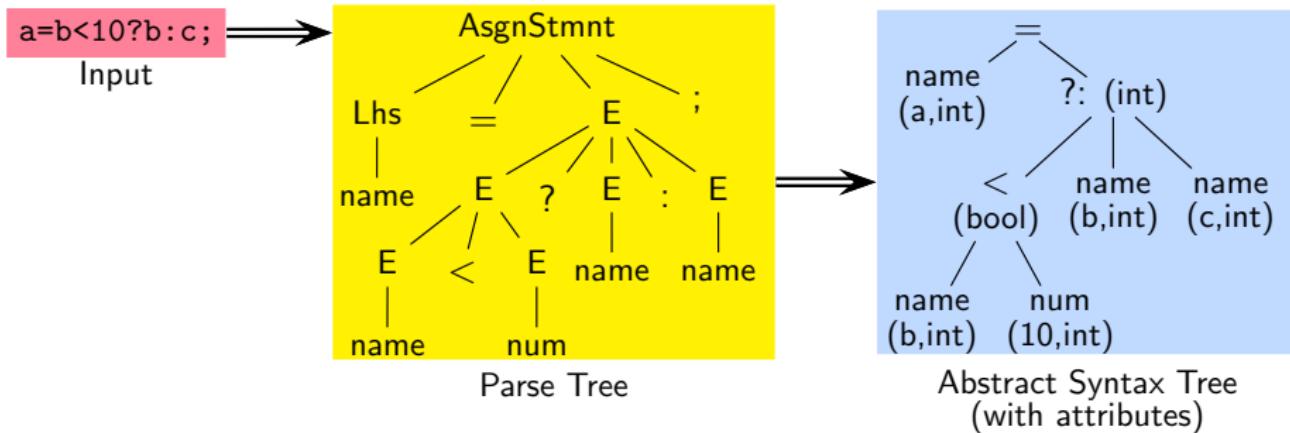


## Translation Sequence in Our Example: Semantic Analysis

a=b<10?b:c;  
Input

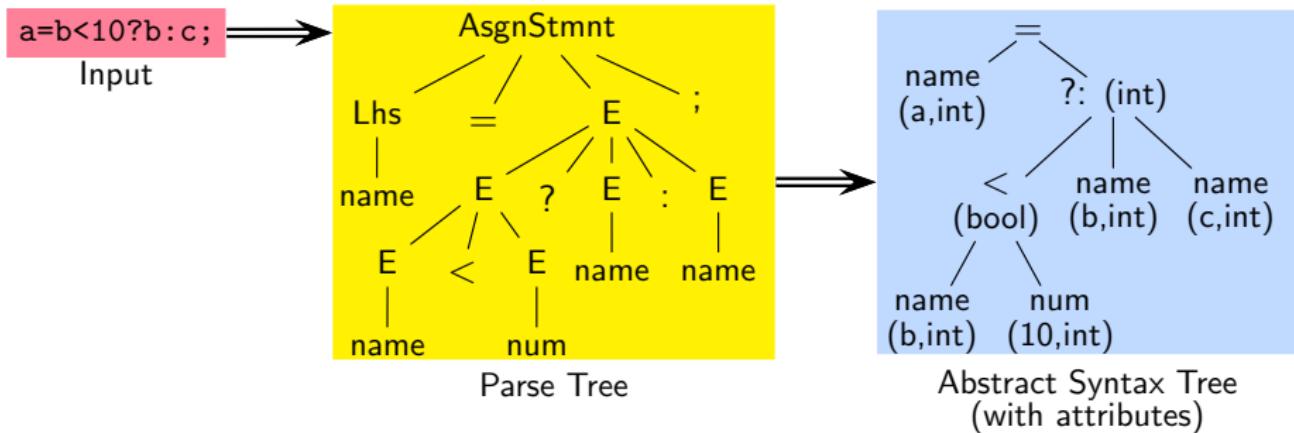


## Translation Sequence in Our Example: Semantic Analysis



Issues:

## Translation Sequence in Our Example: Semantic Analysis

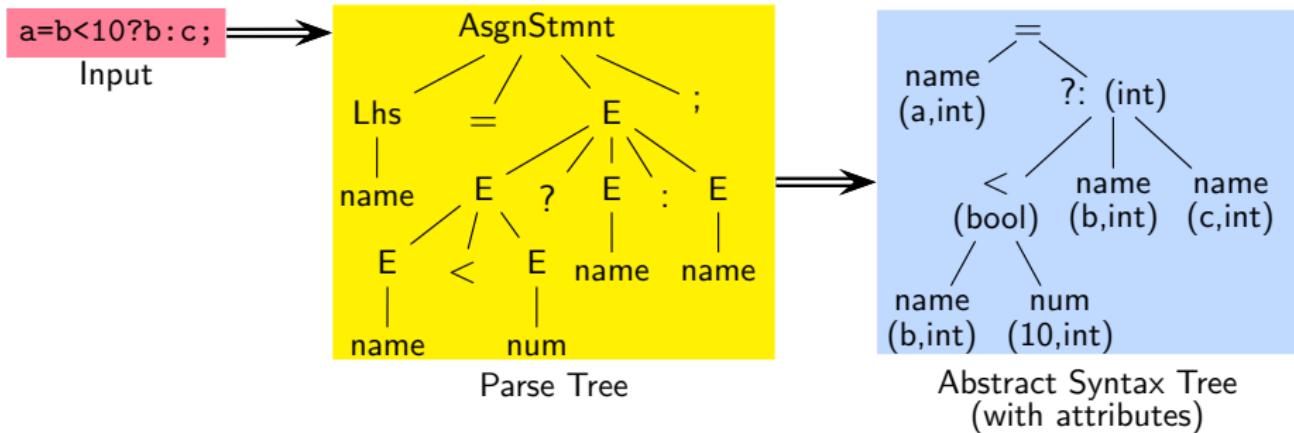


Issues:

- Symbol tables

Have variables been declared? What are their types?  
What is their scope?

## Translation Sequence in Our Example: Semantic Analysis



### Issues:

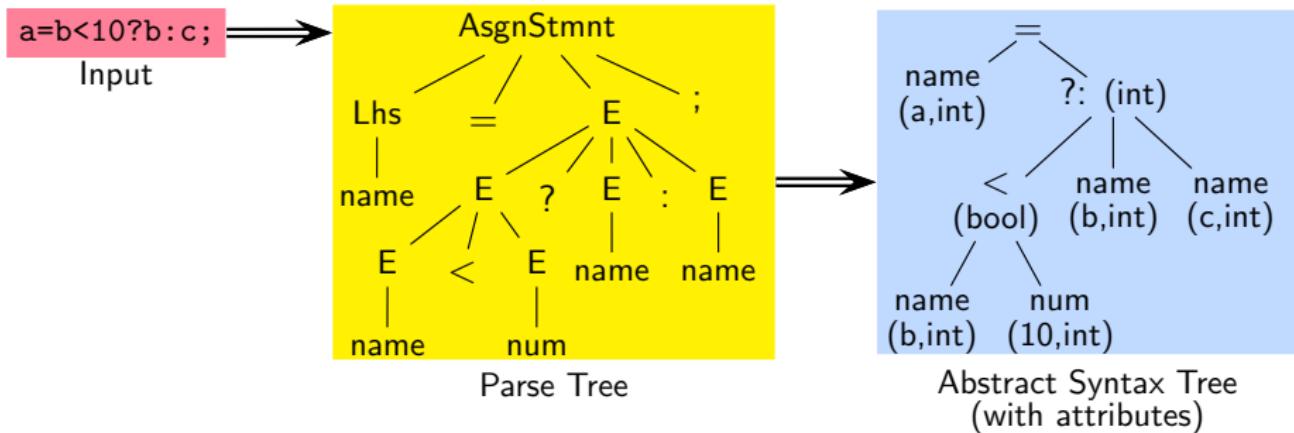
- Symbol tables
 

Have variables been declared? What are their types?  
What is their scope?
- Type consistency of operators and operands
 

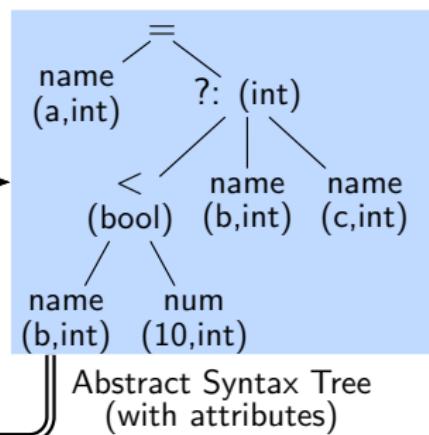
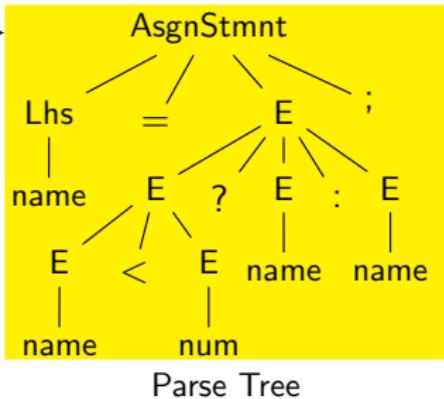
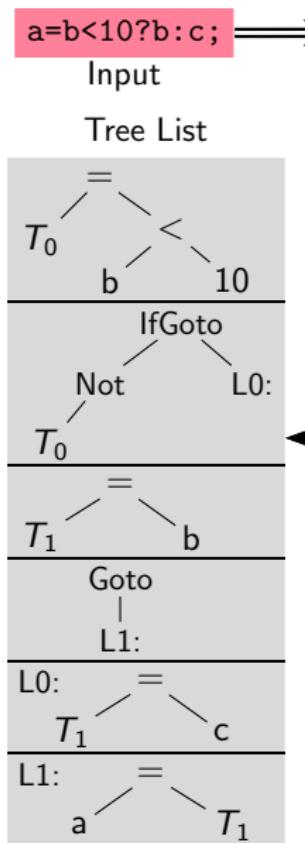
The result of computing `b<10?` is bool and not int



