Introduction to Hacking PostgreSQL

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1 Prerequisites
   ■ Why Should You Hack On PostgreSQL?
   ■ What Skills Will You Need?
   ■ What Tools Should You Use?

2 The Architecture of PostgreSQL
   ■ System Architecture
   ■ Components of the Backend

3 Common Code Conventions
   ■ Memory Management
   ■ Error Handling

4 Community Processes

5 Sample Patch

6 Conclusion
Outline

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Why Hack On PostgreSQL?

Possible Reasons

- Databases are fun!
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- Contribute to the community
  - We need more reviewers
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- Commercial opportunities
Skills

Essential

- Some knowledge of C
  - Fortunately, C is easy
- Some familiarity with Unix and basic Unix programming
  - Postgres development on Win32 is increasingly feasible

Helpful, but not essential

- Unix systems programming
- DBMS internals
- Autotools-foo
- Performance analysis

...depending on what you want to hack on
The Basics

$CC, Bison, Flex, CVS, autotools

- Configure flags: enable-depend, enable-debug, enable-cassert
- Consider `CFLAGS=-00` for easier debugging (and faster builds)
  - With GCC, this suppresses some important warnings
Development Tools

The Basics

$CC, Bison, Flex, CVS, autotools

- Configure flags: enable-depend, enable-debug, enable-cassert
- Consider CFLAGS=-O0 for easier debugging (and faster builds)
  - With GCC, this suppresses some important warnings

Indexing The Source

- A tool like tags, cscope or glimpse is essential when navigating any large code base
  - “What is the definition of this function/type?”
  - “What are all the call-sites of this function?”
  - src/tools/make_[ce]tags
Other Tools

- A debugger is often necessary: most developers use gdb
  - Or a front-end like ddd
  - Even MSVC?
- ccache and distcc are useful, especially on slower machines
- valgrind is useful for debugging memory errors and memory leaks in client apps
  - Not as useful for finding backend memory leaks
Other Tools

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Profiling

- `gprof` is the traditional choice; various bugs and limitations
  - Use `--enable-profiling` to reduce the pain
- `callgrind` works well, nice UI (`kcachegrind`)
- `oprofile` is good at system-level performance analysis
- DTrace
The DocBook toolchain is less than perfect
The DocBook toolchain is less than perfect

I don’t know of a good SGML editor, other than Emacs
- Writing DocBook markup by hand is labour-intensive but not hard: copy conventions of nearby markup
- `make check` does a quick syntax check
- `make draft` is useful for previewing changes
Patch Management

- Most development is done by mailing around patches
  - `echo "diff -c -N -p" >> ~/.cvsrc`
  - `cvs diff > ~/my_patch-vN.patch`
- `interdiff` is a useful tool: “exactly what did I change between v5 and v6?”
- Remote `cvs` is slow: setup a local mirror of the CVS repository
  - `cvsup`, `csup`, `rsync`, `svnsync` (soon!)
- To include newly-added files in a CVS diff, either use a local CVS mirror or `cvsutils`
- For larger projects: akpm’s Quilt, or a distributed VCS
  - Postgres-R uses Monotone
  - Recommended: git tree at repo.or.cz/w/PostgreSQL.git
If you’re not using a good programmer’s text editor, start
Teach your editor to obey the Postgres coding conventions:
  - Hard tabs, with a tab width of 4 spaces
  - Similar to Allman/BSD style; just copy the surrounding code
Using the Postgres coding conventions makes it more likely that your patch will be promptly reviewed and applied
Useful Texts

  - http://www.wiscorp.com/SQLStandards.html (“draft”)
  - There are some books and presentations that are more human-readable
  - There’s a samizdat plaintext version of SQL-92

- SQL references for Oracle, DB2, …

- A textbook on the design of database management systems
  - I personally like *Database Management Systems* by Ramakrishnan and Gehrke

- Books on the toolchain (C, Yacc, autotools, …) and operating system kernels
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6. Conclusion
The Postmaster

Lifecycle

1. Initialize essential subsystems; perform XLOG recovery to restore the database to a consistent state
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2. Attach to shared memory segment (SysV IPC), initialize shared data structures.
3. Fork off daemon processes: autovacuum launcher, stats daemon, bgwriter, syslogger.

