

## Spim Machine Descriptions

GCC Resource Center  
([www.cse.iitb.ac.in/grc](http://www.cse.iitb.ac.in/grc))

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Spim MD: Outline

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### Outline

- Systematic construction of machine descriptions
- Retargetting GCC to spim
  - ▶ spim is mips simulator developed by James Larus
  - ▶ RISC machine
  - ▶ Assembly level simulator: No need of assembler, linkers, or libraries
- Level 0 of spim machine descriptions
- Level 1 of spim machine descriptions

2 July 2012

Spim MD: Outline

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### Outline

Notes

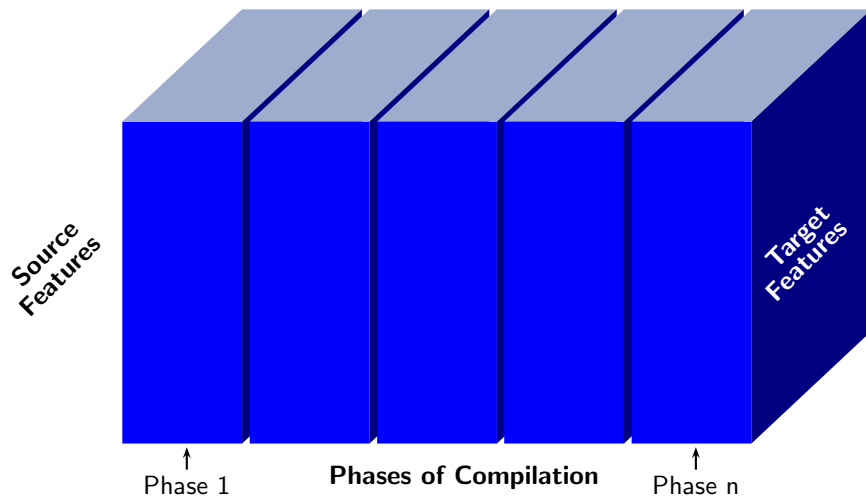


Part 1

## Systematic Construction of Machine Descriptions

Notes

### In Search of Modularity in Retargetable Compilation

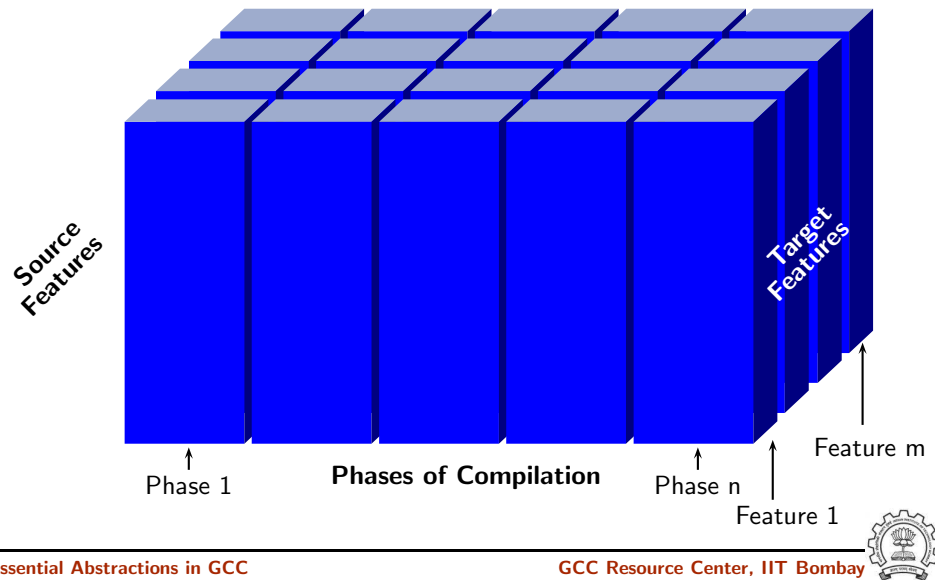


### In Search of Modularity in Retargetable Compilation

Notes



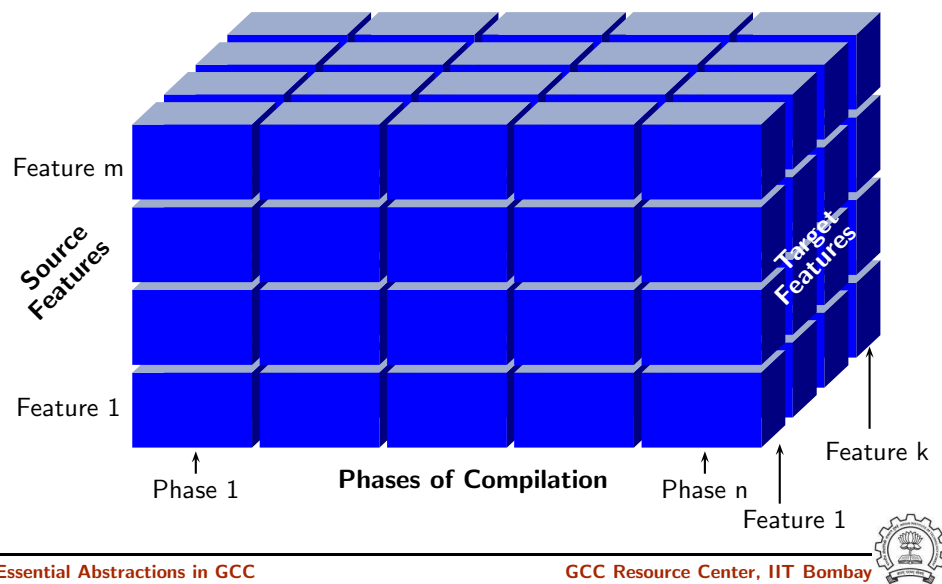
# In Search of Modularity in Retargetable Compilation



# In Search of Modularity in Retargetable Compilation

## Notes

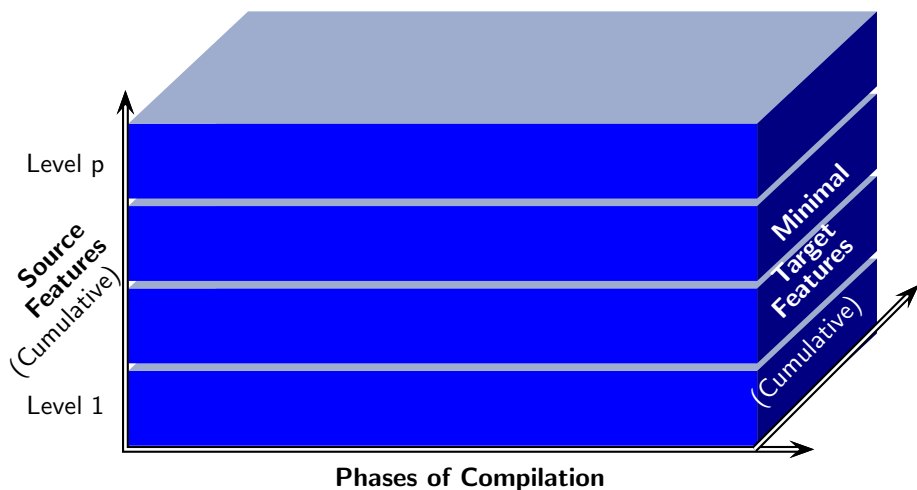
# In Search of Modularity in Retargetable Compilation



# In Search of Modularity in Retargetable Compilation

## Notes

# In Search of Modularity in Retargetable Compilation

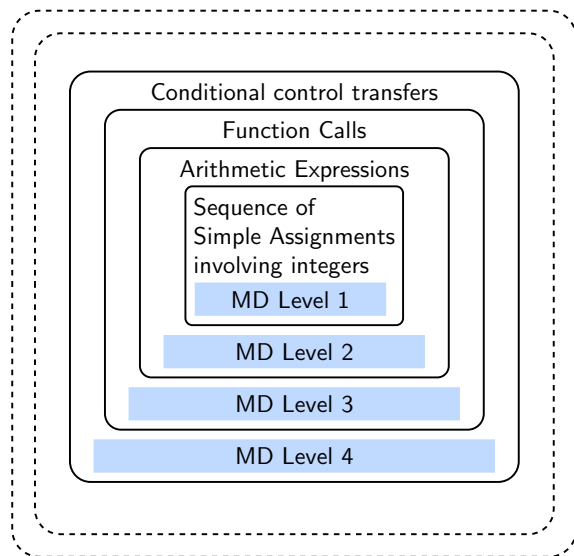


# In Search of Modularity in Retargetable Compilation

# Notes



# Systematic Development of Machine Descriptions



# Systematic Development of Machine Descriptions

# Notes



## Systematic Development of Machine Descriptions

- Define different levels of source language
- Identify the minimal information required in the machine description to support each level
  - ▶ Successful compilation of any program, and
  - ▶ correct execution of the generated assembly program.
- Interesting observations
  - ▶ It is the increment in the source language which results in understandable increments in machine descriptions rather than the increment in the target architecture.
  - ▶ If the levels are identified properly, the increments in machine descriptions are monotonic.



Part 2

*Retargeting GCC to Spim: A Recap*

## Systematic Development of Machine Descriptions

Notes



Notes

## Retargeting GCC to Spim

- Registering spim target with GCC build process
- Making machine description files available
- Building the compiler



## Registering Spim with GCC Build Process

We want to add multiple descriptions:

- Step 1. In the file `$(SOURCE_DIR)/config.sub`  
Add to the `case $basic_machine`
  - ▶ `spim*` in the part following  
`# Recognize the basic CPU types without company name.`
  - ▶ `spim*-*` in the part following  
`# Recognize the basic CPU types with company name.`



## Retargeting GCC to Spim

# Notes



## Registering Spim with GCC Build Process

# Notes



## Registering Spim with GCC Build Process

- Step 2a. In the file `$(SOURCE_D)/gcc/config.gcc`

In case `${target}` used for defining `cpu_type`, i.e. after the line

```
# Set default cpu_type, tm_file, tm_p_file and xm_file ...
```

add the following case

```
spim*-*-*)
  cpu_type=spim
  ;;
```

This says that the machine description files are available in the directory `$(SOURCE_D)/gcc/config/spim`.



## Registering Spim with GCC Build Process

- Step 2b. In the file `$(SOURCE_D)/gcc/config.gcc`

Add the following in the case `${target}` for

```
# Support site-specific machine types.
```

```
spim*-*-*)
  gas=no
  gnu_ld=no
  file_base="`echo ${target} | sed 's/-.*$//`'"
  tm_file="${cpu_type}/${file_base}.h"
  md_file="${cpu_type}/${file_base}.md"
  out_file="${cpu_type}/${file_base}.c"
  tm_p_file="${cpu_type}/${file_base}-protos.h"
  echo ${target}
  ;;
```



## Registering Spim with GCC Build Process

# Notes



## Registering Spim with GCC Build Process

# Notes



## Building a Cross-Compiler for Spim

- Normal cross compiler build process attempts to use the generated `cc1` to compile the emulation libraries (`LIBGCC`) into executables using the assembler, linker, and archiver.
- We are interested in only the `cc1` compiler.  
Add a new target in the `Makefile.in`

```
.PHONY: cc1
cc1:
    make all-gcc TARGET-gcc=cc1$(exeext)
```



## Building a Cross-Compiler for Spim

- Create directories `${BUILD_D}` and in a tree not rooted at `${SOURCE_D}`.
- Change the directory to `${BUILD_D}` and execute the commands

```
$ cd ${BUILD_D}
$ ${SOURCE_D}/configure --target=spim<n>
$ make cc1
```

- Pray for 10 minutes :-)



## Building a Cross-Compiler for Spim

# Notes



## Building a Cross-Compiler for Spim

# Notes





## Level 0 of Spim Machine Descriptions

Notes

### Sub-levels of Level 0

Three sub-levels

- Level 0.0: Merely build GCC for spim simulator  
Does not compile any program (i.e. compilation aborts)
- Level 0.1: Compiles empty void functions

```
void fun(int p1, int p2)
{
    int v1, v2;
}
```

```
void fun()
{
    L: goto L;
}
```

- Level 0.2: Incorporates complete activation record structure  
Required for Level 1

### Sub-levels of Level 0

Notes



## Category of Macros in Level 0

Category	Level 0.0	Level 0.1	Level 0.2
Memory Layout	complete	complete	complete
Registers	partial	partial	complete
Addressing Modes	none	partial	partial
Activation Record Conventions	dummy	dummy	complete
Calling Conventions	dummy	dummy	partial
Assembly Output Format	dummy	partial	partial

- Complete specification of activation record in level 0.2 is not necessary but is provided to facilitate local variables in level 1.
- Complete specification of registers in level 0.2 follows the complete specification of activation record.