## Translation Sequence in Our Example: IR Generation



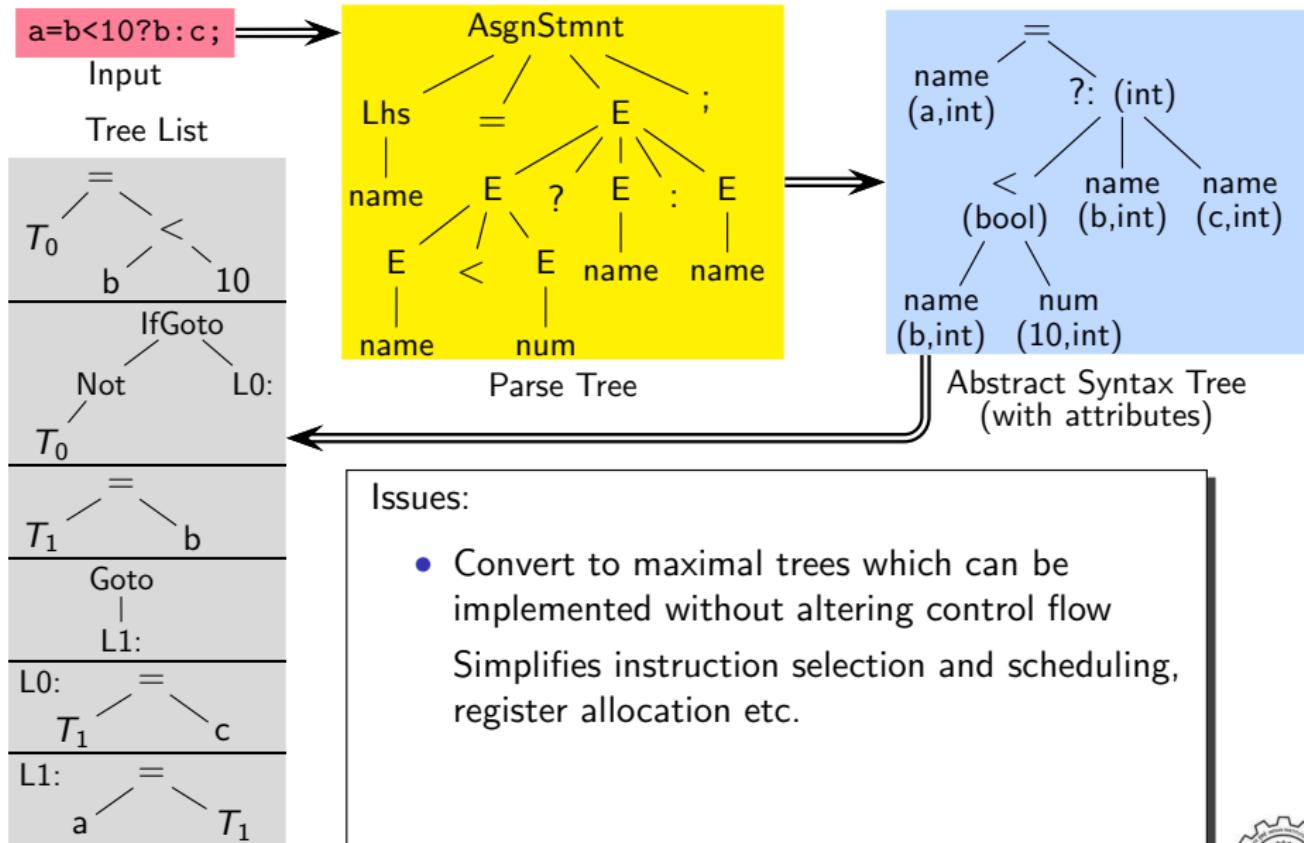
## Translation Sequence in Our Example: IR Generation



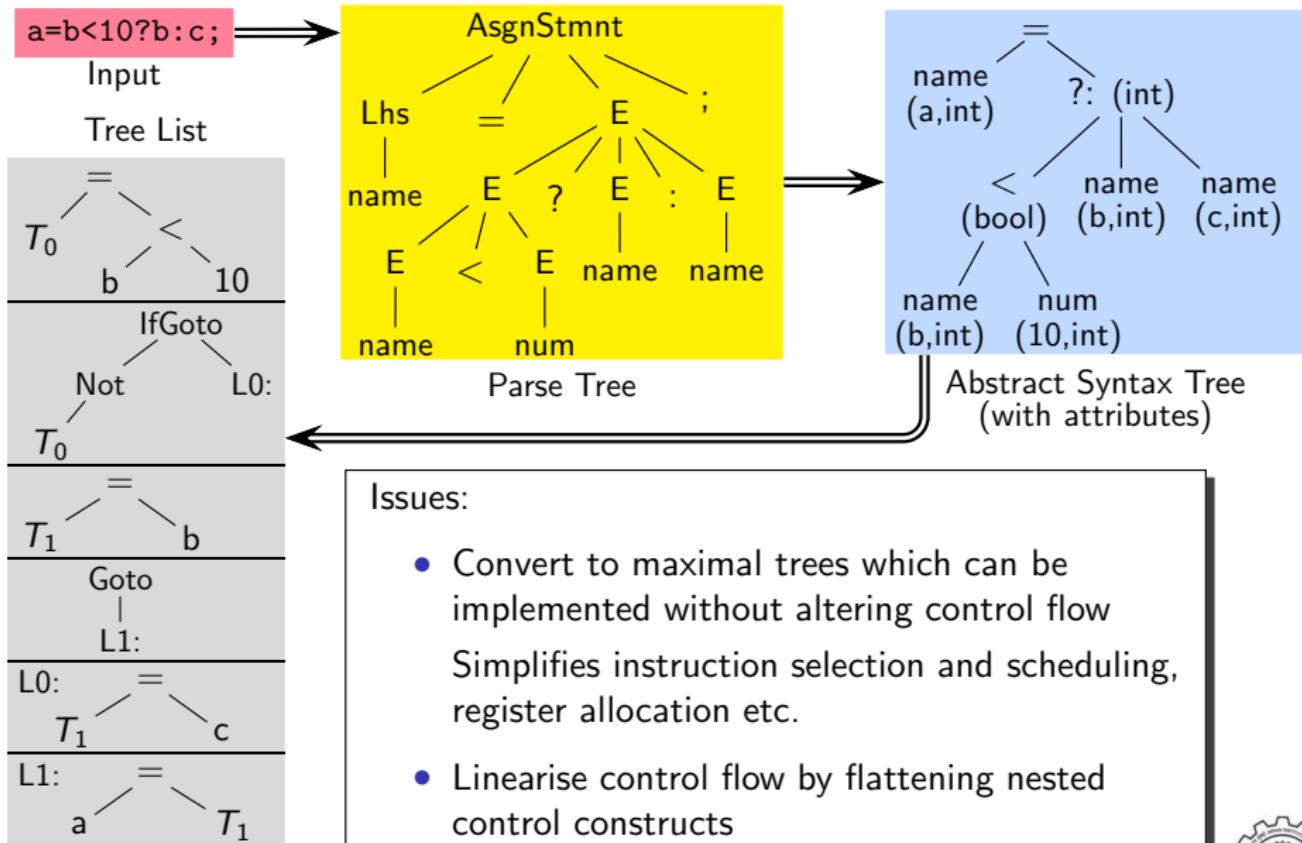
Issues:



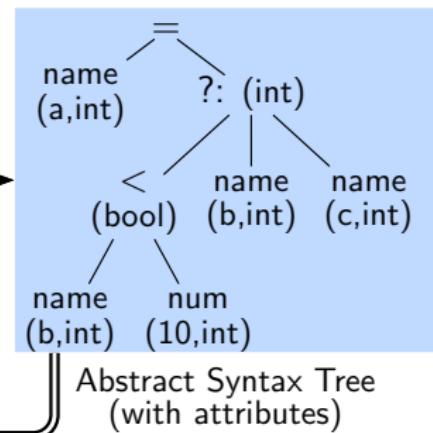
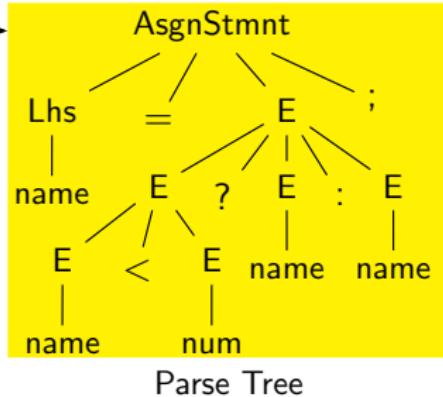
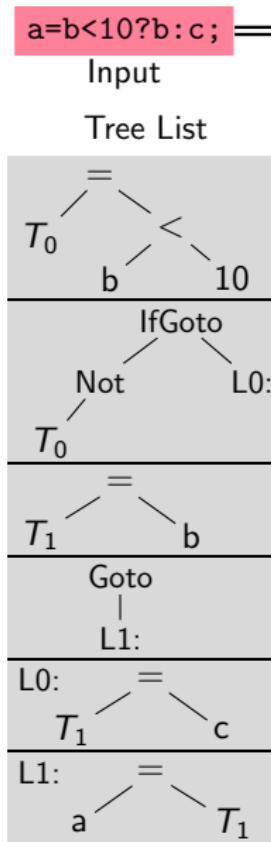
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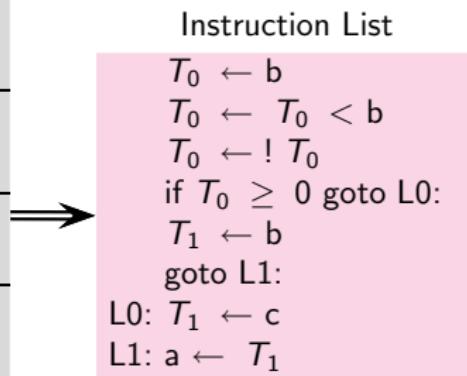
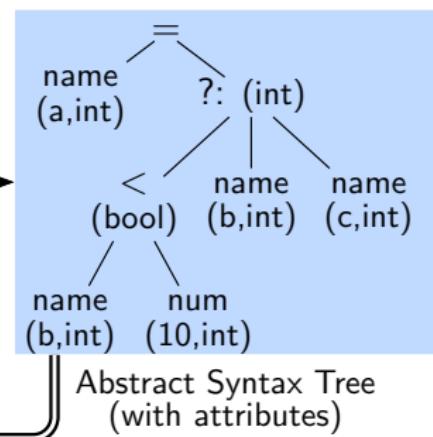
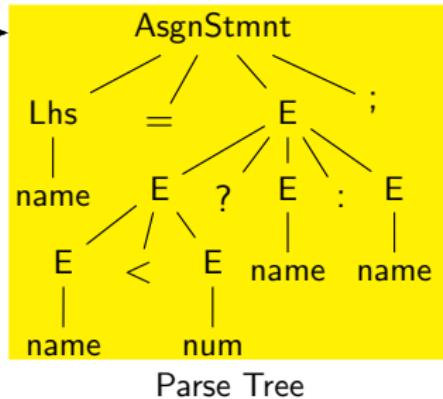
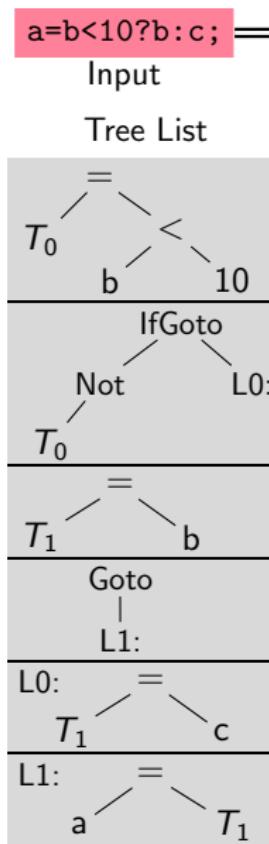
# Translation Sequence in Our Example: Instruction Selection



Abstract Syntax Tree  
(with attributes)



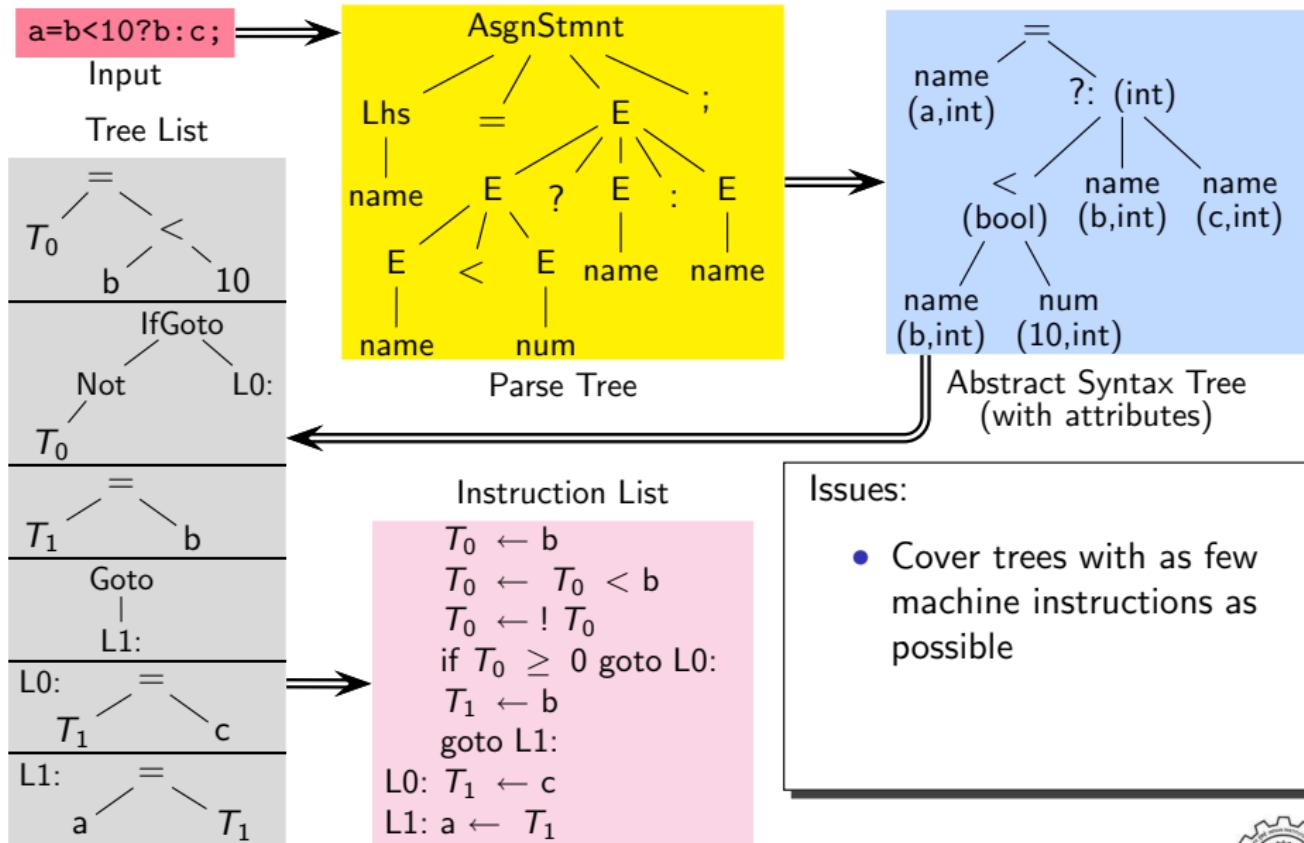
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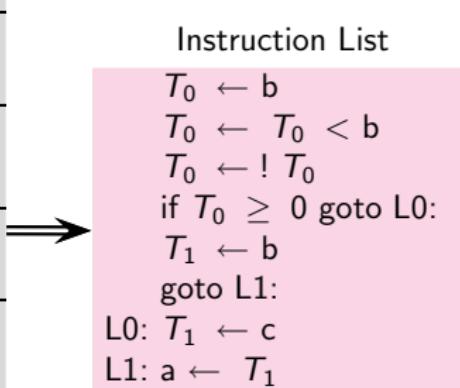
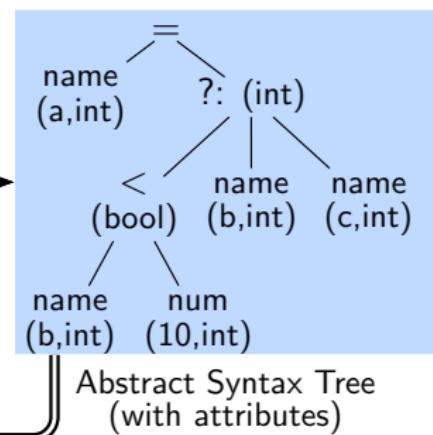
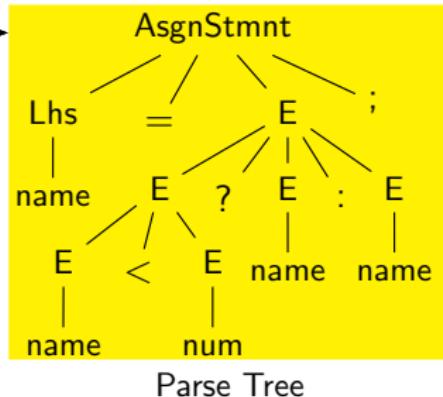
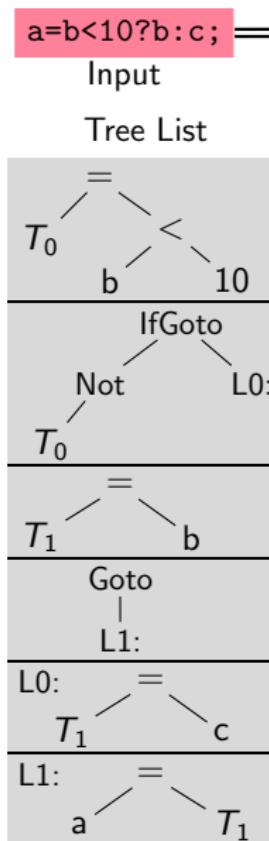
**Issues:**



## Translation Sequence in Our Example: Instruction Selection



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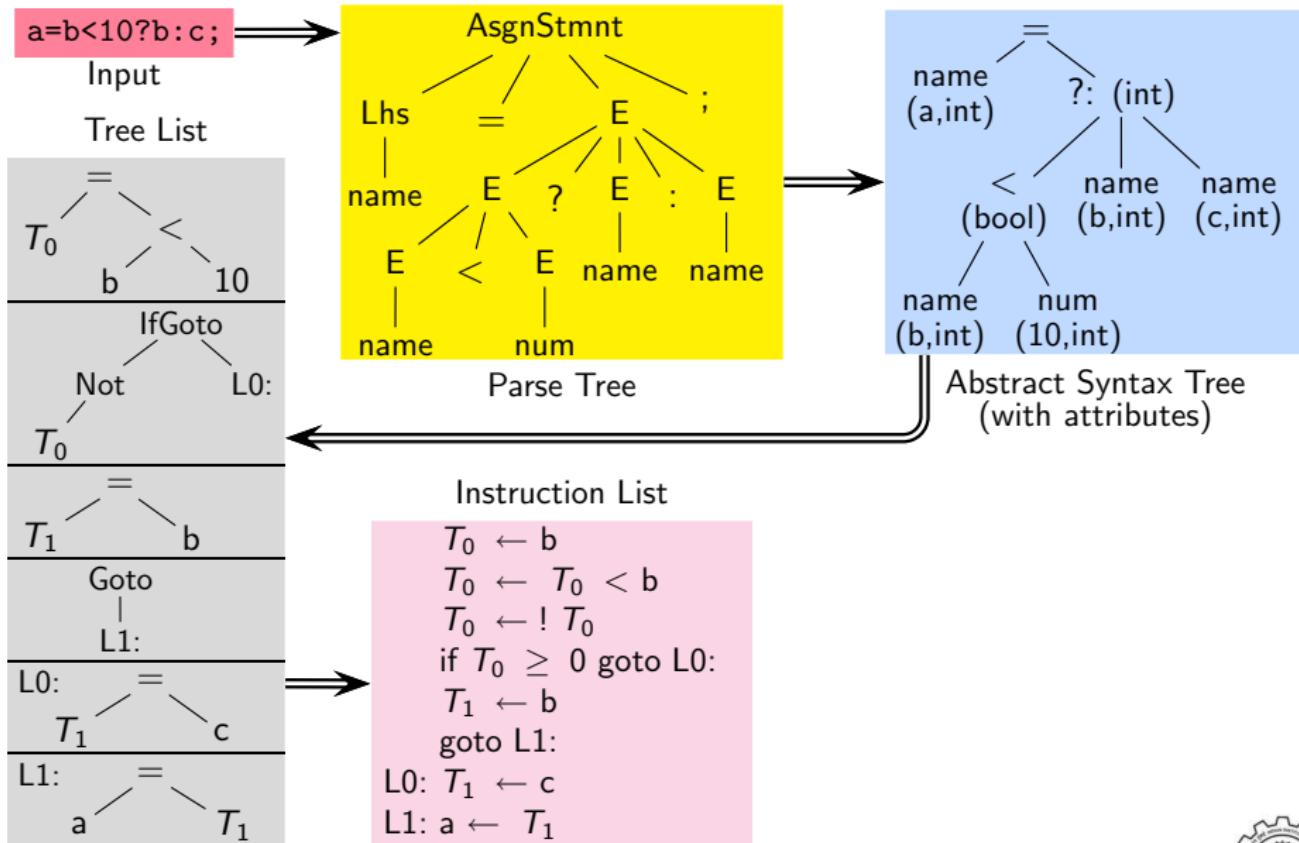


## Issues:

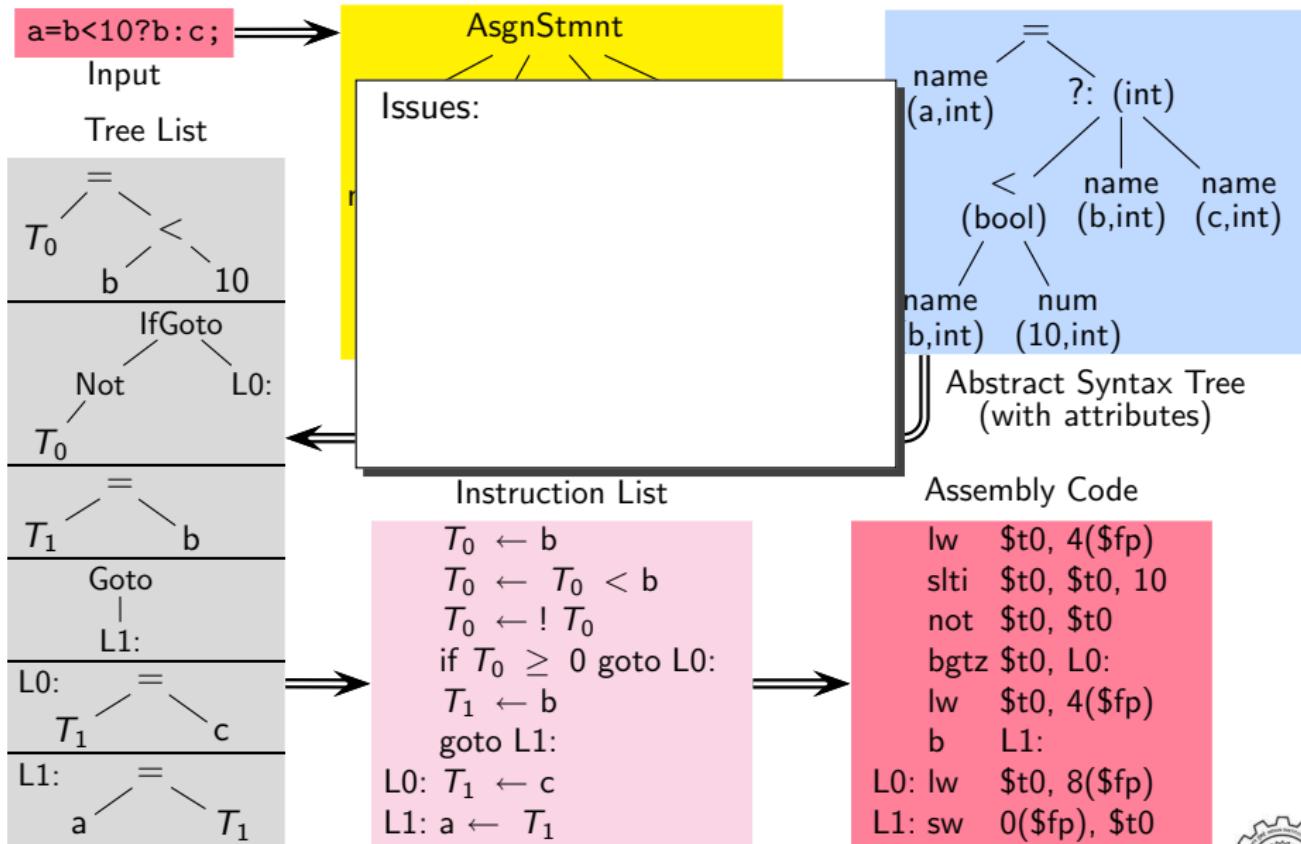
- Cover trees with as few machine instructions as possible
- Use temporaries and local registers



