

**Outline***Workshop on Essential Abstractions in GCC***Gray Box Probing of GCC Translation Sequence**

GCC Resource Center  
([www.cse.iitb.ac.in/grc](http://www.cse.iitb.ac.in/grc))

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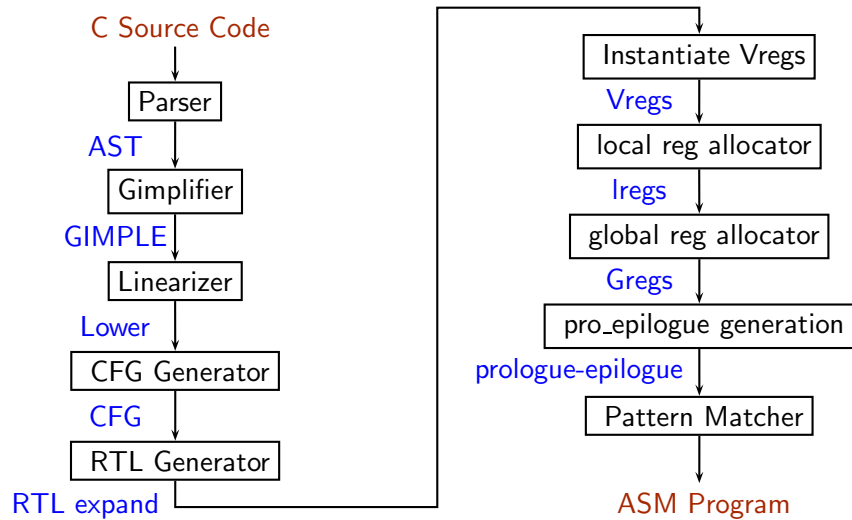
- Overview of translation sequence in GCC
- Overview of intermediate representations
- Intermediate representations of programs across important phases

**What is Gray Box Probing?**

- Black Box probing:  
Examining only the input and output relationship of a system
- White Box probing:  
Examining internals of a system for a given set of inputs
- Gray Box probing:  
Examining input and output of various components/modules of a system

**What is Gray Box Probing?****Notes**

## Important Phases of GCC



*Lowering of abstraction!*



## Important Phases of GCC

Notes



## Phases of GCC

To see output after each pass use the option  
-fdump-<ir>-<pass>

- <ir>
  - ▶ -tree-<pass>
    - ▶ gimple
    - ▶ original
    - ▶ cfg etc.
    - ▶ Use -all to see all dumps
  - ▶ -rtl-<pass>
    - ▶ expand
    - ▶ greg
    - ▶ vreg etc
    - ▶ Use -all to see all dumps

Example: `gcc -fdump-tree-all -fdump-rtl-all test.c`



## Phases of GCC

Notes



## GCC Internal Representation

### Example: test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

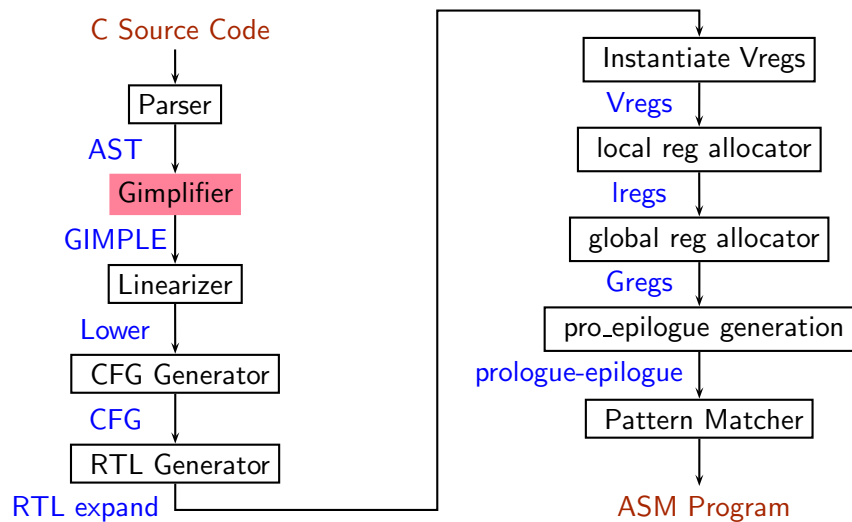


## GCC Internal Representation

# Notes



## Important Phases of GCC



## Important Phases of GCC

# Notes



## Gimplifier

- Three-address representation derived from GENERIC by breaking down into tuples of no more than 3 operands
- Temporaries are introduced to hold intermediate values needed to compute complex expressions.
- Control structures are explicated into conditional jumps.



## Gimple

```
Gimple
goto <D.1197>;
<D.1196>;
a = a + 1;
<D.1197>;
if (a <= 7)
{
    goto <D.1196>;
}
else
{
    goto <D.1198>;
}
<D.1198>;
```

Source

```
while (a <= 7)
{
    a = a+1;
}
```



## Gimplifier

# Notes



## Gimple

# Notes



## Gimple

```

Gimple
if (a <= 12)
{
    D.1199 = a + b;
    a = D.1199 + c;
}
else
{
}

```

Source

```

Source
if (a <= 12)
{
    a = a+b+c;
}

```

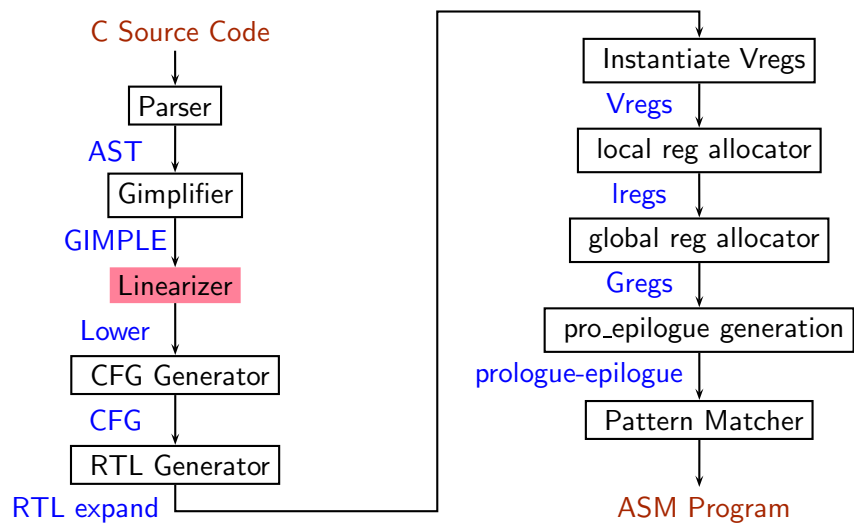


## Gimple

Notes



## Important Phases of GCC



## Important Phases of GCC

Notes



## Gimple - Lower

Lower	Gimple
<pre>if (a &lt;= 12) goto &lt;D.1200&gt;; else goto &lt;D.1201&gt;; &lt;D.1200&gt;; D.1199 = a + b; a = D.1199 + c; &lt;D.1201&gt;; return;</pre>	<pre>if (a &lt;= 12) {     D.1199 = a + b;     a = D.1199 + c; }</pre>

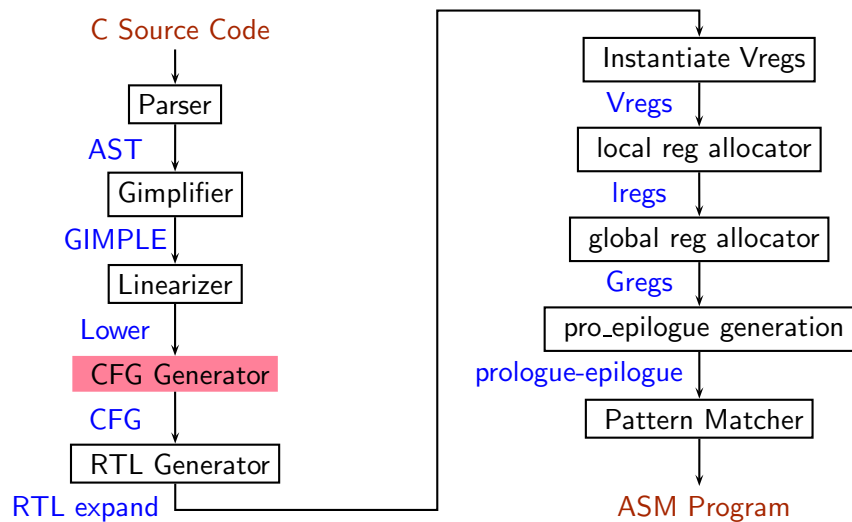


## Gimple - Lower

Notes



## Important Phases of GCC



## Important Phases of GCC

Notes



## Lower - CFG

CFG	Lower
<pre># BLOCK 5 if (a &lt;= 7)   goto &lt;bb 6&gt;; else   goto &lt;bb 7&gt;; # SUCC: 6 (true) 7 (false) # BLOCK 6 D.1199 = a + b; a = D.1199 + c; # SUCC: 7 (fallthru) # BLOCK 7 return; # SUCC: EXIT</pre>	<pre>if (a &lt;= 12) goto &lt;D.1200&gt;; else goto &lt;D.1201&gt;; &lt;D.1200&gt;; D.1199 = a + b; a = D.1199 + c; &lt;D.1201&gt;; return;</pre>