For each new connection, spawn a backend. Periodically check for child death, launch replacements or perform recovery.
The Postmaster

Lifecycle

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3. Fork off daemon processes: autovacuum launcher, stats daemon, bgwriter, syslogger
4. Bind to TCP socket, listen for incoming connections
   - For each new connection, spawn a backend
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## Daemon Processes

### Types of Processes

<table>
<thead>
<tr>
<th>Process</th>
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</tr>
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<td><strong>autovacuum launcher</strong></td>
<td>Periodically start autovacuum workers</td>
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<td><strong>bgwriter</strong></td>
<td>Flush dirty buffers to disk, perform periodic checkpoints</td>
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<tr>
<td><strong>stats collector</strong></td>
<td>Accepts run-time stats from backends via UDP</td>
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<tr>
<td><strong>syslogger</strong></td>
<td>Collect log output from other processes, write to file(s)</td>
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<td><strong>normal backend</strong></td>
<td>Handles a single client session</td>
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- Inter-Process Communication
  - Most shared data is communicated via a shared memory segment
  - Signals, semaphores, and pipes also used as appropriate
  - Stats collector uses UDP on the loopback interface
  - Subprocesses inherit the state of the postmaster after fork()
Daemon Processes

Types of Processes

- **autovacuum launcher**: Periodically start autovacuum workers
- **bgwriter**: Flush dirty buffers to disk, perform periodic checkpoints
- **stats collector**: Accepts run-time stats from backends via UDP
- **syslogger**: Collect log output from other processes, write to file(s)
- **normal backend**: Handles a single client session

Inter-Process Communication

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Consequences

Advantages

- Address space protection: significantly harder for misbehaving processes to crash the entire DBMS
- IPC and modifications to shared data are explicit: all state is process-private by default
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### Advantages

- Address space protection: significantly harder for misbehaving processes to crash the entire DBMS
- IPC and modifications to shared data are **explicit**: all state is process-private by default

### Disadvantages

- Shared memory segment is **statically-sized** at startup
  - Managing arbitrarily-sized shared data is problematic
- Some shared operations can be awkward: e.g. using multiple processors to evaluate a single query
## Backend Lifecycle

1. Postmaster accepts a connection, forks a new backend, then closes its copy of the TCP socket
   - All communication occurs between backend and client
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2. Backend enters the “frontend/backend” protocol:
   1. Authenticate the client
   2. “Simple query protocol”: accept a query, evaluate it, return result set
   3. When the client disconnects, the backend exits
# Stages In Query Processing

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<th>Component</th>
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<td>1</td>
<td>The <em>parser</em></td>
<td>lex &amp; parse the query string</td>
</tr>
<tr>
<td>2</td>
<td>The <em>rewriter</em></td>
<td>apply rewrite rules</td>
</tr>
<tr>
<td>3</td>
<td>The <em>optimizer</em></td>
<td>determine an efficient query plan</td>
</tr>
<tr>
<td>4</td>
<td>The <em>executor</em></td>
<td>execute a query plan</td>
</tr>
<tr>
<td>5</td>
<td>The <em>utility processor</em></td>
<td>process DDL like CREATE TABLE</td>
</tr>
</tbody>
</table>
Lex and parse the query string submitted by the user

- **Lexing**: divide the input string into a sequence of *tokens*
  - Postgres uses GNU Flex

- **Parsing**: construct an abstract syntax tree (AST) from sequence of tokens
  - Postgres uses GNU Bison
  - The elements of the AST are known as *parse nodes*
The Parser

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- **Parsing**: construct an abstract syntax tree (AST) from sequence of tokens
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  - The elements of the AST are known as **parse nodes**
- Produces a “raw parsetree”: a linked list of parse nodes
  - Parse nodes are defined in `include/nodes/parsenodes.h`
- Typically a simple mapping between grammar productions and parse node structure
Semantic Analysis

- In the parser itself, only syntactic analysis is done; basic semantic checks are done in a subsequent “analysis phase”
- `parser/analyze.c` and related code under `parser/`
- Resolve column references, considering schema path and query context
  - `SELECT a, b, c FROM t1, t2, x.t3`  
    `WHERE x IN (SELECT t1 FROM b)`
- Verify that referenced schemas, tables and columns exist
- Check that the types used in expressions are consistent
- In general, check for errors that are impossible or difficult to detect in the parser itself
The analysis phase produces a Query, which is the query’s parse tree (Abstract Syntax Tree) with additional annotations. The rewriter applies rewrite rules, including view definitions. Input is a Query, output is zero or more Queries.

The planner takes a Query and produces a Plan, which encodes how the query should be executed.

- A query plan is a tree of Plan nodes, each describing a physical operation.
- Only needed for “optimizable” statements (INSERT, DELETE, SELECT, UPDATE).
Each node in the plan tree describes a physical operation

- Scan a relation, perform an index scan, join two relations, perform a sort, apply a predicate, perform projection, . . .
Executor

- Each node in the plan tree describes a physical operation:
  - Scan a relation, perform an index scan, join two relations, perform a sort, apply a predicate, perform projection, . . .
- The planner arranges the operations into a plan tree that describes the data flow between operations.
- Tuples flow from the leaves of the tree to the root:
  - Leaf nodes are scans: no input, produce a stream of tuples.
  - Joins are binary operators: accept two inputs (child nodes), produce a single output.
  - The root of the tree produces the query’s result set.
- Therefore, the executor is “trivial”: simply ask the root plan node to repeatedly produce result tuples.
Query Optimization

- SQL is (ostensibly) a declarative query language
  - The query specifies the properties the result set must satisfy, not the procedure the DBMS must follow to produce the result set
- For a typical SQL query, there are many equivalent query plans
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  - **scan types:** Seq scan, index scan, bitmap index scan
  - **join order:** Inner joins are commutative: reordered freely
  - **join types:** Sort-merge join, hash join, nested loops
  - **aggregation:** Hashed aggregation, aggregation by sorting
  - **predicates:** Predicate push down, evaluation order
  - **rewrites:** Subqueries and set operations → joins, outer joins → inner joins, function inlining, . . .
Tasks Of The Query Optimizer

Basic Optimizer Task

Of the many ways in which we could evaluate a query, which would be the cheapest to execute?
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Of the many ways in which we could evaluate a query, which would be the cheapest to execute?

Two Distinct Subproblems

1. Enumerate all the possible plans for a given query
2. Estimate the cost of a given query plan

In practice, too slow → do both steps at the same time
Stages in Query Optimization

The System R Algorithm

1. Rewrite the query to make it more amenable to optimization: pull up subqueries, rewrite IN clauses, simplify constant expressions, reduce outer joins, . . .
## Stages in Query Optimization

### The System R Algorithm

1. **Rewrite the query to make it more amenable to optimization:** pull up subqueries, rewrite `IN` clauses, simplify constant expressions, reduce outer joins, . . .

2. **Determine the *interesting* ways to access each base relation**
   - Remember the cheapest estimated access path, plus the cheapest path for each distinct sort order
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4. . . .
### Storage Management

#### Tables → Files

- Tables and indexes are stored in normal operating-system files.
- Each table/index divided into “segments” of at most 1GB.
- Tablespaces just control the filesystem location of segments.
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Files → Blocks

- Each file is divided into blocks of BLCKSZ bytes each
  - 8192 by default; compile-time constant
- Blocks consist of items, such as heap tuples (in tables), or index entries (in indexes), along with metadata
- Tuple versions uniquely identified by triple \((r, p, i)\): relation OID, block number, offset within block; known as “ctid”
Almost all I/O is not done directly: to access a page, a process asks the buffer manager for it.
The Buffer Manager

- Almost all I/O is not done directly: to access a page, a process asks the buffer manager for it
- The buffer manager implements a hash table in shared memory, mapping page identifiers → buffers
  - If the requested page is in `shared_buffers`, return it
  - Otherwise, ask the kernel for it and stash it in `shared_buffers`
    - If no free buffers, replace an existing one (which one?)
    - The kernel typically does its own I/O caching as well
- Keep a pin on the page, to ensure it isn’t replaced while in use
Concurrency Control

Table-level Locks

- Also known as “1mgr locks”, “heavyweight locks”
- Protect entire tables against concurrent DDL operations
- Many different lock modes; matrix for determining if two locks conflict
- Automatic deadlock detection and resolution
Concurrency Control

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Row-level Locks

- Writers don’t block readers: MVCC
- Writers must block writers: implemented via row-level locks
- Implemented by marking the row itself (on disk)
- Also used for SELECT FOR UPDATE, FOR SHARE
LWLocks ("Latches")

- Protect shared data structures against concurrent access
- Two lock modes: shared and exclusive (reader/writer)
- *No* deadlock detection: should only be held for short durations
## Concurrency Control: Low-Level Locks

### LWLocks ("Latches")

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- Two lock modes: shared and exclusive (reader/writer)
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### Spinlocks

- LWLocks are implemented on top of spinlocks, which are in turn a thin layer on top of an atomic test-and-set (TAS) primitive provided by the platform
- If an LWLock is contended, waiting is done via blocking on a SysV semaphore; spinlocks just busy wait, then micro-sleep
Organization of Source Tree

- **doc/**: documentation, FAQs
- **src/**
  - **bin/**: client programs (psql, pg_dump, ...)
  - **include/**: headers
    - **catalog/**: system catalog definitions
  - **interfaces/**: libpq, ecpg
  - **pl/**: procedural languages (PL/PgSQL, PL/Perl, ...)
  - **test/regress/**: SQL regression tests

Makefiles

- Makefile per directory (recursive make)
- **src/makefiles** has platform-specific Makefiles
- **src/Makefile.global.in** is the top-level Makefile
Content of src/backend

- **access/**: index implementations, heap access manager, transaction management, write-ahead log
- **commands/**: implementation of DDL commands
- **executor/**: executor logic, implementation of executor nodes
- **libpq/**: implementation of backend side of FE/BE protocol
- **optimizer/**: query planner
- **parser/**: lexer, parser, analysis phase
Content of src/backend, cont.

- **postmaster/**: postmaster, stats daemon, AV daemon, ...
- **rewrite/**: application of query rewrite rules
- **storage/**: shmem, locks, bufmgr, storage management, ...
- **tcop/**: “traffic cop”, FE/BE query loop, dispatching from protocol commands → implementation
- **utils/**:
  - **adt/**: builtin data types, functions, operators
  - **cache/**: caches for system catalog lookups, query plans
  - **hash/**: in-memory hash tables
  - **mmgr/**: memory management
  - **sort/**: external sorting, TupleStore
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Postgres uses a simple object system with support for single inheritance. The root of the class hierarchy is Node:

```c
typedef struct {
    NodeTag type;
    Parent parent;
} Node;

typedef struct {
    int a_field;
} Parent;

typedef struct {
    int b_field;
} Child;
```

This relies on a C trick: you can treat a Child * like a Parent * since their initial fields are the same.

Unfortunately, this can require a lot of ugly casting.

The first field of any Node is a NodeTag, which can be used to determine a Node’s specific type at runtime.
Nodes, Cont.

Basic Node Utility Functions

- Create a new Node: `makeNode()`
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- Serialise a node to text: `nodeToString()`
Basic Node Utility Functions

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- Test if two nodes are equal: `equal()`
- Deep-copy a node: `copyObject()`
- Serialise a node to text: `nodeToString()`
- Deserialise a node from text: `stringToNode()`
When you modify a node or add a new node, remember to update

- nodes/equalfuncs.c
- nodes/copyfuncs.c

You may have to update nodes/outfuncts.c and nodes/readfuncts.c if your Node is to be serialised/deserialised

Grep for references to the node’s type to make sure you don’t forget to update anything

- When adding a new node, look at how similar nodes are treated
Memory Management

- Postgres uses **hierarchical, region-based** memory management, and it absolutely rocks
  - `backend/util/mmgr`
  - Similar concept to Ttridge’s `talloc()`, “arenas”, ...
- All memory allocations are made in a **memory context**
- Default context of allocation: `CurrentMemoryContext`
- `palloc()` allocates in CMC
- `MemoryContextAlloc()` allocates in a given context
Allocations can be freed individually via `pfree()`.