# Translation Sequence in Our Example: Emitting Instructions



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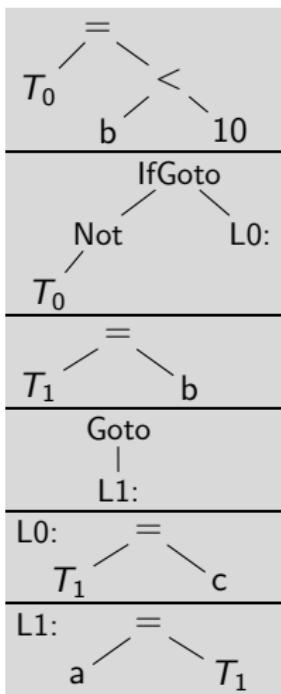


# Translation Sequence in Our Example: Emitting Instructions

a=b<10?b:c;

Input

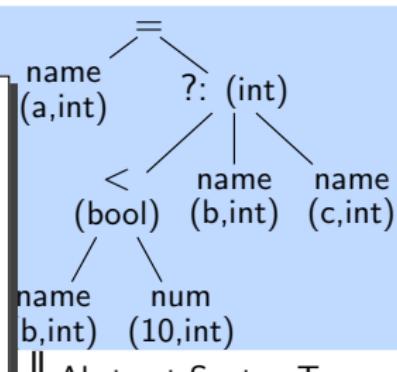
Tree List



AsgnStmnt

Issues:

- Offsets of variables in the stack frame



Abstract Syntax Tree  
(with attributes)

Instruction List

```

 $T_0 \leftarrow b$ 
 $T_0 \leftarrow T_0 < b$ 
 $T_0 \leftarrow !T_0$ 
if  $T_0 \geq 0$  goto L0:
 $T_1 \leftarrow b$ 
goto L1:
L0:  $T_1 \leftarrow c$ 
L1:  $a \leftarrow T_1$ 

```

Assembly Code

```

lw    $t0, 4($fp)
slti $t0, $t0, 10
not $t0, $t0
bgtz $t0, L0:
lw    $t0, 4($fp)
b    L1:
L0: lw    $t0, 8($fp)
L1: sw    0($fp), $t0

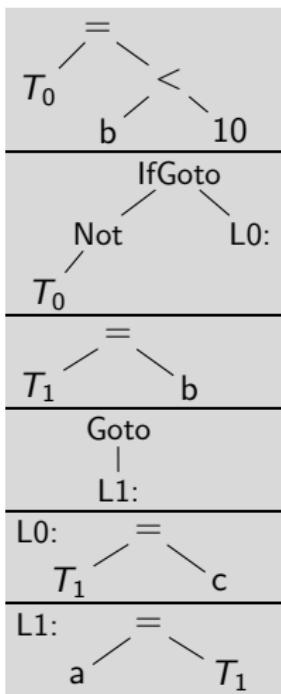
```

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Input

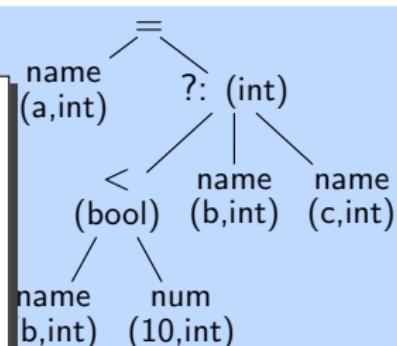
Tree List



AsgnStmnt

Issues:

- Offsets of variables in the stack frame
- Actual register numbers and assembly mnemonics



Abstract Syntax Tree  
(with attributes)

Instruction List

```

T0 ← b
T0 ← T0 < b
T0 ← ! T0
if T0 ≥ 0 goto L0:
T1 ← b
goto L1:
L0: T1 ← c
L1: a ← T1
  
```

Assembly Code

```

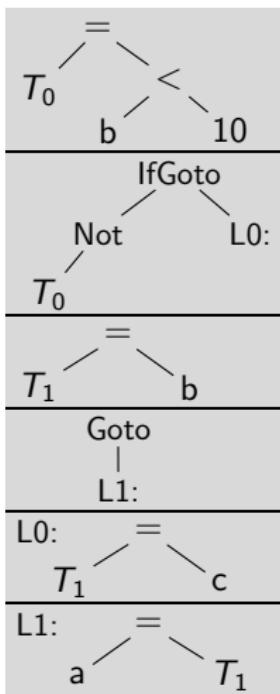
lw $t0, 4($fp)
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bgtz $t0, L0:
lw $t0, 4($fp)
b L1:
L0: lw $t0, 8($fp)
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```

# Translation Sequence in Our Example: Emitting Instructions

a=b<10?b:c;

Input

Tree List



AsgnStmnt

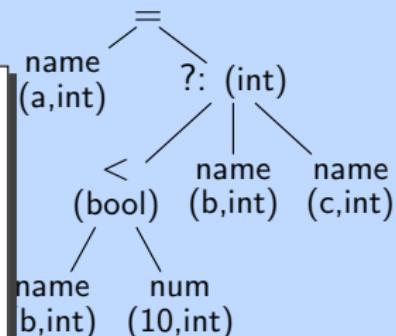
Issues:

- Offsets of variables in the stack frame
- Actual register numbers and assembly mnemonics
- Code to construct and discard activation records

Instruction List

```

T0 ← b
T0 ← T0 < b
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if T0 ≥ 0 goto L0:
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goto L1:
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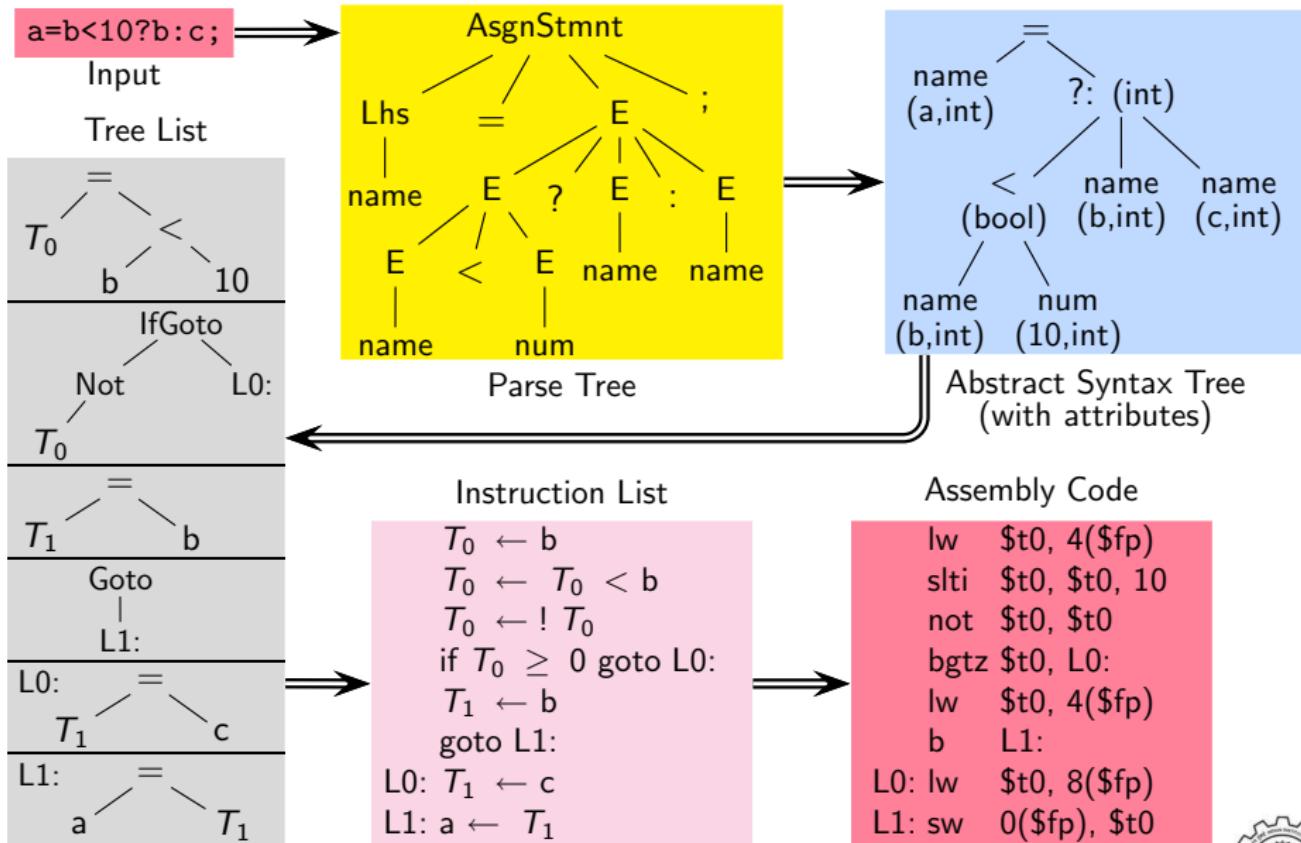
Abstract Syntax Tree  
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Assembly Code

```

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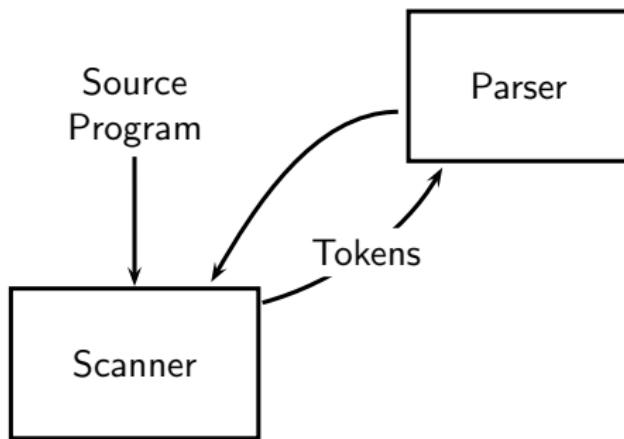
# Translation Sequence in Our Example: Emitting Instructions



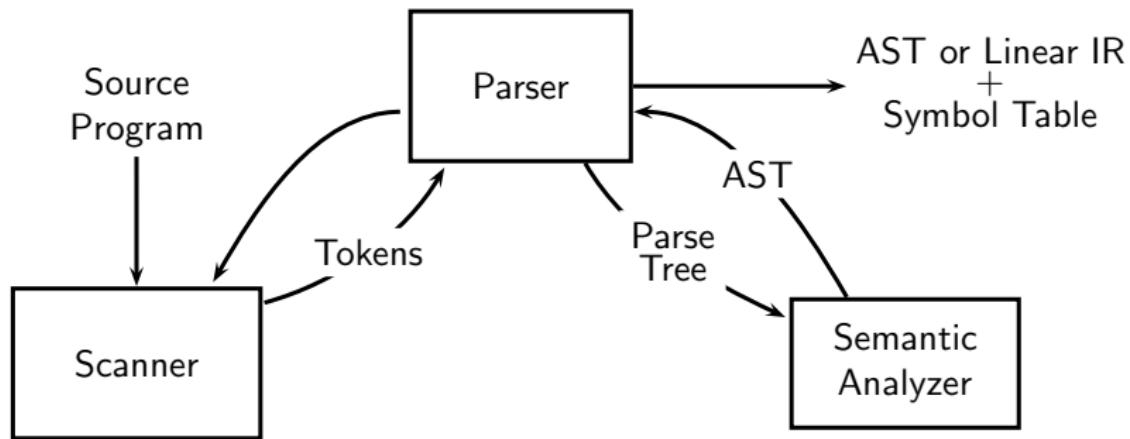
# Typical Front Ends

Parser

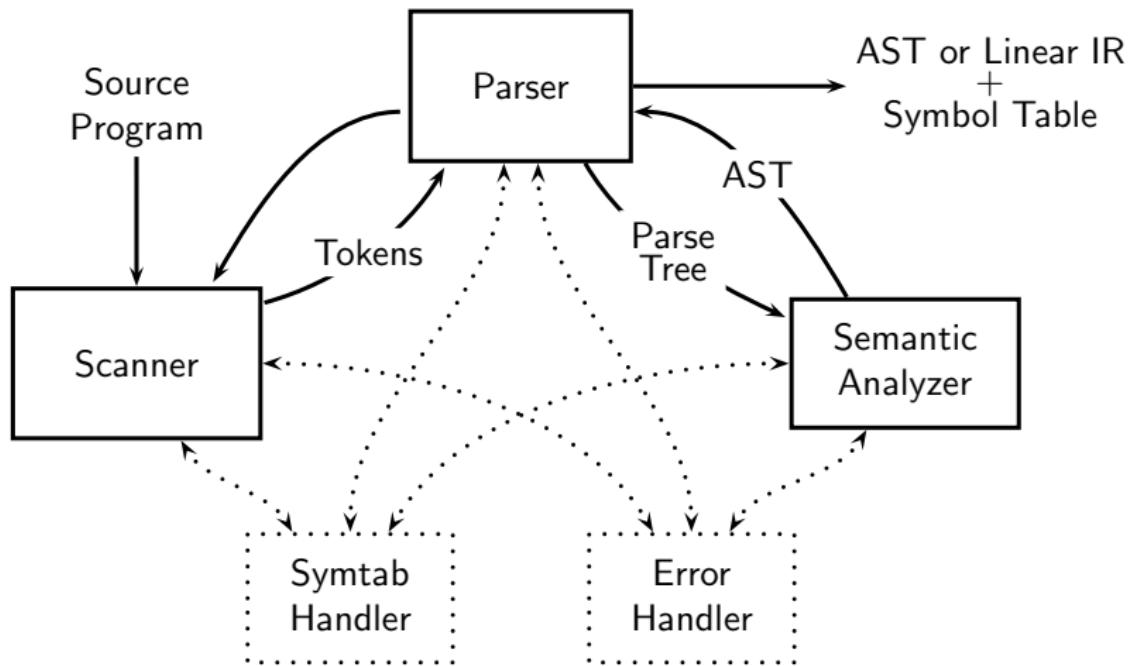
## Typical Front Ends



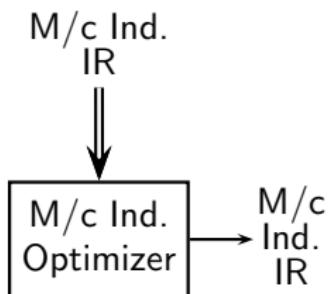
## Typical Front Ends



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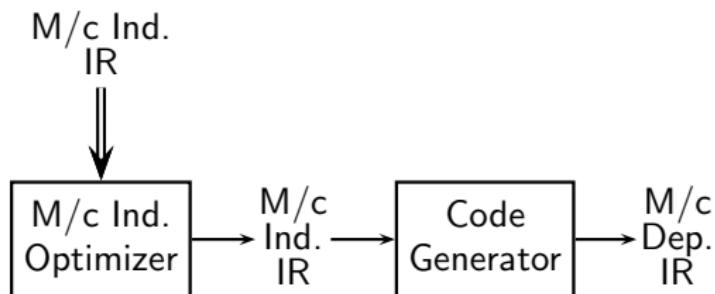


## Typical Back Ends



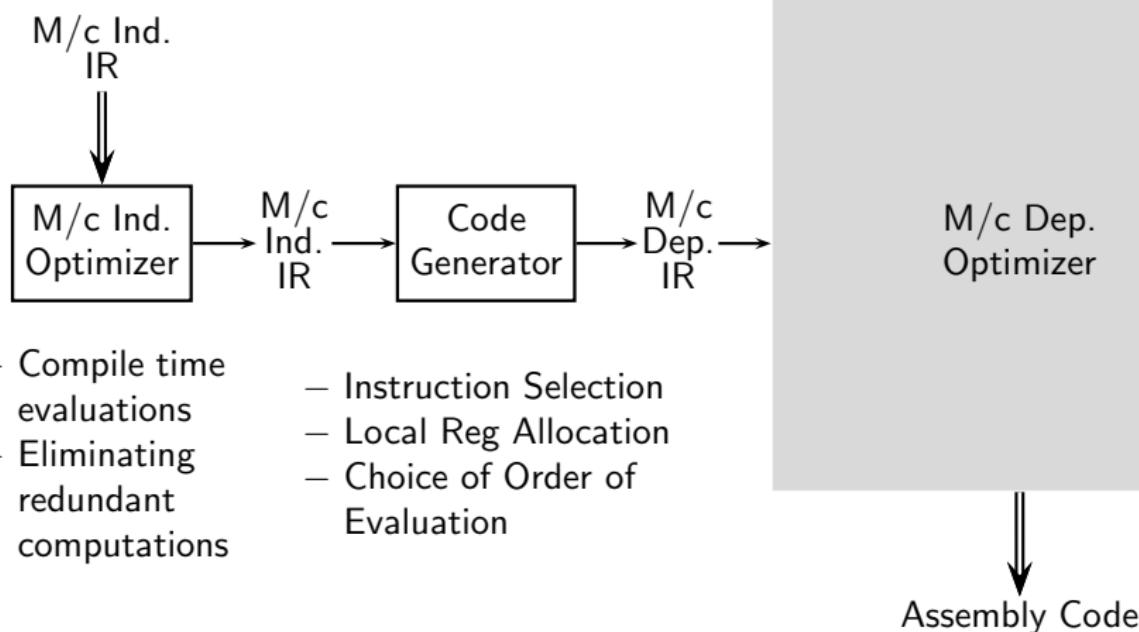
- Compile time evaluations
- Eliminating redundant computations

## Typical Back Ends

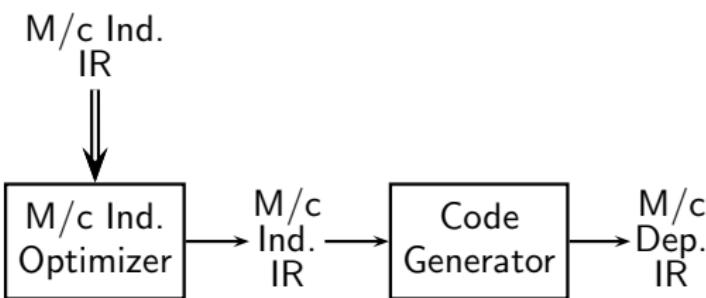


- Compile time evaluations
- Eliminating redundant computations
- Instruction Selection
- Local Reg Allocation
- Choice of Order of Evaluation

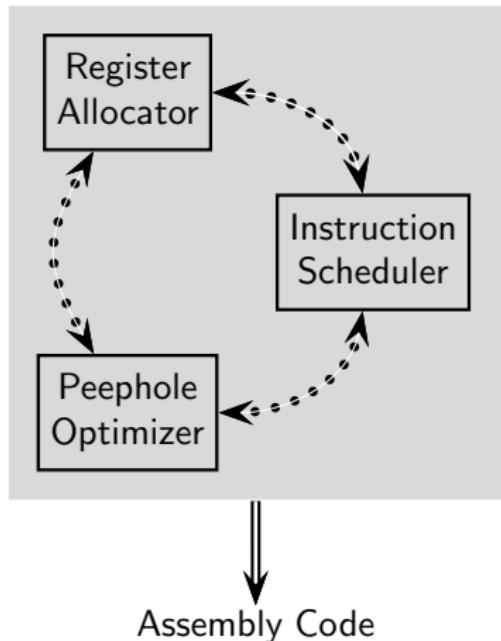
## Typical Back Ends



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- Compile time evaluations
- Eliminating redundant computations
- Instruction Selection
- Local Reg Allocation
- Choice of Order of Evaluation



*Part 4*

## *Introduction to GCC*

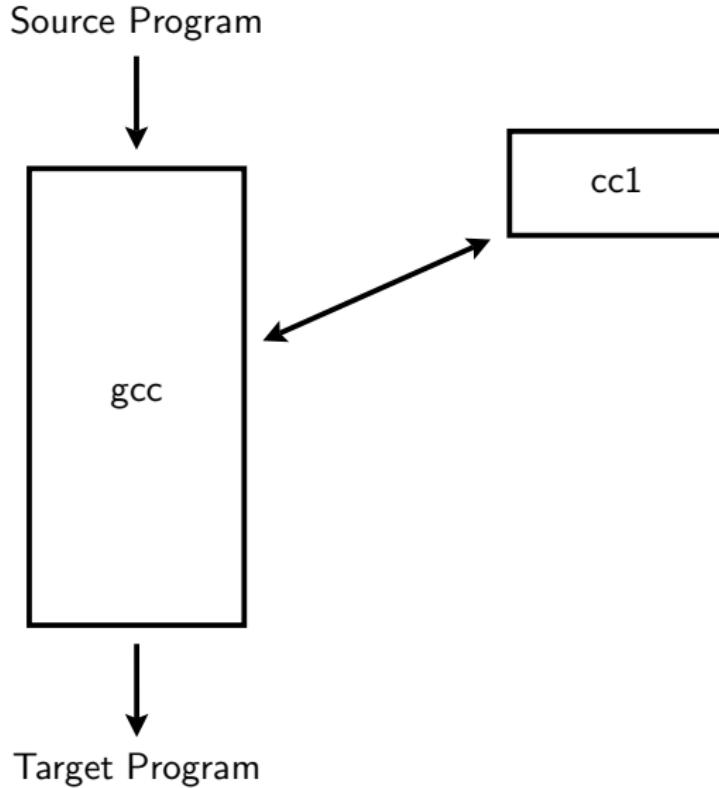
# The GNU Tool Chain

Source Program

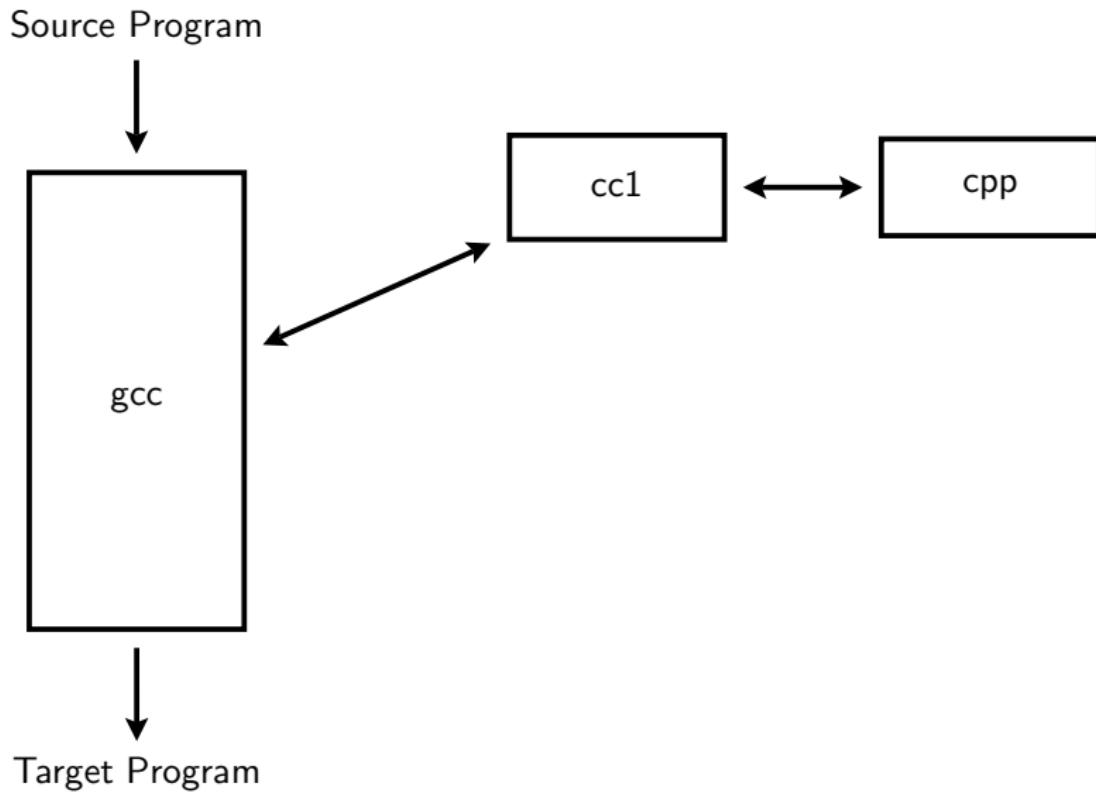


Target Program

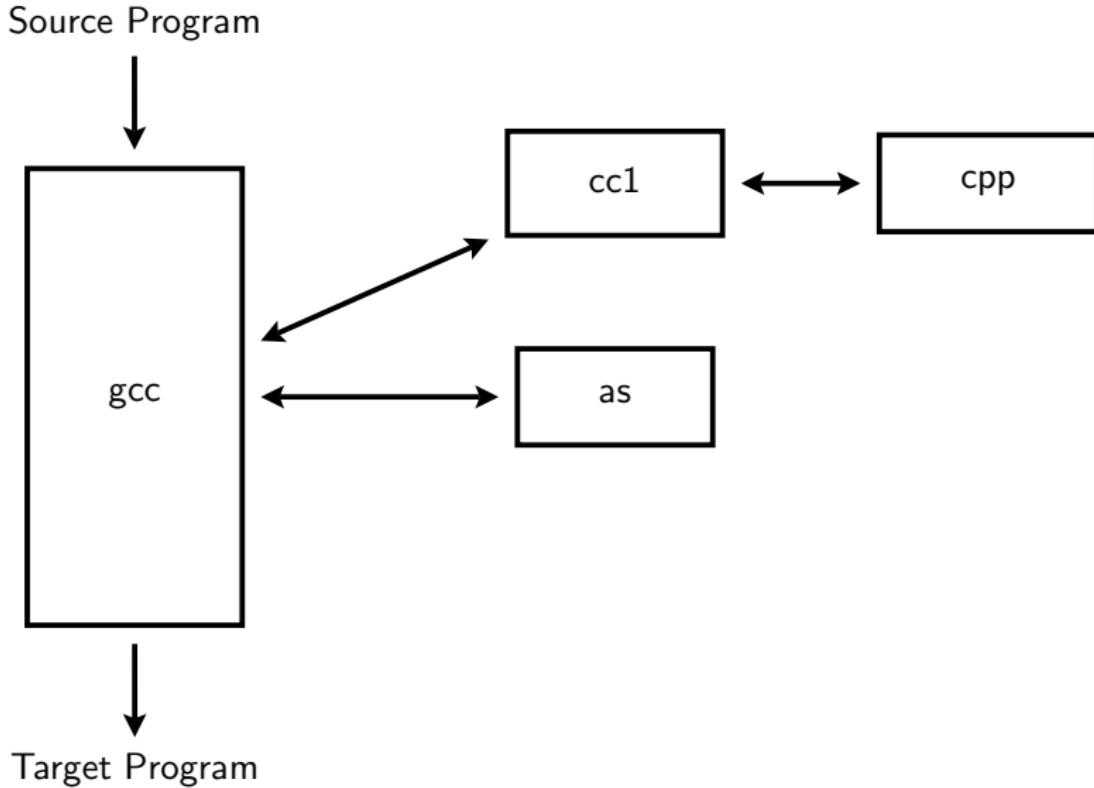
# The GNU Tool Chain



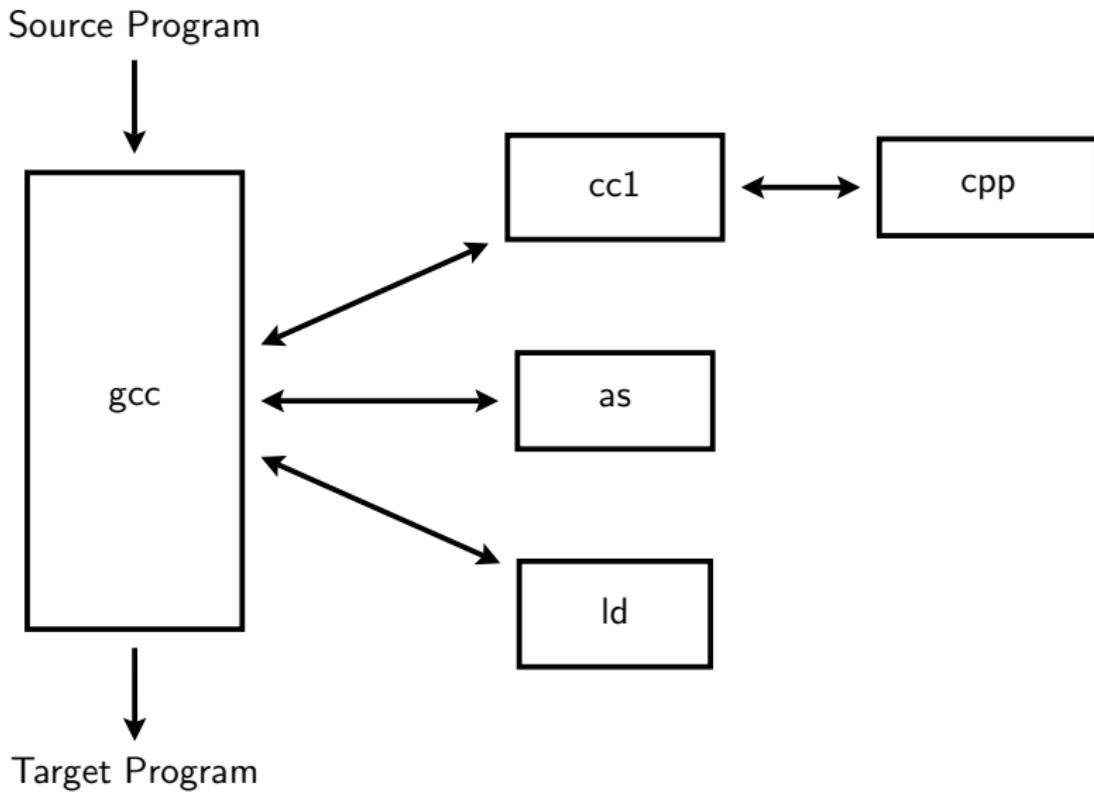
# The GNU Tool Chain



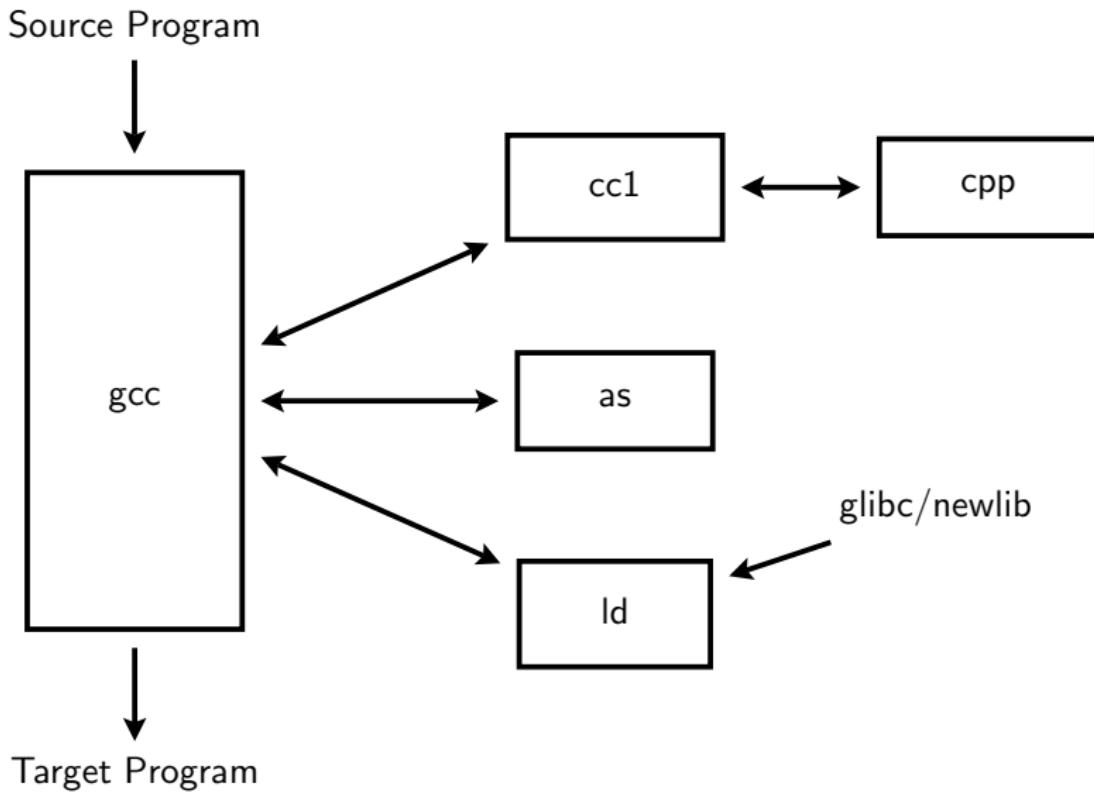
# The GNU Tool Chain



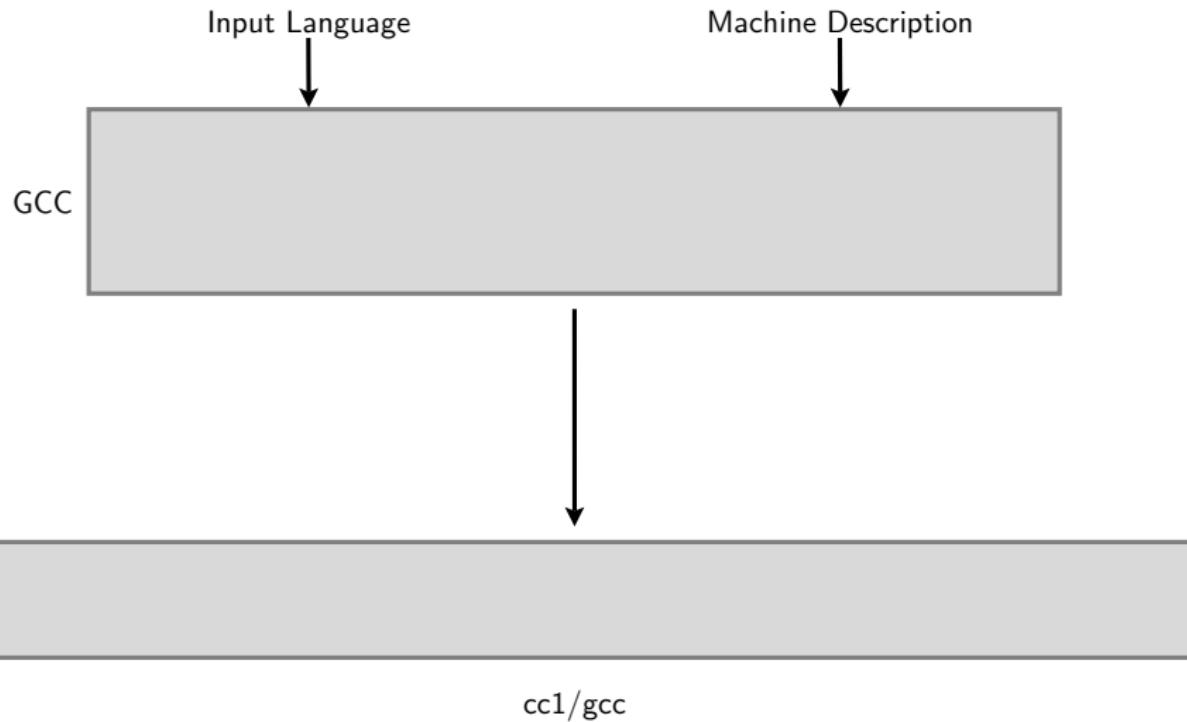
# The GNU Tool Chain



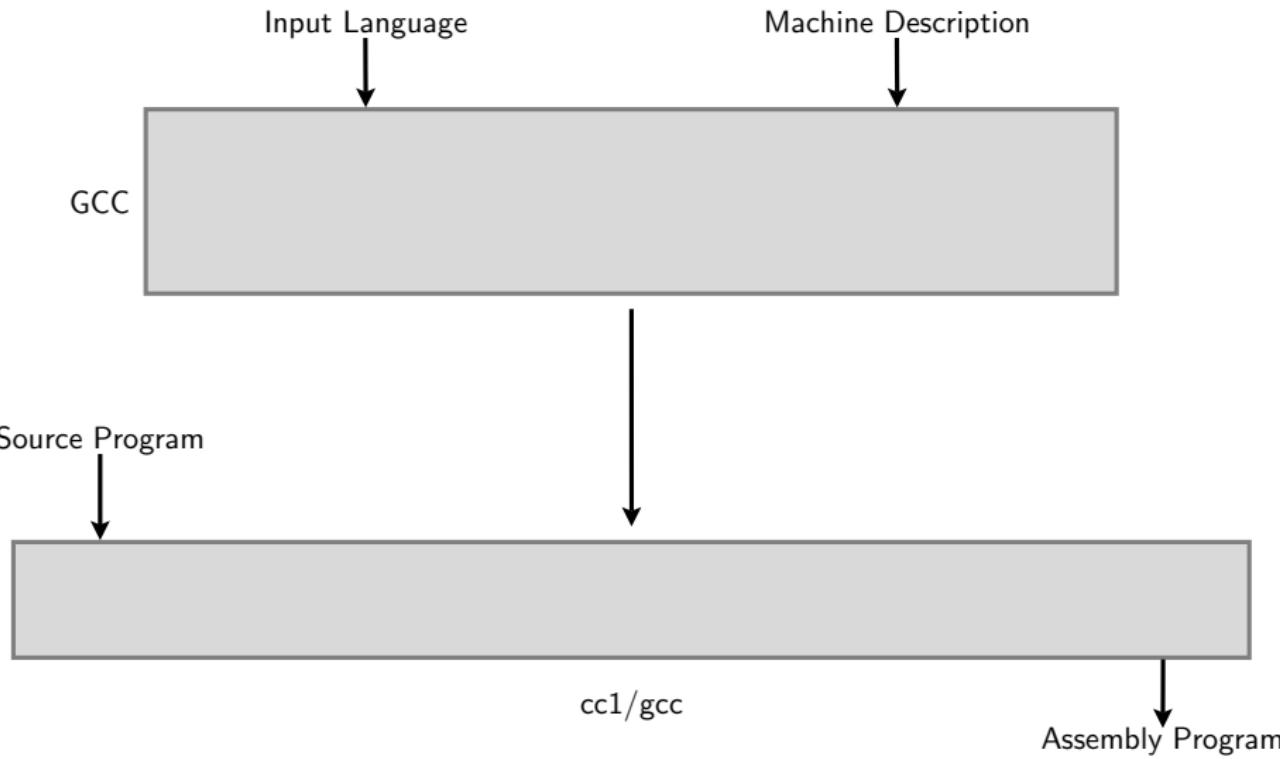
# The GNU Tool Chain



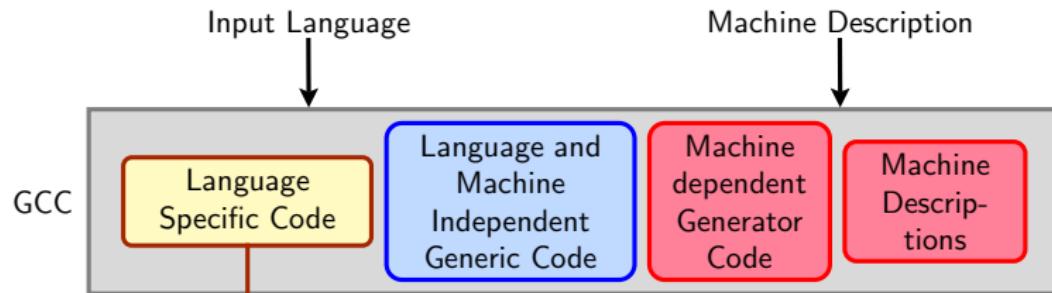
# The GCC Framework



# The GCC Framework

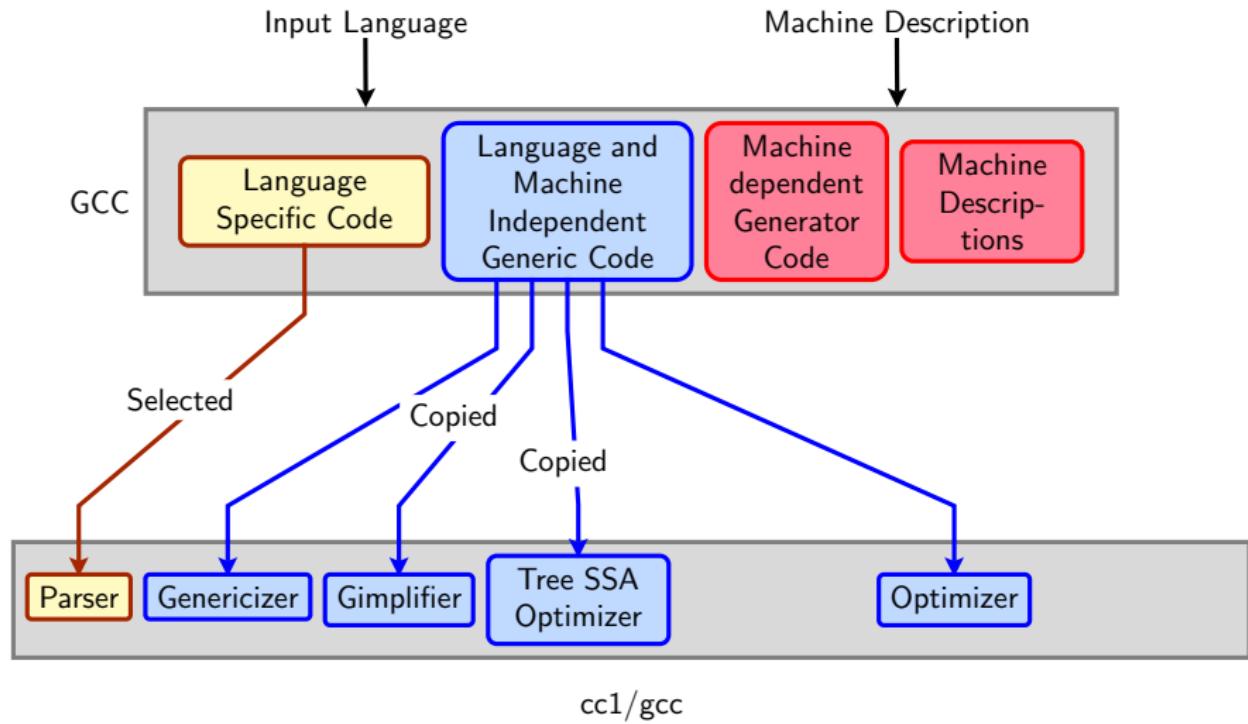


# The GCC Framework

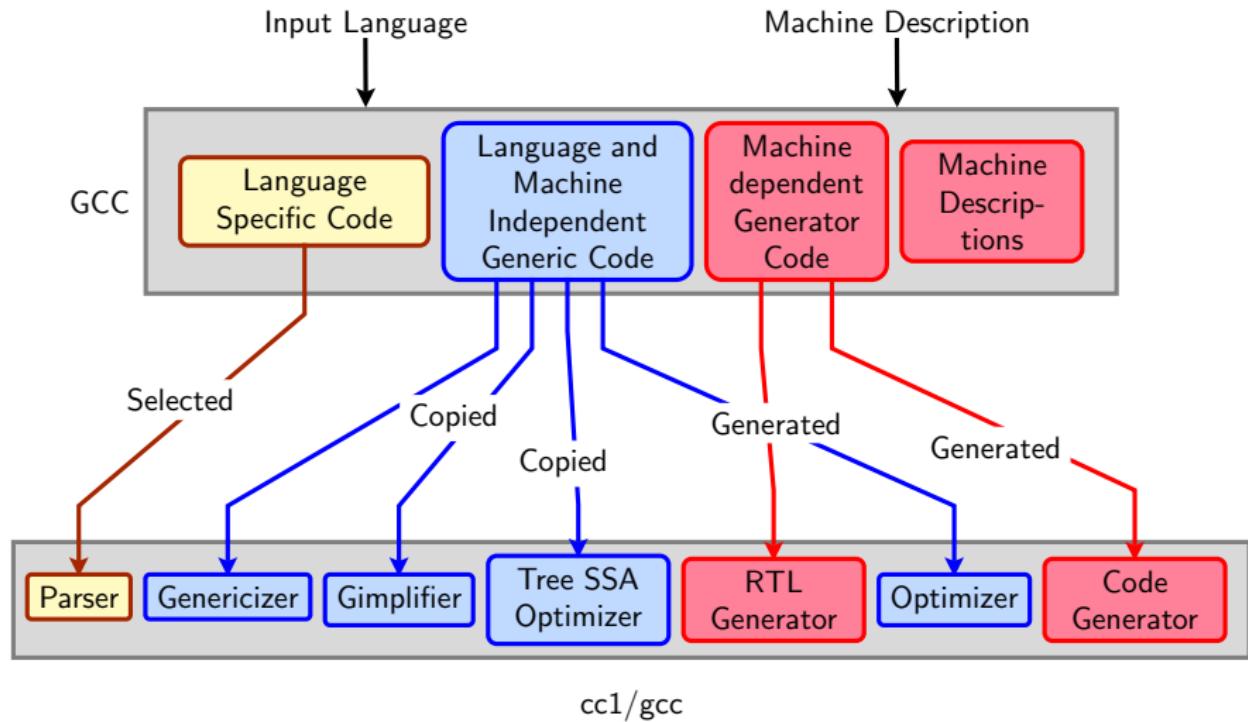


cc1/gcc

# The GCC Framework

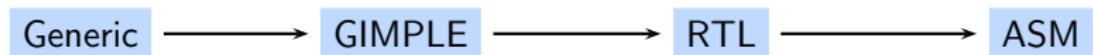


# The GCC Framework



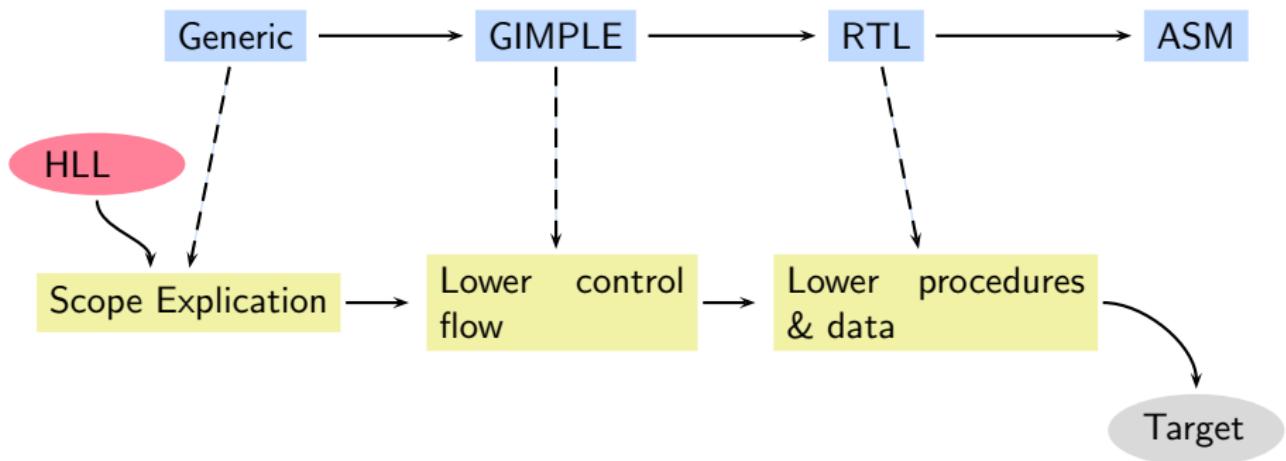
# The cc1 Phase Sequence as IR Chain

The GCC Phase Sequence