## Memory Layout Related Macros for Level 0

```
#define BITS_BIG_ENDIAN 0
#define BYTES_BIG_ENDIAN 0
#define WORDS_BIG_ENDIAN 0
#define UNITS_PER_WORD 4
#define PARM_BOUNDARY 32
#define STACK_BOUNDARY 64
#define FUNCTION_BOUNDARY 32

#define BIGGEST_ALIGNMENT 64
#define STRICT_ALIGNMENT 0
#define MOVE_MAX 4
#define Pmode SImode
#define FUNCTION_MODE SImode
#define SLOW_BYTE_ACCESS 0
#define CASE_VECTOR_MODE SImode
```



## Category of Macros in Level 0

# Notes

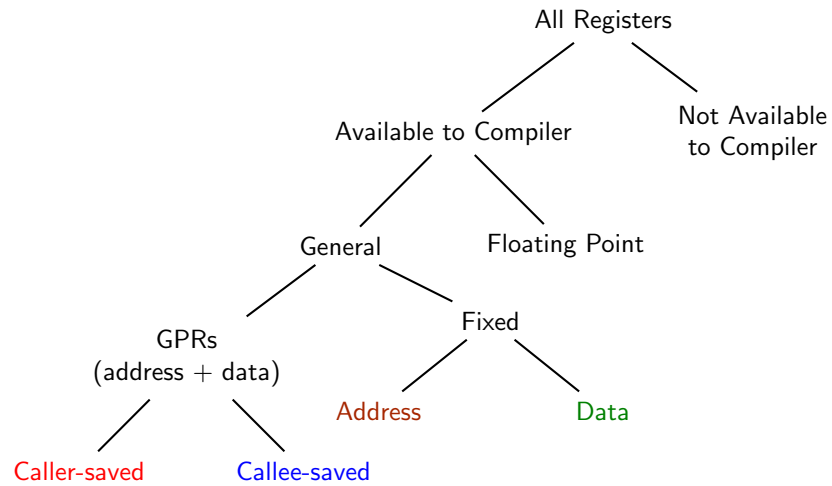


## Memory Layout Related Macros for Level 0

# Notes



## Register Categories for Spim



## Register Categories for Spim

# Notes



## Registers in Spim

\$zero	00	32		constant data
\$at	01	32		NA
\$v0	02	32,64	result	caller
\$v1	03	32	result	caller
\$a0	04	32,64	argument	caller
\$a1	05	32	argument	caller
\$a2	06	32,64	argument	caller
\$a3	07	32	argument	caller
\$t0	08	32,64	temporary	caller
\$t1	09	32	temporary	caller
\$t2	10	32,64	temporary	caller
\$t3	11	32	temporary	caller
\$t4	12	32,64	temporary	caller
\$t5	13	32	temporary	caller
\$t6	14	32,64	temporary	caller
\$t7	15	32	temporary	caller
\$s0	16	32,64	temporary	callee
\$s1	17	32	temporary	callee
\$s2	18	32,64	result	callee
\$s3	19	32	result	callee
\$s4	20	32,64	temporary	callee
\$s5	21	32	temporary	callee
\$s6	22	32,64	temporary	callee
\$s7	23	32	temporary	callee
\$t8	24	32,64	temporary	caller
\$t9	25	32	temporary	caller
\$k0	26	32,64		NA
\$k1	27	32		NA
\$gp	28	32,64	global pointer	address
\$sp	29	32	stack pointer	address
\$fp	30	32,64	frame pointer	address
\$ra	31	32	return address	address



## Registers in Spim

# Notes



## Register Information in Level 0.2

```

$zero,$at
#define FIRST_PSEUDO_REGISTER 32

#define FIXED_REGISTERS \
/* not for global */ \
/* register allocation */ \
{ 1,1,0,0, 0,0,0,0, \
  0,0,0,0, 0,0,0,0, \
  0,0,0,0, 0,0,0,0, \
  0,0,1,1,1,1,1,1 }

#define CALL_USED_REGISTERS \
/* Caller-saved registers */ \
{ 1,1,1,1, 1,1,1,1, \
  1,1,1,1, 1,1,1,1, \
  0,0,0,0, 0,0,0,0, \
  1,1,1,1, 1,1,1,1 }

/* Register sizes */
#define HARD_REGNO_NREGS(R,M) \
((GET_MODE_SIZE (M) + \
  UNITS_PER_WORD - 1) \
 / UNITS_PER_WORD)

#define HARD_REGNO_MODE_OK(R,M) \
hard_regno_mode_ok (R, M)

#define MODES_TIEABLE_P(M1,M2) \
modes_tieable_p (M1,M2)

$s0 to $s7