When a memory context is reset or deleted, all allocations in the context are released.

- Resetting contexts is both faster and less error-prone than releasing individual allocations.

Contexts are arranged in a tree; deleting/resetting a context deletes/resets its child contexts.
You should sometimes `pfree()` your allocations

- If the context of allocation is known to be short-lived, don’t bother with `pfree()`
- If the code might be invoked in an arbitrary memory context (e.g. utility functions), you should `pfree()`
- You can’t `pfree()` an arbitrary Node (no “deep free”)

The exact rules are a bit hazy :-(
Memory Leaks

- Be aware of the memory allocation assumptions made by functions you call
- Memory leaks, *per se*, are rare in the backend
  - All memory is released eventually
  - A “leak” occurs when memory is allocated in a too-long-lived memory context: e.g. allocating some per-tuple resource in a per-txn context
    - `MemoryContextStats()` useful for locating the guilty context
- (Almost) never use `malloc()` in the backend
Error Handling

- Most errors reported by \texttt{ereport()} or \texttt{elog()}
  - \texttt{ereport()} is for user-visible errors, and allows more fields to be specified (SQLSTATE, detail, hint, etc.)
- Implemented via \texttt{longjmp}; conceptually similar to exceptions in other languages
  - \texttt{elog(ERROR)} walks back up the stack to the closest error handling block; that block can either handle the error or re-throw it
  - The top-level error handler aborts the current transaction and resets the transaction’s memory context
    - Releases all resources held by the transaction, including files, locks, memory, and buffer pins
Custom error handlers can be defined via PG_TRY().

Think about error handling!

- *Never* ignore the return values of system calls.

Should your function return an error code, or ereport() on failure?

- Probably ereport() to save callers the trouble of checking for failure.
- *Unless* the caller can provide a better (more descriptive) error message, or might not consider the failure to be an actual error.

Use assertions (Assert) liberally to detect programming mistakes, but *never* errors the user might encounter.
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Mailing Lists

- The vast majority of communication occurs on mailing lists
  - \texttt{pgsql-hackers} is the main list
  - \texttt{pgsql-patches} and \texttt{pgsql-committers} can be useful to learn from

- Written communication skills are important
  - Good developers are often good writers

- Some developers are on IRC; internals questions are welcome
  - \texttt{irc.freenode.net}, \#postgresql
Your First Patch

- **Step 1: Research and preparation**
  - Is your new feature actually useful? Does it just scratch your itch, or is it of general value?
  - Does it need to be implemented in the backend, or can it live in pgfoundry, contrib/, or elsewhere?
  - Does the SQL standard define similar or equivalent functionality?
    - What about Oracle, DB2, ...?
  - Has someone suggested this idea in the past?
    - Search the archives and TODO list
  - Most ideas are bad
  - **Don’t** take the TODO list as gospel
Sending A Proposal

- **Step 2**: Send a proposal for your feature to `pgsql-hackers`
  - Patches that appear without prior discussion risk wasting your time
- Discuss your proposed syntax and behaviour
  - Consider corner cases, and how the feature will relate to other parts of PostgreSQL (consistency is good)
  - Will any system catalog changes be required?
  - Backward-compatibility?
- Try to reach a consensus with `pgsql-hackers` on how the feature ought to behave
Step 3: Begin implementing the feature

A general strategy is to look at how similar parts of the system function

- Don’t copy and paste (IMHO)
  - Common source of errors
- Instead, read through similar sections of code to try to understand how they work, and the APIs they are using
- Implement (just) what you need, refactoring the existed APIs if required