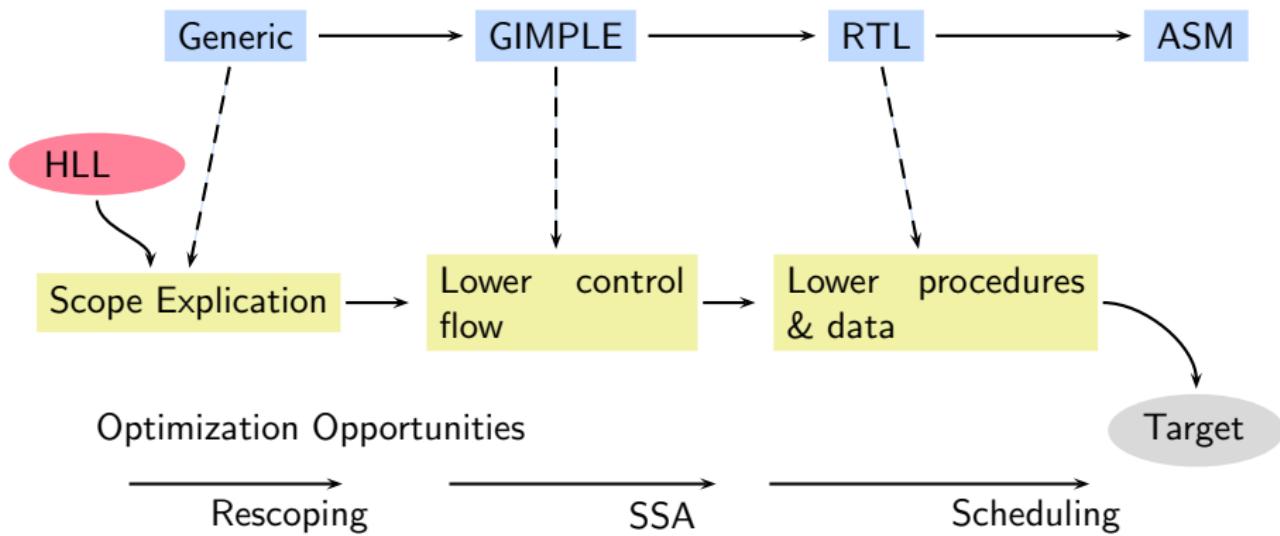
# The cc1 Phase Sequence as IR Chain

## The GCC Phase Sequence



# The cc1 Phase Sequence as IR Chain

## The GCC Phase Sequence



*Part 5*

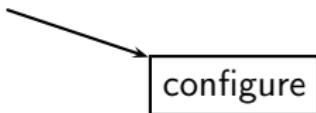
## *Configuration and Building*

# Configuring GCC

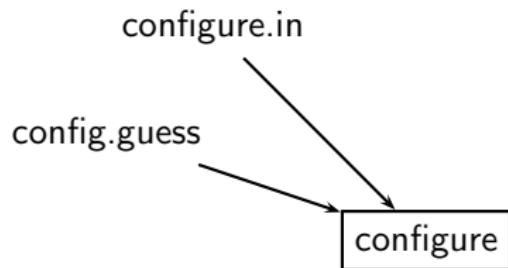
configure

# Configuring GCC

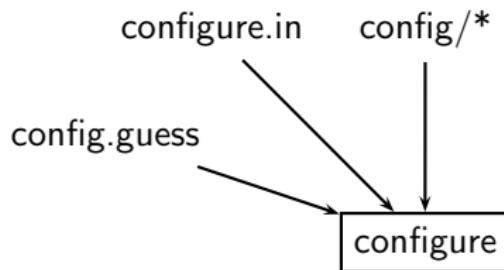
config.guess



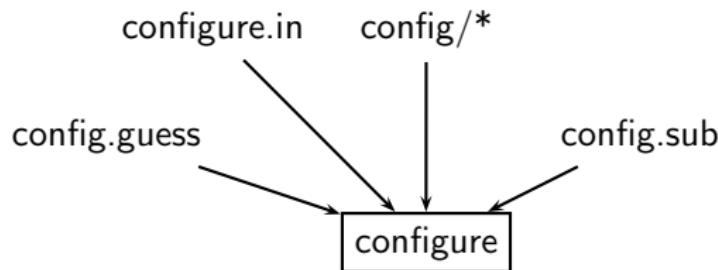
# Configuring GCC



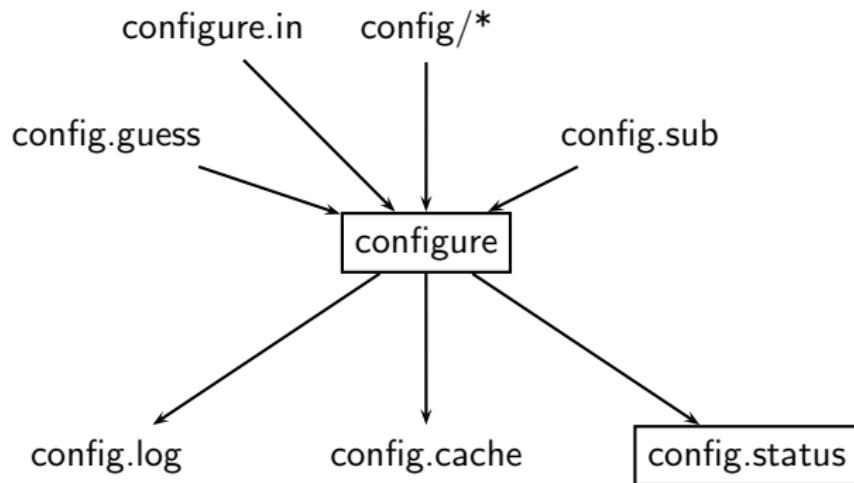
# Configuring GCC



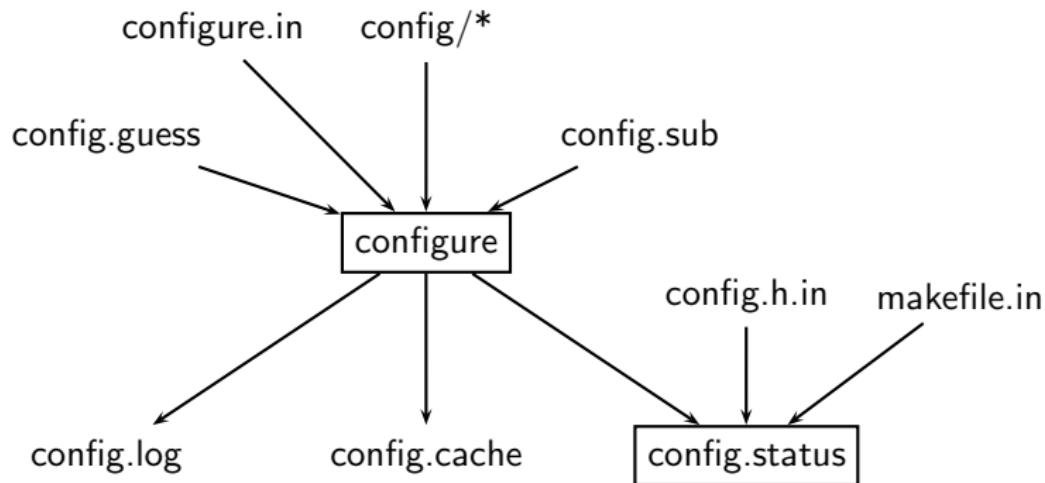
# Configuring GCC



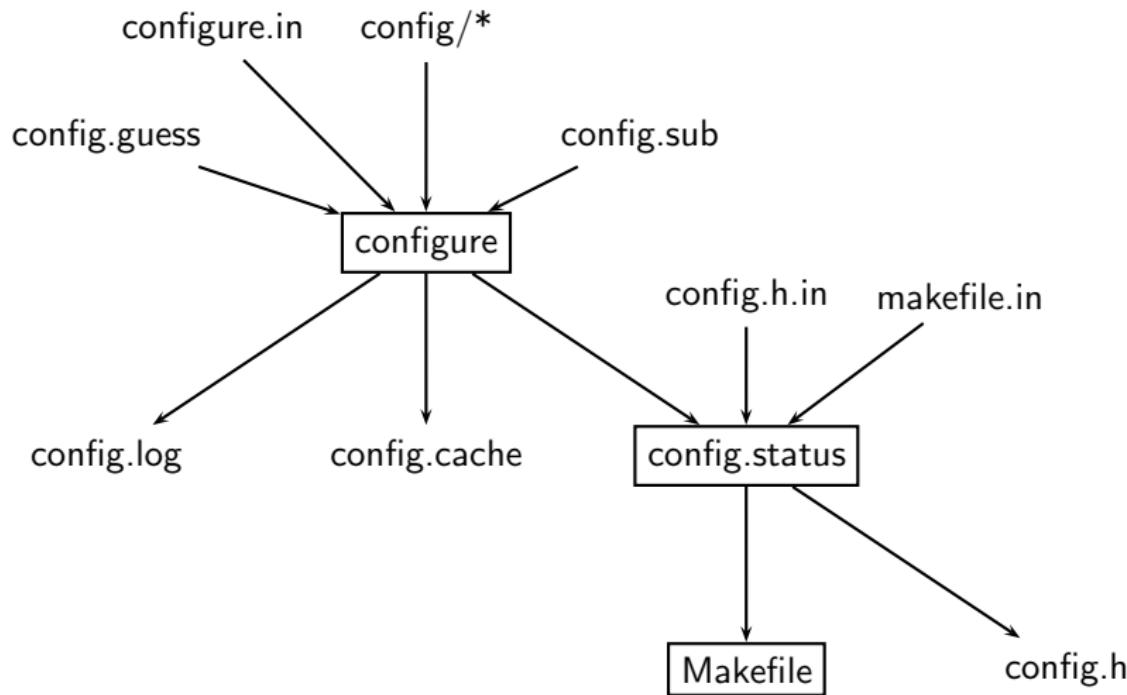
# Configuring GCC



# Configuring GCC



# Configuring GCC



*Part 6*

## *About The Course*