```



## Register Information in Level 0.2

Notes



## Register Classes in Level 0.2

```

enum reg_class \
{ NO_REGS, CALLER_SAVED_REGS, \
  CALLEE_SAVED_REGS, BASE_REGS, \
  GENERAL_REGS, ALL_REGS, \
  LIM_REG_CLASSES \
};
#define N_REG_CLASSES \
LIM_REG_CLASSES
#define REG_CLASS_NAMES \
{ "NO_REGS", "CALLER_SAVED_REGS", \
  "CALLEE_SAVED_REGS", \
  "BASE_REGS", "GEN_REGS", \
  "ALL_REGS" \
}
#define REG_CLASS_CONTENTS \
/* Register numbers */ \
{ 0x00000000,0xff00ffff, \
  0x00ff0000,0xf0000000, \
  0x0cfffff3,0xfffffff3 }

address registers
#define REGNO_REG_CLASS(REGNO) \
regno_reg_class(REGNO)
#define BASE_REG_CLASS \
BASE_REGS
#define INDEX_REG_CLASS NO_REGS
#define REG_CLASS_FROM_LETTER(c) \
NO_REGS
#define REGNO_OK_FOR_BASE_P(R) 1
#define REGNO_OK_FOR_INDEX_P(R) 0
#define PREFERRED_RELOAD_CLASS(X,C) \
CLASS
/* Max reg required for a class */
#define CLASS_MAX_NREGS(C, M) \
((GET_MODE_SIZE (M) + \
  UNITS_PER_WORD - 1) \
 / UNITS_PER_WORD)
#define LEGITIMATE_CONSTANT_P(x) \
legitimate_constant_p(x)

```



## Register Classes in Level 0.2

Notes



## function calling conventions

pass arguments on stack. return values goes in register \$v0 (in level 1).

```
#define RETURN_POPS_ARGS(FUN, TYPE, SIZE) 0
#define FUNCTION_ARG(CUM, MODE, TYPE, NAMED) 0
#define FUNCTION_ARG_REGNO_P(r) 0
/*Data structure to record the information about args passed in
 *registers. Irrelevant in this level so a simple int will do. */
#define CUMULATIVE_ARGS int
#define INIT_CUMULATIVE_ARGS(CUM, FNTYPE, LIBNAME, FNDECL, NAMED_ARGS) \
{ CUM = 0; }
#define FUNCTION_ARG_ADVANCE(cum, mode, type, named) cum++
#define FUNCTION_VALUE(valtype, func) function_value()
#define FUNCTION_VALUE_REGNO_P(REGN) ((REGN) == 2)
```

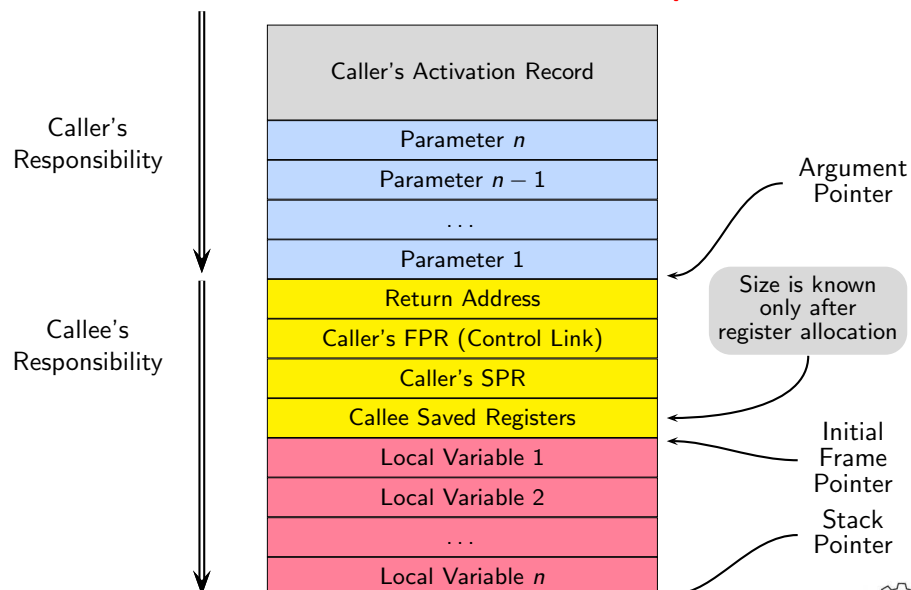


## function calling conventions

## Notes



## Activation Record Structure in Spim



## Activation Record Structure in Spim

## Notes



## Minimizing Registers for Accessing Activation Records

Reduce four pointer registers (stack, frame, args, and hard frame) to fewer registers.

```
#define ELIMINABLE_REGS
  {{FRAME_POINTER_REGNUM,    STACK_POINTER_REGNUM},
   {FRAME_POINTER_REGNUM,    HARD_FRAME_POINTER_REGNUM},
   {ARG_POINTER_REGNUM,     STACK_POINTER_REGNUM},
   {HARD_FRAME_POINTER_REGNUM, STACK_POINTER_REGNUM}}
}
```

/Recomputes new offsets, after eliminating./

```
#define INITIAL_ELIMINATION_OFFSET(FROM, TO, VAR)
  (VAR) = initial_elimination_offset(FROM, TO)
```



## Specifying Activation Record

```
#define STARTING_FRAME_OFFSET  starting_frame_offset ()
```

```
#define FIRST_PARM_OFFSET(FUN) 0
```

```
#define STACK_POINTER_REGNUM 29
```

```
#define FRAME_POINTER_REGNUM 1
```

```
#define HARD_FRAME_POINTER_REGNUM 30
```

```
#define ARG_POINTER_REGNUM HARD_FRAME_POINTER_REGNUM
```

```
#define FRAME_POINTER_REQUIRED 0
```



## Minimizing Registers for Accessing Activation Records

Notes



## Specifying Activation Record

Notes



## Level 0.0 Machine Description File

Empty :-)



## Operations in Level 0

Operations	Level 0.0	Level 0.1	Level 0.2
JUMP direct	dummy	actual	actual
JUMP indirect	dummy	dummy	dummy
NOP	dummy	actual	actual
MOV	not required	partial	partial
RETURN	not required	partial	partial

spim0.0.c

```

rtx gen_jump (...)
{ return 0; }
rtx gen_indirect_jump (...)
{ return 0; }
rtx gen_nop ()
{ return 0; }

```

spim0.0.h

```
#define CODE_FOR_indirect_jump 8
```

spim0.2.md

```

(define_insn "jump"
  [(set (pc)
        (label_ref (match_operand 0 "" ""))
        )]
  ""
  "j %10"
  )

```



## Level 0.0 Machine Description File

Notes



## Operations in Level 0

Notes



## Operations in Level 0

Operations	Level 0.0	Level 0.1	Level 0.2
JUMP direct	dummy	actual	actual
JUMP indirect	dummy	dummy	dummy
NOP	dummy	actual	actual
MOV	not required	partial	partial
RETURN	not required	partial	partial

Only define\_expand. No define\_insn.

```
(define_expand "movsi"
  [(set (match_operand:SI 0 "nonimmediate_operand" "")
        (match_operand:SI 1 "general_operand" ""))]
  ""
  {
    if (GET_CODE(operands[0]) == MEM && GET_CODE(operands[1]) != REG)
    {
      if (can_create_pseudo_p())
      {
        operands[1] = force_reg(SImode, operands[1]);
      }
    }
  })
```

spim0.2.md



## Operations in Level 0

Notes



## Operations in Level 0

Operations	Level 0.0	Level 0.1	Level 0.2
JUMP direct	dummy	actual	actual
JUMP indirect	dummy	dummy	dummy
NOP	dummy	actual	actual
MOV	not required	partial	partial
RETURN	not required	partial	partial

```
(define_expand "epilogue"
  [(clobber (const_int 0))]
  ""
  {
    spim_epilogue();
    DONE;
  })
(define_insn "IITB_return"
  [(return)]
  ""
  "jr \\$ra"
)
```

spim0.2.md

Only return.  
No epilogue code.

spim0.2.c

```
void spim_epilogue()
{
  emit_insn(gen_IITB_return());
}
```



## Operations in Level 0

Notes





## Operations in Level 0

Operations	Level 0.0	Level 0.1	Level 0.2
JUMP direct	dummy	actual	actual
JUMP indirect	dummy	dummy	dummy
NOP	dummy	actual	actual
MOV	not required	partial	partial
RETURN	not required	partial	partial

```
(define_insn "nop"
  [(const_int 0)]
  ""
  "nop"
)
```



Part 4

*Level 1 of Spim Machine Descriptions*

## Operations in Level 0

Notes



Notes

## Increments for Level 1

- Addition to the source language
  - ▶ Assignment statements involving integer constant, integer local or global variables.
  - ▶ Returning values. (No calls, though!)
- Changes in machine descriptions
  - ▶ Minor changes in macros required for level 0  
\$zero now belongs to new class Assembly output needs to change
  - ▶ Some function bodies expanded
  - ▶ New operations included in the .md file

diff -w shows the changes!



## Operations Required in Level 1

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow Src$	$R_i \leftarrow R_j$	move rj, ri	
	$R \leftarrow M_{global}$	lw r, m	
	$R \leftarrow M_{local}$	lw r, c(\$fp)	
	$R \leftarrow C$	li r, c	
	$M \leftarrow R$	sw r, m	
RETURN $Src$	RETURN $Src$	$\$v0 \leftarrow Src$ j \$ra	level 0
$Dest \leftarrow Src_1 + Src_2$	$R_i \leftarrow R_j + R_k$	add ri, rj, rk	
	$R_i \leftarrow R_j + C$	addi ri, rj, c	



## Increments for Level 1

# Notes



## Operations Required in Level 1

# Notes



## Move Operations in spim1.md

- Multiple primitive variants require us to map a single operation in IR to multiple RTL patterns  
⇒ use `define_expand`
- Ensure that the second operand is in a register

```
(define_expand "movsi"
  [(set (match_operand:SI 0 "nonimmediate_operand" "")
        (match_operand:SI 1 "general_operand" ""))
  ]
  ""
  { if (GET_CODE(operands[0]) == MEM &&
        GET_CODE(operands[1]) != REG &&
        (can_create_pseudo_p()) /* force conversion only */
        /* before register allocation */
        { operands[1] = force_reg(SImode, operands[1]); }
    }
  )
```



## Move Operations in spim1.md

# Notes



## Move Operations in spim1 Compiler for Assignment a = b

```
(define_expand "movsi"
  [(set (match_operand:SI 0 "nonimmediate_operand" "")
        (match_operand:SI 1 "general_operand" ""))
  ]
  ""
  { if (GET_CODE(operands[0]) == MEM &&
        GET_CODE(operands[1]) != REG &&
        (can_create_pseudo_p()) /* force conversion only */
        /* before register allocation */
        { operands[1] = force_reg(SImode, operands[1]); }
    }
  )

(insn 6 5 7 3 t.c:25 (set (reg:SI 38)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [0xffffffff])) [0 b+0 S4 A32])) -1 (nil))
(insn 7 6 8 3 t.c:25 (set (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
  (const_int -8 [0xffffffff8])) [0 a+0 S4 A32])
  (reg:SI 38)) -1 (nil))
```



## Move Operations in spim1 Compiler for Assignment a = b

# Notes



## Move Operations in spim1.md

- Load from Memory  $R \leftarrow M$

```
(define_insn "*load_word"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (match_operand:SI 1 "memory_operand" "m"))]
  ""
  "lw \t%0, %m1"
  )
```

- Load Constant  $R \leftarrow C$

```
(define_insn "*constant_load"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (match_operand:SI 1 "const_int_operand" "i"))]
  ""
  "li \t%0, %c1"
  )
```



## Move Operations in spim1.md

# Notes



## Move Operations in spim1.md

- Register Move  $R_i \leftarrow R_j$

```
(define_insn "*move_regs"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (match_operand:SI 1 "register_operand" "r"))]
  ""
  "move \t%0,%1"
  )
```

- Store into  $M \leftarrow R$

```
(define_insn "*store_word"
  [(set (match_operand:SI 0 "memory_operand" "=m")
        (match_operand:SI 1 "register_operand" "r"))]
  ""
  "sw \t%1, %m0"
  )
```



## Move Operations in spim1.md

# Notes



## Code Generation in spim1 Compiler for Assignment a = b

- RTL statements

```
(insn 6 5 7 3 t.c:25 (set (reg:SI 38)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [0xffffffffc])) [0 b+0 S4 A32])) -1 (nil))
(insn 7 6 8 3 t.c:25 (set (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-
  (const_int -8 [0xffffffff8])) [0 a+0 S4 A32])
  (reg:SI 38)) -1 (nil))
```

- Generated Code

```
lw $v0, -16($fp)
sw $v0, -20($fp)
```

# Notes



## Using register \$zero for constant 0

- Introduce new register class `zero_register_operand` in `spim1.h` and define `move_zero`

```
(define_insn "IITB_move_zero"
  [(set (match_operand:SI 0 "nonimmediate_operand" "=r,m")
    (match_operand:SI 1 "zero_register_operand" "z,z"))
  ])
  ""
  "@
  move \t%0,%1
  sw \t%1, %m0"
)
```

- How do we get `zero_register_operand` in an RTL?

# Notes

## Code Generation in spim1 Compiler for Assignment a = b



## Using register \$zero for constant 0



## Using register \$zero for constant 0

- Use `define_expand "movsi"` to get `zero_register_operand` in an RTL

```
if (GET_CODE(operands[1]) == CONST_INT && INTVAL(operands[1]) == 0)
{
  emit_insn(gen_IITB_move_zero(operands[0],
                               gen_rtx_REG(SImode, 0)));
  DONE;
}
else /* Usual processing */
```

- DONE says do not generate the RTL template associated with "movsi"
- required template is generated by `emit_insn(gen_IITB_move_zero(...))`



## Supporting Addition in Level 1

```
(define_insn "addsi3"
  [(set (match_operand:SI 0 "register_operand" "=r,r")
        (plus:SI (match_operand:SI 1 "register_operand" "r,r")
                 (match_operand:SI 2 "nonmemory_operand" "r,i")))
  ])
  ""
  "@
  add \t%0, %1, %2
  addi \t%0, %1, %c2"
)
```

- Constraints combination 1 of three operands: R, R, R
- Constraints combination 2 of three operands: R, R, C



## Using register \$zero for constant 0

# Notes



## Supporting Addition in Level 1

# Notes



## Comparing movsi and addsi3

- movsi uses define\_expand whereas addsi3 uses combination of operands
- Why not use constraints for movsi too?
- Combination of operands is used during pattern matching and not during expansion
  - ▶ We will need to support memory as both source and destination
  - ▶ Will also allow memory to memory move in RTL  
We will not know until assembly emission which one is a load instruction and which one is a store instruction



Part 5

*Conclusions*

## Comparing movsi and addsi3

Notes



Notes

## Conclusions

- Incremental construction of machine description files is very instructive
- Increments in machine descriptions is governed by increments in source language
- Machine characteristics need to be specified in C macros and C functions
  - ▶ Does not seem amenable to incremental construction
  - ▶ Seems difficult to a novice
- Specifying instructions seems simpler and more systematic
  - ▶ Is amenable to incremental construction
  - ▶ The concept of minimal machine descriptions is very useful
- `define_insn` and `define_expand` are the main constructs used on machine descriptions



Part 6

*Constructs Supported in Level 2*

## Conclusions

Notes



Notes



## Arithmetic Operations Required in Level 2

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow Src_1 - Src_2$	$R_i \leftarrow R_j - R_k$	sub ri, rj, rk	level 2
$Dest \leftarrow -Src$	$R_i \leftarrow -R_j$	neg ri, rj	
$Dest \leftarrow Src_1 / Src_2$	$R_i \leftarrow R_j / R_k$	div rj, rk mflo ri	
$Dest \leftarrow Src_1 \% Src_2$	$R_i \leftarrow R_j \% R_k$	rem ri, rj, rk	
$Dest \leftarrow Src_1 * Src_2$	$R_i \leftarrow R_j * R_k$	mul ri, rj, rk	



## Bitwise Operations Required in Level 2

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow Src_1 \ll Src_2$	$R_i \leftarrow R_j \ll R_k$	sllv ri, rj, rk	level 2
	$R_i \leftarrow R_j \ll C_5$	sll ri, rj, c	
$Dest \leftarrow Src_1 \gg Src_2$	$R_i \leftarrow R_j \gg R_k$	srav ri, rj, rk	
	$R_i \leftarrow R_j \gg C_5$	sra ri, rj, c	
$Dest \leftarrow Src_1 \& Src_2$	$R_i \leftarrow R_j \& R_k$	and ri, rj, rk	
	$R_i \leftarrow R_j \& C$	andi ri, rj, c	
$Dest \leftarrow Src_1   Src_2$	$R_i \leftarrow R_j   R_k$	or ri, rj, rk	
	$R_i \leftarrow R_j   C$	ori ri, rj, c	
$Dest \leftarrow Src_1 \wedge Src_2$	$R_i \leftarrow R_j \wedge R_k$	xor ri, rj, rk	
	$R_i \leftarrow R_j \wedge C$	xori ri, rj, c	
$Dest \leftarrow \sim Src$	$R_i \leftarrow \sim R_j$	not ri, rj	



## Arithmetic Operations Required in Level 2

Notes



## Bitwise Operations Required in Level 2

Notes



## Divide Operation in spim2.md using define\_insn

- For division, the spim architecture imposes use of multiple asm instructions for single operation.
- Two ASM instructions are emitted using single RTL pattern

```
(define_insn "divsi3"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (div:SI (match_operand:SI 1 "register_operand" "r")
                (match_operand:SI 2 "register_operand" "r"))
        )]
  ""
  "div\t%1, %2\n\tmflo\t%0"
)
```



## Advantages/Disadvantages of using define\_insn

- Very simple to add the pattern
- Primitive target feature represented as single insn pattern in .md
- Unnecessary atomic grouping of instructions
- May hamper optimizations in general, and instruction scheduling, in particular



