B^{ed}-Tree: An All-Purpose Index Structure for String Similarity Search Based on Edit Distance

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Outline



Motivation and B^{ed}-Tree Framework

- String Orders
 - Dictionary order
 - