

*Workshop on Essential Abstractions in GCC***Outline****Incremental Machine Descriptions for Spim:  
Levels 2, 3, and 4**

GCC Resource Center  
([www.cse.iitb.ac.in/grc](http://www.cse.iitb.ac.in/grc))

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- Constructs supported in level 2
- Constructs supported in level 3
- Constructs supported in level 4

*Part 1**Constructs Supported in Level 2*

## Arithmetic Operations Required in Level 2

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow Src_1 - Src_2$	$R_i \leftarrow R_j - R_k$	sub ri, rj, rk	level 2
$Dest \leftarrow -Src$	$R_i \leftarrow -R_j$	neg ri, rj	
$Dest \leftarrow Src_1 / Src_2$	$R_i \leftarrow R_j / R_k$	div rj, rk mflo ri	
$Dest \leftarrow Src_1 \% Src_2$	$R_i \leftarrow R_j \% R_k$	rem ri, rj, rk	
$Dest \leftarrow Src_1 * Src_2$	$R_i \leftarrow R_j * R_k$	mul ri, rj, rk	



## Bitwise Operations Required in Level 2

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow Src_1 \ll Src_2$	$R_i \leftarrow R_j \ll R_k$	sllv ri, rj, rk	level 2
	$R_i \leftarrow R_j \ll C_5$	sll ri, rj, c	
$Dest \leftarrow Src_1 \gg Src_2$	$R_i \leftarrow R_j \gg R_k$	srav ri, rj, rk	
	$R_i \leftarrow R_j \gg C_5$	sra ri, rj, c	
$Dest \leftarrow Src_1 \& Src_2$	$R_i \leftarrow R_j \& R_k$	and ri, rj, rk	
	$R_i \leftarrow R_j \& C$	andi ri, rj, c	
$Dest \leftarrow Src_1   Src_2$	$R_i \leftarrow R_j   R_k$	or ri, rj, rk	
	$R_i \leftarrow R_j   C$	ori ri, rj, c	
$Dest \leftarrow Src_1 \wedge Src_2$	$R_i \leftarrow R_j \wedge R_k$	xor ri, rj, rk	
	$R_i \leftarrow R_j \wedge C$	xori ri, rj, c	
$Dest \leftarrow \sim Src$	$R_i \leftarrow \sim R_j$	not ri, rj	



## Arithmetic Operations Required in Level 2

Notes



## Bitwise Operations Required in Level 2

Notes



## Divide Operation in spim2.md using define\_insn

- For division, the spim architecture imposes use of multiple asm instructions for single operation.
- Two ASM instructions are emitted using single RTL pattern

```
(define_insn "divsi3"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (div:SI (match_operand:SI 1 "register_operand" "r")
                 (match_operand:SI 2 "register_operand" "r"))
        )]
  ""
  "div\t%1, %2\n\tmflo\t%0"
)
```



## Advantages/Disadvantages of using define\_insn

- Very simple to add the pattern
- Primitive target feature represented as single insn pattern in .md
- Unnecessary atomic grouping of instructions
- May hamper optimizations in general, and instruction scheduling, in particular



## Divide Operation in spim2.md using define\_insn

# Notes



## Advantages/Disadvantages of using define\_insn

# Notes



## Divide Operation in spim2.md using define\_expand

- The RTL pattern can be expanded into two different RTLs.

```
(define_expand "divsi3"
  [(parallel[(set (match_operand:SI 0 "register_operand" "")
    (div:SI (match_operand:SI 1 "register_operand" "")
      (match_operand:SI 2 "register_operand" ""))
    )
  (clobber (reg:SI 26))
  (clobber (reg:SI 27))]]]
  ""
  {
    emit_insn(gen_IITB_divide(gen_rtx_REG(SImode,26),
                              operands[1], operands[2]));
    emit_insn(gen_IITB_move_from_lo(operands[0],
                                     gen_rtx_REG(SImode,26)));
    DONE;
  }
)
```



## Divide Operation in spim2.md using define\_expand

- Divide pattern equivalent to div instruction in architecture.

```
(define_insn "IITB_divide"
  [(parallel[(set (match_operand:SI 0 "LO_register_operand" "=q")
    (div:SI (match_operand:SI 1 "register_operand" "r")
      (match_operand:SI 2 "register_operand" "r"))
    )
  (clobber (reg:SI 27))]]]
  ""
  "div t%1, %2"
)
```



## Divide Operation in spim2.md using define\_expand

Notes



## Divide Operation in spim2.md using define\_expand

Notes



## Divide Operation in spim2.md using define\_expand

- Moving contents of special purpose register LO to/from general purpose register

```
(define_insn "IITB_move_from_lo"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (match_operand:SI 1 "LO_register_operand" "q"))]
  ""
  "mflo \t\t%0"
)

(define_insn "IITB_move_to_lo"
  [(set (match_operand:SI 0 "LO_register_operand" "=q")
        (match_operand:SI 1 "register_operand" "r"))]
  ""
  "mtlo \t\t%1"
)
```