## Divide Operation in spim2.md using define\_insn

# Notes



## Advantages/Disadvantages of using define\_insn

# Notes



## Divide Operation in spim2.md using define\_expand

- The RTL pattern can be expanded into two different RTLs.

```
(define_expand "divsi3"
  [(parallel[(set (match_operand:SI 0 "register_operand" "")
                  (div:SI (match_operand:SI 1 "register_operand" "")
                          (match_operand:SI 2 "register_operand" ""))
                  )
            (clobber (reg:SI 26))
            (clobber (reg:SI 27))]]]
  ""
  {
    emit_insn(gen_IITB_divide(gen_rtx_REG(SImode,26),
                             operands[1], operands[2]));
    emit_insn(gen_IITB_move_from_lo(operands[0],
                                     gen_rtx_REG(SImode,26)));

    DONE;
  }
)
```



## Divide Operation in spim2.md using define\_expand

# Notes



## Divide Operation in spim2.md using define\_expand

- Divide pattern equivalent to div instruction in architecture.

```
(define_insn "IITB_divide"
  [(parallel[(set (match_operand:SI 0 "LO_register_operand" "=q")
                  (div:SI (match_operand:SI 1 "register_operand" "r")
                          (match_operand:SI 2 "register_operand" "r"))
                  )
            (clobber (reg:SI 27))]]]
  ""
  "div t%1, %2"
)
```



## Divide Operation in spim2.md using define\_expand

# Notes



## Divide Operation in spim2.md using define\_expand

- Moving contents of special purpose register LO to/from general purpose register

```
(define_insn "IITB_move_from_lo"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (match_operand:SI 1 "LO_register_operand" "q"))]
  ""
  "mflo \\t%0"
  )

(define_insn "IITB_move_to_lo"
  [(set (match_operand:SI 0 "LO_register_operand" "=q")
        (match_operand:SI 1 "register_operand" "r"))]
  ""
  "mtlo \\t%1"
  )
```



## Divide Operation in spim2.md using define\_expand

# Notes



## Divide Operation in spim2.md using define\_expand

- Divide pattern equivalent to div instruction in architecture.

```
(define_insn "modsi3"
  [(parallel [(set (match_operand:SI 0 "register_operand" "=r")
                  (mod:SI (match_operand:SI 1 "register_operand" "r")
                          (match_operand:SI 2 "register_operand" "r")))
             (clobber (reg:SI 26))
             (clobber (reg:SI 27))])]
  ""
  "rem \\t%0, %1, %2"
  )
```



## Divide Operation in spim2.md using define\_expand

# Notes



## Advantages/Disadvantages of Using `define_expand` for Division

- Two instructions are separated out at GIMPLE to RTL conversion phase
- Both instructions can undergo all RTL optimizations independently
- C interface is needed in md
- Compilation becomes slower and requires more space



## Divide Operation in `spim2.md` using `define_split`

```
(define_split
  [(parallel [(set (match_operand:SI 0 "register_operand" "")
    (div:SI (match_operand:SI 1 "register_operand" "")
      (match_operand:SI 2 "register_operand" ""))
    )
    (clobber (reg:SI 26))
    (clobber (reg:SI 27))]]]
  ""
  [(parallel [(set (match_dup 3)
    (div:SI (match_dup 1)
      (match_dup 2)))
    (clobber (reg:SI 27))]]]
    [(set (match_dup 0)
      (match_dup 3))
    [(parallel [
      (set (match_operand:SI 0 "LO_register_operand" "=q")
        (div:SI (match_operand:SI 1 "register_operand" "r")
          (match_operand:SI 2 "register_operand" "r")))
      (clobber (reg:SI 27))]]]
    [(set (match_operand:SI 0 "register_operand" "=r")
      (match_operand:SI 1 "LO_register_operand" "q"))]]
  ]
  "operands[3]=gen_rtx_REG(SImode,26); "
```



## Advantages/Disadvantages of Using `define_expand` for Division

# Notes



## Divide Operation in `spim2.md` using `define_split`

# Notes



## Constructs Supported in Level 3

### Operations Required in Level 3

Operation	Primitive Variants	Implementation	Remark
$Dest \leftarrow fun(P_1, \dots, P_n)$	call $L_{fun}, n$	lw $r_i, [SP+c1]$	Level 1
		sw $r_i, [SP]$	
		:	
		lw $r_i, [SP+c2]$	
		sw $r_i, [SP-n*4]$	
		jal L	New
		$Dest \leftarrow \$v0$	level 1
$fun(P_1, P_2, \dots, P_n)$	call $L_{fun}, n$	lw $r_i, [SP+c1]$	Level 1
		sw $r_i, [SP]$	
		:	
		lw $r_i, [SP+c2]$	
		sw $r_i, [SP-n*4]$	
		jal L	New



### Operations Required in Level 3



## Call Operation in spim3.md

```
(define_insn "call"
  [(call (match_operand:SI 0 "memory_operand" "=m")
         (match_operand:SI 1 "immediate_operand" "i"))
   (clobber (reg:SI 31))]
  ""
  "*"
  return emit_asm_call(operands,0);
  ""
  )
```



## Call Operation in spim3.md

```
(define_insn "call_value"
  [(set (match_operand:SI 0 "register_operand" "=r")
        (call (match_operand:SI 1 "memory_operand" "m")
              (match_operand:SI 2 "immediate_operand" "i"))))
   (clobber (reg:SI 31))]
  ""
  "*"
  return emit_asm_call(operands,1);
  ""
  )
```



## Call Operation in spim3.md

# Notes



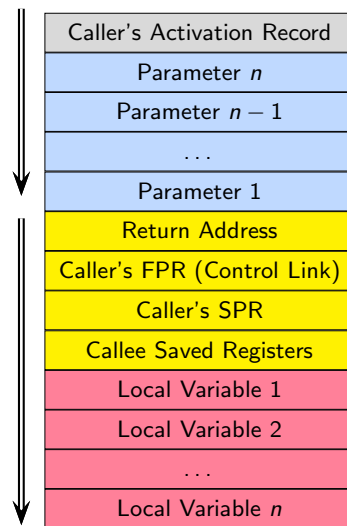
## Call Operation in spim3.md

# Notes



## Activation Record Generation during Call

- Operations performed by caller
  - ▶ Push parameters on stack.
  - ▶ Load return address in return address register.
  - ▶ Transfer control to Callee.
- Operations performed by callee
  - ▶ Push Return address stored by caller on stack.
  - ▶ Push caller's Frame Pointer Register.
  - ▶ Push caller's Stack Pointer.
  - ▶ Save callee saved registers, if used by callee.
  - ▶ Create local variables frame.
  - ▶ Start callee body execution.