Ask for advice as necessary (−hackers or IRC)

- Write down the issues you encounter as you write the code, include the list when you submit the patch

Consider posting work-in-progress versions of the patch
Testing, Documentation

- **Step 4: Update tools**
  - For example, if you’ve modified DDL syntax, update psql’s tab completion
  - Add pg_dump support if necessary

- **Step 5: Testing**
  - Make sure the existing regression tests don’t fail
  - No compiler warnings
  - Add new regression tests for the new feature

- **Step 6: Update documentation**
  - Writing good documentation is more important than getting the DocBook details completely correct
  - Add new index entries, if appropriate
  - Check documentation changes visually in a browser
Submitting The Patch

- **Step 7: Submit the patch**
  - Use context diff format: `diff -c`
    - Unified diffs are okay for SGML changes
  - First, review every hunk of the patch
    - Is this hunk necessary?
    - Does it needlessly change existing code or whitespace?
    - Does it have any errors? Does it fail in corner cases? Is there a more elegant way to do this?
  - Work with a code reviewer to make any necessary changes
  - If your patch falls through the cracks, *be persistent*
    - The developers are busy and reviewing patches is difficult, time-consuming, and unglamourous work
The TABLESAMPLE Clause

- The TABLESAMPLE clause is defined by SQL:2003 and implemented by SQL Server and DB2
  - Oracle calls it SAMPLE, slightly different syntax
- Example query:
  