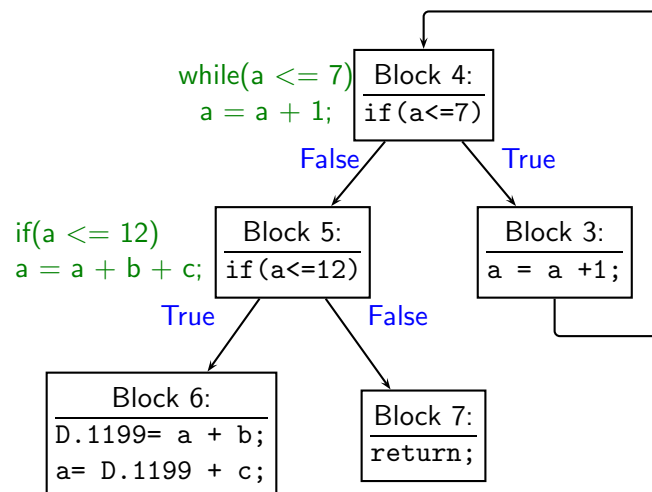


## Lower - CFG

## Notes



## Control Flow Graph



## Control Flow Graph

## Notes

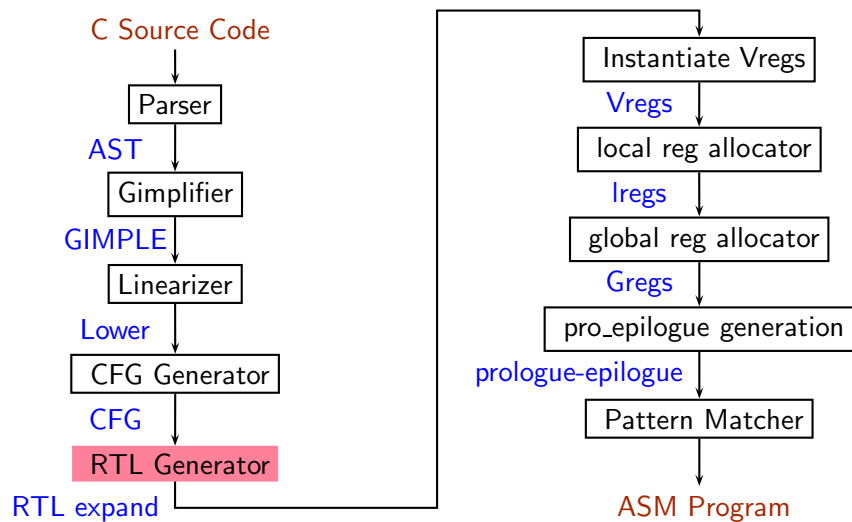


## Decisions that have been taken

- Three-address representation is generated
- All high level control flow structures are made explicit.
- Source code divided into interconnected blocks of sequential statements.
- This is a convenient structure for later analysis.



## Important Phases of GCC



## Decisions that have been taken

# Notes



## Important Phases of GCC

# Notes

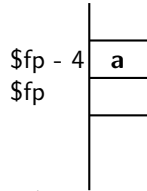




## CFG - RTL Expand

```
stack($fp - 4) = stack($fp - 4) + 1
|| flags=?
```

```
(insn 12 11 0 test.c:6 (parallel [
  (set (mem/c/i:SI (plus:SI
    (reg/f:SI 54 virtual-stack-vars)
    (const int -4 [...])) [...]))
    (plus:SI
      (mem/c/i:SI (plus:SI
        (reg/f:SI 54 virtual-stack-vars)
        (const int -4 [...])) [...]))
      (const int 1 [...]))))
(clobber (reg:CC 17 flags))
]) -1 (nil))
```