## Divide Operation in spim2.md using define\_expand

- Divide pattern equivalent to div instruction in architecture.

```
(define_insn "modsi3"
  [(parallel [(set (match_operand:SI 0 "register_operand" "=r")
                  (mod:SI (match_operand:SI 1 "register_operand" "r")
                        (match_operand:SI 2 "register_operand" "r")))
             (clobber (reg:SI 26))
             (clobber (reg:SI 27))])]
  ""
  "rem \t\t%0, %1, %2"
)
```



## Divide Operation in spim2.md using define\_expand

Notes



## Divide Operation in spim2.md using define\_expand

Notes



## Advantages/Disadvantages of Using `define_expand` for Division

- Two instructions are separated out at GIMPLE to RTL conversion phase
- Both instructions can undergo all RTL optimizations independently
- C interface is needed in md
- Compilation becomes slower and requires more space



## Divide Operation in `spim2.md` using `define_split`

```
(define_split
  [(parallel [(set (match_operand:SI 0 "register_operand" "")
    (div:SI (match_operand:SI 1 "register_operand" "")
      (match_operand:SI 2 "register_operand" ""))
    )
    (clobber (reg:SI 26))
    (clobber (reg:SI 27))]]]
  ""
  [(parallel [(set (match_dup 3)
    (div:SI (match_dup 1)
      (match_dup 2)))
    (clobber (reg:SI 27))]]
    (set (match_dup 0)
      (match_dup 3))
  ])
  "operands[3]=gen_rtx_REG(SImode,26); "
```



## Advantages/Disadvantages of Using `define_expand` for Division

Notes



## Divide Operation in `spim2.md` using `define_split`

Notes



Part 2

## Constructs Supported in Level 3

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Spim MD Levels 2,3,4: Constructs Supported in Level 3

12/24

### Operations Required in Level 3

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow fun(P_1, \dots, P_n)$	call $L_{fun, n}$	lw $r_i$ , [SP+c1] sw $r_i$ , [SP] :	Level 1
		lw $r_i$ , [SP+c2] sw $r_i$ , [SP-n*4]	
		jal L	New
		$Dest \leftarrow \$v0$	level 1
$fun(P_1, P_2, \dots, P_n)$	call $L_{fun, n}$	lw $r_i$ , [SP+c1] sw $r_i$ , [SP] :	Level 1
		lw $r_i$ , [SP+c2] sw $r_i$ , [SP-n*4]	
		jal L	New



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Spim MD Levels 2,3,4: Constructs Supported in Level 3

12/24

### Operations Required in Level 3

Notes



**Call Operation in spim3.md**

```
(define_insn "call"
  [(call (match_operand:SI 0 "memory_operand" "m")
        (match_operand:SI 1 "immediate_operand" "i"))
   (clobber (reg:SI 31))]
  ""
  "*"
  return emit_asm_call(operands,0);
  ""
)
```

**Call Operation in spim3.md**

```
(define_insn "call_value"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (call (match_operand:SI 1 "memory_operand" "m")
              (match_operand:SI 2 "immediate_operand" "i"))))
   (clobber (reg:SI 31))]
  ""
  "*"
  return emit_asm_call(operands,1);
  ""
)
```

**Call Operation in spim3.md**

# Notes

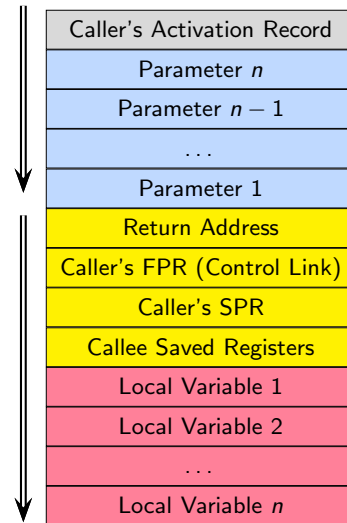
**Call Operation in spim3.md**

# Notes



## Activation Record Generation during Call

- Operations performed by caller
  - ▶ Push parameters on stack.
  - ▶ Load return address in return address register.
  - ▶ Transfer control to Callee.
- Operations performed by callee
  - ▶ Push Return address stored by caller on stack.
  - ▶ Push caller's Frame Pointer Register.
  - ▶ Push caller's Stack Pointer.
  - ▶ Save callee saved registers, if used by callee.
  - ▶ Create local variables frame.
  - ▶ Start callee body execution.



## Activation Record Generation during Call

Notes

## Prologue in spim3.md

```
(define_expand "prologue"
  [(clobber (const_int 0))]
  "")
{
  spim_prologue();
  DONE;
})
```

```
(set (mem:SI (reg:SI $sp))
     (reg:SI 31 $ra))

(set (mem:SI (plus:SI (reg:SI $sp)
                     (const_int -4 )))
     (reg:SI $sp))

(set (mem:SI (plus:SI (reg:SI $sp)
                     (const_int -8 )))
     (reg:SI $fp))

(set (reg:SI $fp)
     (reg:SI $sp))

(set (reg:SI $sp)
     (plus:SI (reg:SI $fp)
              (const_int -36)))
```



## Prologue in spim3.md

Notes



## Epilogue in spim3.md

```
(define_expand "epilogue"
  [(clobber (const_int 0))]
  ""
  (spim_epilogue();
   DONE;
  )

  (set (reg:SI $sp)
       (reg:SI $fp))

  (set (reg:SI $fp)
       (mem:SI (plus:SI (reg:SI $sp)
                        (const_int -8 ))))

  (set (reg:SI $ra)
       (mem:SI (reg:SI $sp)))

  (parallel [
    (return)
    (use (reg:SI $ra))])])
```



## Epilogue in spim3.md

# Notes



Part 3

*Constructs Supported in Level 4*

## Operations Required in Level 4

Operation	Primitive Variants	Implementation	Remark
$Src_1 < Src_2 ?$ goto L : PC	$CC \leftarrow R_i < R_j$ $CC < 0 ?$ goto L : PC	blt $r_i, r_j, L$	
$Src_1 > Src_2 ?$ goto L : PC	$CC \leftarrow R_i > R_j$ $CC > 0 ?$ goto L : PC	bgt $r_i, r_j, L$	
$Src_1 \leq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \leq R_j$ $CC \leq 0 ?$ goto L : PC	ble $r_i, r_j, L$	
$Src_1 \geq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \geq R_j$ $CC \geq 0 ?$ goto L : PC	bge $r_i, r_j, L$	



## Operations Required in Level 4

Operation	Primitive Variants	Implementation	Remark
$Src_1 == Src_2 ?$ goto L : PC	$CC \leftarrow R_i == R_j$ $CC == 0 ?$ goto L : PC	beq $r_i, r_j, L$	
$Src_1 \neq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \neq R_j$ $CC \neq 0 ?$ goto L : PC	bne $r_i, r_j, L$	



## Operations Required in Level 4

Notes



## Operations Required in Level 4

Notes



## Conditional Branch Instruction in spim4.md

```

(define_insn "cbranchsi4"
  [(set (pc)
        (if_then_else
          (match_operator:SI 0 "comparison_operator"
            [(match_operand:SI 1 "register_operand" "")
             (match_operand:SI 2 "register_operand" "")])
          (label_ref (match_operand 3 "" ""))
          (pc)))]
  ""
  "*"
  return conditional_insn(GET_CODE(operands[0]), operands);
  ""
)

```



## Support for Branch pattern in spim4.c

```

char *
conditional_insn (enum rtx_code code, rtx operands[])
{
  switch (code)
  {
    case EQ: return "beq %1, %2, %l3";
    case NE: return "bne %1, %2, %l3";
    case GE: return "bge %1, %2, %l3";
    case GT: return "bgt %1, %2, %l3";
    case LT: return "blt %1, %2, %l3";
    case LE: return "ble %1, %2, %l3";
    case GEU: return "bgeu %1, %2, %l3";
    case GTU: return "bgtn %1, %2, %l3";
    case LTU: return "bltn %1, %2, %l3";
    case LEU: return "bleu %1, %2, %l3";
    default: /* Error. Issue ICE */
  }
}

```



## Conditional Branch Instruction in spim4.md

Notes



## Support for Branch pattern in spim4.c

Notes



## Alternative for Branch: Conditional compare in spim4.md

```
(define_code_iterator cond_code
  [lt ltu eq ge geu gt gtu le leu ne])

(define_expand "cmpsi"
  [(set (cc0) (compare
    (match_operand:SI 0 "register_operand" "")
    (match_operand:SI 1 "nonmemory_operand" "")))]
  ""
  {
    compare_op0=operands[0];
    compare_op1=operands[1];
    DONE;
  }
)
```



## Alternative for Branch: Branch pattern in spim4.md

```
(define_expand "b<code>"
  [(set (pc) (if_then_else (cond_code:SI (match_dup 1)
    (match_dup 2))
    (label_ref (match_operand 0 "" ""))
    (pc)))]
  ""
  {
    operands[1]=compare_op0;
    operands[2]=compare_op1;
    if (immediate_operand(operands[2],SImode))
    {
      operands[2]=force_reg(SImode,operands[2]);
    }
  }
)
```



## Alternative for Branch: Conditional compare in spim4.md

Notes



## Alternative for Branch: Branch pattern in spim4.md

Notes



**Alternative for Branch: Branch pattern in spim4.md**

```
(define_insn "*insn_b<code>"
  [(set (pc)
        (if_then_else
         (cond_code:SI
          (match_operand:SI 1 "register_operand" "r")
          (match_operand:SI 2 "register_operand" "r"))
         (label_ref (match_operand 0 "" ""))
         (pc)))]
  ""
  "*"
  return conditional_insn(<CODE>,operands);
  ""
  )
```

**Alternative for Branch: Branch pattern in spim4.md**

Notes

