
CS206 Quiz 2

Max marks: 55

Time: 1.5 hours

- *Be brief, complete and stick to what has been asked.*
- *If needed, you may cite results/proofs covered in class without reproducing them.*
- *If you need to make any assumptions, state them clearly.*
- *Do not copy solutions from others.*

1. Let ϕ_1 be $\forall x \neg S(x, x)$, ϕ_2 be $\exists x P(x)$, ϕ_3 be $\forall x \exists y S(x, y)$, ϕ_4 be $\forall x (P(x) \rightarrow \exists y S(y, x))$.

- (a) [5 marks] Give a model M such that $M \models \phi_1 \wedge \phi_2 \wedge \phi_3 \wedge \neg \phi_4$.
(b) [5 marks] Give a model M' such that $M' \models \phi_1 \wedge \phi_2 \wedge \phi_3 \wedge \phi_4$.

Domain in your model must not have more than three elements for each subpart.

2. Prove using natural deduction.

- (a) [10 marks] $\forall x \exists y (P(x) \rightarrow Q(y)) \vdash \exists y \forall x (P(x) \rightarrow Q(y))$
(b) [10 marks] $\exists x \exists y (H(x, y) \vee H(y, x)), \neg \exists x H(x, x) \vdash \exists x \exists y \neg (x = y)$

3. Let $\phi(x, y, z)$ and $\psi(x, y, z)$ be predicate logic formulae with free variables x, y, z and containing the binary function symbol $+$ and $*$ and the binary predicate symbols $=$ and $>$. Let M be a model in which the domain is a set of integers, the predicates $=$ and $>$ are interpreted as the usual “equal to” and “greater than” predicates over pairs of integers and functions $+$ and $*$ are interpreted as usual “addition” and “multiplication” functions over pairs of integers.

Now consider the following program fragment in C or C++ or Java, where x, y and z are variables of type integer.

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x = x + y + (z*x*y);
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Thus if the values of x, y and z prior to the execution of this statement were 1, 2 and 3 respectively, their values after execution of this statement would be 9, 2 and 3 respectively.

- (a) [10 marks] You are told that $\phi(x, y, z)$ evaluates to True in the model M described above if and only if the values of x, y and z **prior to** the execution of program statement are used for the free variables x, y and z in $\phi(x, y, z)$. Give a predicate logic formula $\psi(x, y, z)$ that evaluates to True in the same model M described above if and only if the values of x, y and z **after** the execution of program statement are used for the free variables x, y and z in $\psi(x, y, z)$. Your formula $\psi(x, y, z)$ must necessarily depend on $\phi(x, y, z)$ since the values of x, y and z prior to the execution of program statement determine the values of the variables after the execution of program statement.

(b) [10 marks] Now consider the inverse problem. You are told that $\psi'(x, y, z)$ evaluates to True in the model M described above if and only if the values of x, y and z **after** the execution of program statement are used for the free variables x, y and z in $\psi'(x, y, z)$. Give a predicate logic formula $\phi'(x, y, z)$ that evaluates to True in the same model M described above if and only if the values of x, y and z **prior to** the execution of program statement are used for the free variables x, y and z in $\phi'(x, y, z)$.

Your formula $\phi'(x, y, z)$ must necessarily depend on $\psi'(x, y, z)$ for the same reason as in the previous subquestion.

(c) [5 marks] Starting from $\phi(x, y, z)$, suppose we derive $\psi(x, y, z)$ as in subpart (a). Now letting $\psi(x, y, z) \equiv \psi'(x, y, z)$, suppose we derive $\phi'(x, y, z)$ as in subpart (b). Indicate which of the following relations hold and briefly explain why.

- $\phi(x, y, z) \equiv \phi'(x, y, z)$
- $\phi(x, y, z) \rightarrow \phi'(x, y, z)$
- $\phi'(x, y, z) \rightarrow \phi(x, y, z)$
- None of the above.