## CFG - RTL Expand

Notes



## RTL in Spim

```
a = a + 1;
```

```
Expr:
r39=stack($fp - 4)
r40=r39+1
stack($fp - 4)=r40
```

```
(insn 7 6 8 test.c:6 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])) -1 (nil))
(insn 8 7 9 test.c:6 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...])))) -1 (nil))
(insn 9 8 0 test.c:6 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...]))
  (reg:SI 40)) -1 (nil))
```



## RTL in Spim

Notes



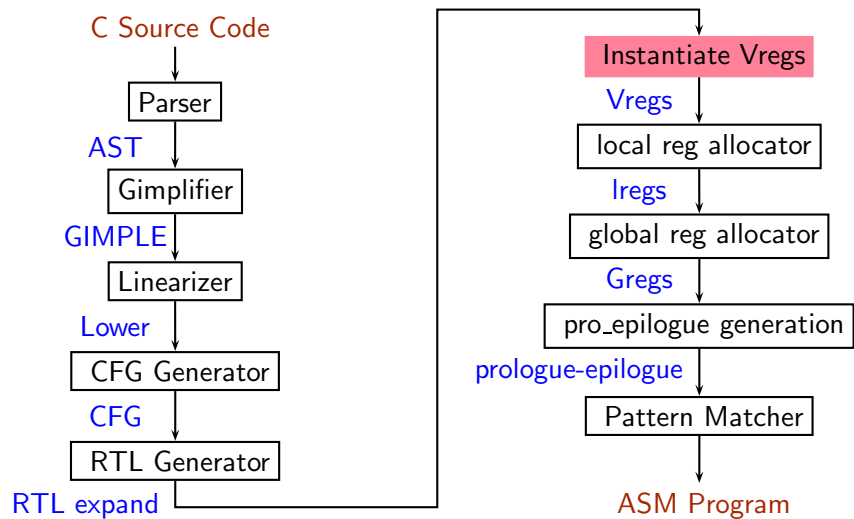
## RTL generation

Decisions that are already taken

- A virtual stack is generated.
- RTL instruction is dependant upon architecture.
- Position of each variable is fixed with reference to frame pointer.



## Important Phases of GCC



## RTL generation

Notes



## Important Phases of GCC

Notes



## RTL Expand - Vreg

```
(insn 7 6 8 3 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 1 $at )
    (const_int -4 [...])) [...])) 4 *IITB_move_from_mem (nil))
(insn 8 7 9 3 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...])))) 12 addsi3 (nil))
(insn 9 8 18 3 (set (mem/c/i:SI (plus:SI (reg/f:SI 1 $at )
  (const_int -4 [...])) []))
  (reg:SI 40)) 5 *IITB_move_to_mem (nil))
```



## Vregs generation

## Decisions that are already taken

- Virtual stack is assigned to a hard register.
- Each insn is committed to an instruction in machine description.
- This commitment helps in scheduling, peephole, and assembly code generation.



## RTL Expand - Vreg

## Notes

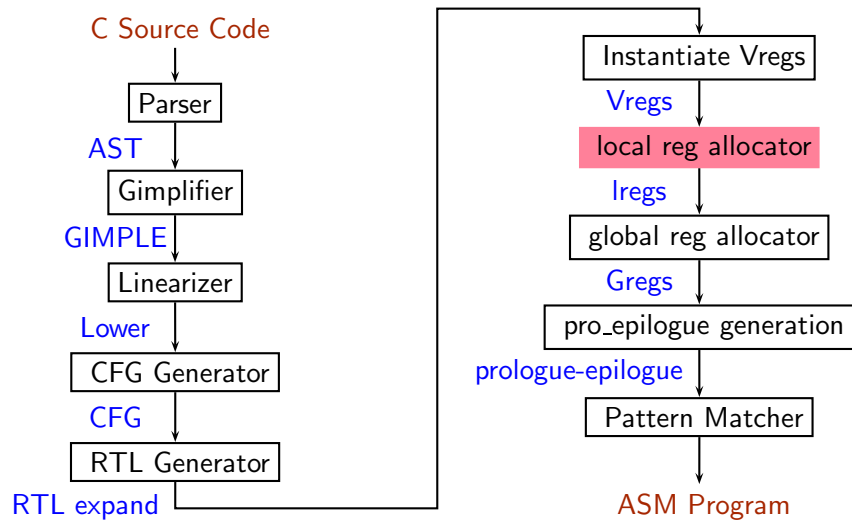


## Vregs generation

## Notes



## Important Phases of GCC



## Important Phases of GCC

# Notes



## Vregs - Lregs

```

(insn 7 6 8 2 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 1 $at)
    (const_int -4 [...])) [...]))
  4 *IITB_move_from_mem (nil))
(insn 8 7 9 2 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...])) 12 addsi3
  (expr_list:REG DEAD (reg:SI 39) (nil)))
(insn 9 8 12 2 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 1 $at)
    (const_int -4 [...])) [...])
  (reg:SI 40) 5 *IITB_move_to_mem
  (expr_list:REG DEAD(reg:SI 40) (nil)))
  