  ```sql
  SELECT avg(salary)
  FROM emp TABLESAMPLE SYSTEM (50);
  ```
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  - Oracle calls it SAMPLE, slightly different syntax
- Example query:
  ```sql
  SELECT avg(salary)
  FROM emp TABLESAMPLE SYSTEM (50);
  ```
- TODO item: “estimated_count(*)”
  ```sql
  SELECT count(*) * 10
  FROM t TABLESAMPLE SYSTEM (10);
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  ```
- Straightforward to implement, but requires modifying some interesting parts of the system
Deciphering the SQL standard is notoriously difficult
- I usually start with the index
- The BERNULLI sample method sounds hard to implement
- REPEATABLE provides a way to seed the random number generator
How Should We Implement Sampling?

- Simple approach: sequentially walk the heap, decide whether to skip a block using `random()` and the sampling percentage.
- Therefore, add “sample scan” as a new scan type, analogous to sequential scan or index scan.
Implementation Ideas

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Deficiencies

1. Non-uniform sampling when either
   - row size is non-uniform
   - distribution of live tuples is non-uniform
2. Consumes a lot of entropy
3. Could be optimized to reduce random I/O
Behavioral Questions

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4. How does this interact with inheritance? Joins?
Implementation Plan

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2. Modify the nodes of the parse tree to allow TABLESAMPLE to be encoded in the AST
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5. Implement the guts of the SampleScan executor node
6. Add support for REPEATABLE
7. Add support for DELETE and UPDATE
8. Update documentation
   - Can’t easily add regression tests
Grammar Modifications

- Parsing TABLESAMPLE itself is quite easy
  - Add some new keywords: TABLESAMPLE and REPEATABLE must be made semi-reserved to avoid shift-reduce conflicts

- Checking SelectStmt reveals that relation_expr is the production for a base relation in the FROM clause with an optional alias and inheritance spec

- Unfortunately, relation_expr is also used by DDL commands, so create a new production and use it in the places we want to allow TABLESAMPLE
New parse node for the data TABLESAMPLE clause

Need to attach new parse node to the AST somehow

- The parser constructs a RangeVar for each FROM clause entry, so use that
Parse Node Updates

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Range Table

The parse-analysis phase constructs a “range table” consisting of the FROM clause elements

- When converting the FROM clause RVs into range table entries (RTEs), attach the TableSampleInfo
Optimiser Terminology

**RelOptInfo**: Per-relation planner state. For each base relation or join, stores the estimated row count, row width, cheapest path, ... 

**Path**: Planner state for a particular way accessing a relation (or join relation); each RelOptInfo has a list of candidate paths
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**Plan**: A “finalized” output path: a node of the plan tree passed to the executor

- Once the planner has decided on the optimal Path tree, produce a corresponding Plan tree
Optimizer Modifications

- We need only modify stage 1 of the System R algorithm: finding the cheapest interesting paths for each base relation
  - Joins between sample scans not fundamentally different than normal joins
  - We don’t need a SamplePath node; just use Path

- Only consider sample scans when a TABLESAMPLE clause is specified

- Simple cost estimation: assume we need to do a single I/O for each sampled page
Plan Trees

- Review: the planner produces a tree of Plan nodes
  - Plan nodes are treated as immutable by the executor
- The executor constructs a tree of PlanState nodes to describe the run-time state of a plan-in-execution
  - Each PlanState is associated with exactly one Plan node
  - PlanState.plan holds a PlanState’s associated Plan node
# The “Iterator” API
Implemented By Each Executor Node

<table>
<thead>
<tr>
<th>Mandatory</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>InitNode</strong>: Given a Plan tree, construct a PlanState tree</td>
</tr>
<tr>
<td><strong>ProcNode</strong>: Given a PlanState tree, return next result tuple</td>
</tr>
<tr>
<td>■ Some plan nodes support bidirectional scans</td>
</tr>
<tr>
<td><strong>EndNode</strong>: Shutdown a PlanState tree, releasing resources</td>
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<tr>
<td><strong>ReScan:</strong></td>
</tr>
<tr>
<td><strong>MarkPos:</strong></td>
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<tr>
<td><strong>RestrPos:</strong></td>
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Executor Terminology

**Block**: A page on disk. Identified by a `BlockNumber`

**Buffer**: A page in memory. The buffer manager loads blocks from disk into buffers (shared buffers)

**OffsetNumber**: Identifies an item within a page

**Datum**: An instance of a data type in memory

**HeapTuple**: A collection of Datums with a certain schema

**EState**: Run-time state for a single instance of the executor

**Projection**: The act of applying a target list
The Executor’s TupleTable

- Tuples are passed around the executor using TupleTableSlots
- Different kinds of tuples:
  - Pointers into buffer pages
    - The output of a scan node, no projection
    - Need to drop pin on buffer when finished with tuple
  - Pointers into heap-allocated memory
    - Result of applying an expression: projection, SRFs, ...
    - Can be “minimal” tuples: no MVCC metadata needed
    - Need to pfree() tuple when finished
  - “Virtual” tuples
- The TupleTableSlot abstraction papers over all these details
Implementing The Executor Node

Initialization

- Most of this is boilerplate code :-(
- Initialize executor machinery needed to evaluate quals and do projection
- Read-lock the relation: no DDL changes allowed while we’re scanning
Implementing REPEATABLE

- Simple implementation: pass the repeat seed to `srandom()`
Implementing **REPEATABLE**

- Simple implementation: pass the repeat seed to `srandom()`
- **Wrong**: if the execution of multiple sample scans is interleaved, they will stomp on the other’s PRNG state
- Therefore, use `initstate()` to give each sample scan its own private PRNG state
Supporting UPDATE and DELETE

Implementation of UPDATE and DELETE

- Run the executor to get “result tuples”
- Mark the result tuples as expired (“deleted by my transaction”) on disk
- If UPDATE, insert a new tuple
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TABLESAMPLE support

- Quite easy: basically comes for free!
- relation_expr is already used by the DELETE and UPDATE
  - Modify to use relation_expr_opt_sample
- Hackup parse-analysis to attach TableSampleInfo
Possible Improvements

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5. Provide information about the degree of confidence in the sampled results
6. “Page at a time” scan mode
Next Steps

1. Sign up to the development lists
2. Setup your local development environment
3. Participate in development discussions
   - Read design proposals, ask questions/give feedback
   - Try to reproduce (and fix!) reported bugs
   - Look at proposed patches
   - Help out with administrativia, contribute to the documentation
4. Read the code!
5. Look for a small project that piques your interest, and get started!
Any questions?