## Activation Record Generation during Call

# Notes

## Prologue in spim3.md

```

(define_expand "prologue"
  [(clobber (const_int 0))]
  "")
{
  spim_prologue();
  DONE;
})

(set (mem:SI (reg:SI $sp))
  (reg:SI 31 $ra))

(set (mem:SI (plus:SI (reg:SI $sp)
  (const_int -4 )))
  (reg:SI $sp))

(set (mem:SI (plus:SI (reg:SI $sp)
  (const_int -8 )))
  (reg:SI $fp))

(set (reg:SI $fp)
  (reg:SI $sp))

(set (reg:SI $sp)
  (plus:SI (reg:SI $fp)
  (const_int -36)))

```



## Prologue in spim3.md

# Notes





**Epilogue in spim3.md**

```
(define_expand "epilogue"
  [(clobber (const_int 0))]
  ""
  spim_epilogue();
  DONE;
)

(set (reg:SI $sp)
  (reg:SI $fp))

(set (reg:SI $fp)
  (mem:SI (plus:SI (reg:SI $sp)
    (const_int -8 ))))

(set (reg:SI $ra)
  (mem:SI (reg:SI $sp)))

(parallel [
  (return)
  (use (reg:SI $ra))])
```

**Epilogue in spim3.md****Notes****Notes**

## Operations Required in Level 4

Operation	Primitive Variants	Implementation	Remark
$Src_1 < Src_2 ?$ goto L : PC	$CC \leftarrow R_i < R_j$ $CC < 0 ?$ goto L : PC	blt $r_i, r_j, L$	
$Src_1 > Src_2 ?$ goto L : PC	$CC \leftarrow R_i > R_j$ $CC > 0 ?$ goto L : PC	bgt $r_i, r_j, L$	
$Src_1 \leq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \leq R_j$ $CC \leq 0 ?$ goto L : PC	ble $r_i, r_j, L$	
$Src_1 \geq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \geq R_j$ $CC \geq 0 ?$ goto L : PC	bge $r_i, r_j, L$	



## Operations Required in Level 4

Operation	Primitive Variants	Implementation	Remark
$Src_1 == Src_2 ?$ goto L : PC	$CC \leftarrow R_i == R_j$ $CC == 0 ?$ goto L : PC	beq $r_i, r_j, L$	
$Src_1 \neq Src_2 ?$ goto L : PC	$CC \leftarrow R_i \neq R_j$ $CC \neq 0 ?$ goto L : PC	bne $r_i, r_j, L$	



## Operations Required in Level 4

Notes



## Operations Required in Level 4

Notes



**Conditional Branch Instruction in spim4.md**

```
(define_insn "cbranchsi4"
  [(set (pc)
        (if_then_else
         (match_operator:SI 0 "comparison_operator"
          [(match_operand:SI 1 "register_operand" "")
           (match_operand:SI 2 "register_operand" "")])
         (label_ref (match_operand 3 "" ""))
         (pc)))]
  ""
  "*"
  return conditional_insn(GET_CODE(operands[0]),operands);
  ""
  )
```

**Support for Branch pattern in spim4.c**

```
char *
conditional_insn (enum rtx_code code, rtx operands[])
{
  switch(code)
  {
    case EQ: return "beq %1, %2, %l3";
    case NE: return "bne %1, %2, %l3";
    case GE: return "bge %1, %2, %l3";
    case GT: return "bgt %1, %2, %l3";
    case LT: return "blt %1, %2, %l3";
    case LE: return "ble %1, %2, %l3";
    case GEU: return "bgeu %1, %2, %l3";
    case GTU: return "bg tu %1, %2, %l3";
    case LTU: return "bltu %1, %2, %l3";
    case LEU: return "bleu %1, %2, %l3";
    default: /* Error. Issue ICE */
  }
}
```

**Conditional Branch Instruction in spim4.md**

Notes

**Support for Branch pattern in spim4.c**

Notes



**Alternative for Branch: Conditional compare in spim4.md**

```
(define_code_iterator cond_code
  [lt ltu eq ge geu gt gtu le leu ne])

(define_expand "cmpsi"
  [(set (cc0) (compare
             (match_operand:SI 0 "register_operand" "")
             (match_operand:SI 1 "nonmemory_operand" "")))]
  ""
  {
    compare_op0=operands[0];
    compare_op1=operands[1];
    DONE;
  }
)
```

**Alternative for Branch: Branch pattern in spim4.md**

```
(define_expand "b<code>"
  [(set (pc) (if_then_else (cond_code:SI (match_dup 1)
                             (match_dup 2))
                           (label_ref (match_operand 0 "" ""))
                           (pc)))]
  ""
  {
    operands[1]=compare_op0;
    operands[2]=compare_op1;
    if (immediate_operand(operands[2], SImode))
    {
      operands[2]=force_reg(SImode, operands[2]);
    }
  }
)
```

**Alternative for Branch: Conditional compare in spim4.md**

Notes

**Alternative for Branch: Branch pattern in spim4.md**

Notes



**Alternative for Branch: Branch pattern in spim4.md**

```
(define_insn "*insn_b<code>"
  [(set (pc)
        (if_then_else
         (cond_code:SI
          (match_operand:SI 1 "register_operand" "r")
          (match_operand:SI 2 "register_operand" "r"))
         (label_ref (match_operand 0 "" ""))
         (pc)))]
  ""
  "*"
  return conditional_insn(<CODE>,operands);
  "
```

**Alternative for Branch: Branch pattern in spim4.md**

# Notes