```

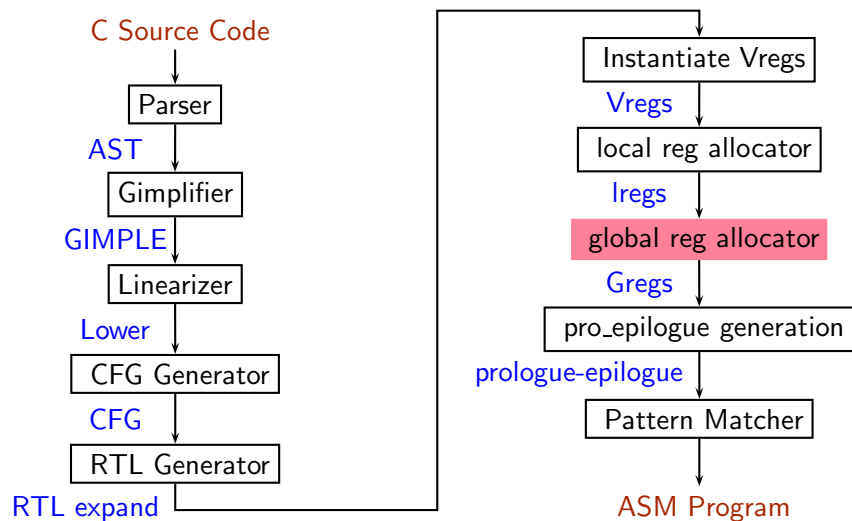


## Vregs - Lregs

# Notes



## Important Phases of GCC



## Important Phases of GCC

# Notes



## Lregs - Gregs

```

(insn 7 6 8 3 test.c:4 (set (reg:SI 2 $v0 [39])
  (mem/c/i:SI (plus:SI (reg/f:SI 1 $fp )
    (const_int -4 [...])) [...]))
  4 *IITB_move_from_mem (nil))
(insn 8 7 9 3 test.c:4 (set (reg:SI $v0 [40])
  (plus:SI (reg:SI $v0 [39])
    (const_int 1 [...]))) 12 addsi3 (nil))
(insn 9 8 18 3 test.c:4 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 1 $fp )
    (const_int -4 [...])) [...])
  (reg:SI $v0 40)) 5 *IITB_move_to_mem (nil))
  
```



## Lregs - Gregs

# Notes



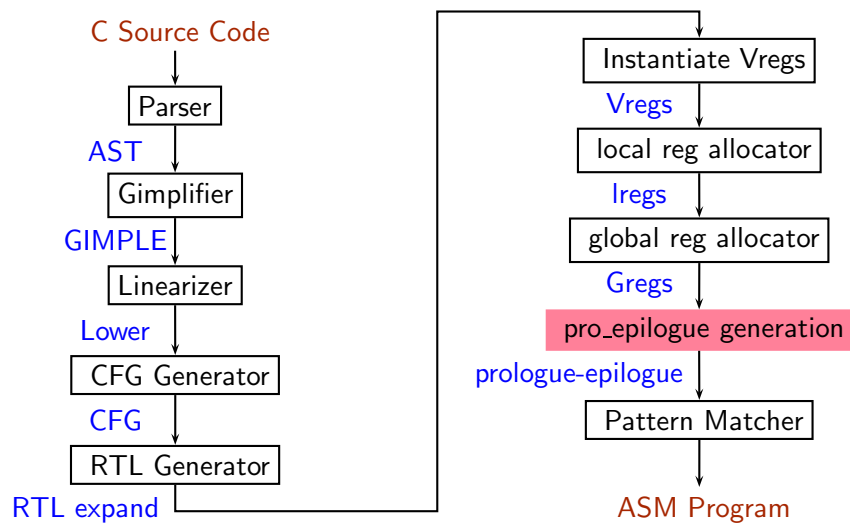
## After Greg generation

### Decisions that are already taken

- Local and global register are being allocated.
- Register usage has been annotated by REG\_DEAD etc.



## Important Phases of GCC



## After Greg generation

# Notes



## Important Phases of GCC

# Notes



## RTL for Function Calls in Spim

Calling function	Called function
<ul style="list-style-type: none"> <li>Allocate memory for activation records (push)</li> <li>Copy arguments into stack arguments</li> <li><b>Call function</b></li> <li>Get result from stack (pop)</li> <li>Deallocate memory for activation record (pop)</li> </ul>	<ul style="list-style-type: none"> <li>Allocate memory for return value (push)</li> <li>Store mandatory callee save registers (push)</li> <li>Set frame pointer</li> <li>Allocate local variables (push)</li> <li><b>Execute code</b></li> <li>Put result in return value space</li> <li>Deallocate local variables (pop)</li> <li>Load callee save registers (pop)</li> <li>Return</li> </ul>



## RTL for Function Calls in Spim

# Notes



## Prologue and Epilogue

```

(insn 17 3 18 2 test.c:2
  (set (mem:SI (reg/f:SI 29 $sp) [0 S4 A8])
    (reg:SI 31 $ra)) -1 (nil))
(insn 18 17 19 2 test.c:2
  (set (mem:SI (plus:SI (reg/f:SI 29 $sp)
    (const_int -4 [...])) [...])
    (reg/f:SI 29 $sp)) -1 (nil))
(insn 19 18 20 2 test.c:2 (set
  (mem:SI (plus:SI (reg/f:SI 29 $sp)
    (const_int -8 [...])) [...])
  (reg/f:SI 30 $fp)) -1 (nil))
(insn 20 19 21 2 test.c:2 (set
  (reg/f:SI 30 $fp)
  (reg/f:SI 29 $sp)) -1 (nil))
(insn 21 20 22 2 test.c:2 (set
  (reg/f:SI 29 $sp)
  (plus:SI (reg/f:SI 30 $fp)
    (const_int -32 [...])) -1 (nil))

```

```

sw $ra, 0($sp)
sw $sp, 4($sp)
sw $fp, 8($sp)
move $fp,$sp
addi $sp,$fp,32

```

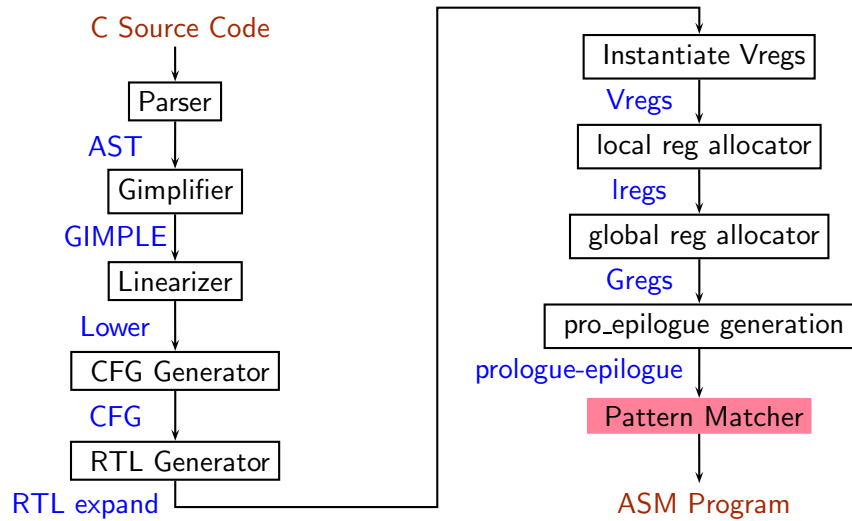


## Prologue and Epilogue

# Notes



## Important Phases of GCC



## Important Phases of GCC

# Notes



## Assembly

Assembly Code for  $a = a + 1$ ;

For spim

```
lw $v0, -8($fp)
addi $v0, $v0, 1
sw $v0, -8($fp)
```

For i386

```
addl $1, -8(%ebp)
```



## Assembly

# Notes





## Conclusion

- Source code is transformed into assembly by moving it step by step close to machine architecture.
- This transformation can be understood to a large extent by observing their inputs and output.
- Fortunately in gcc, output of all the passes can be seen using fdump.
- Complete list of dumps can be seen by doing man gcc.



## Conclusion

# Notes

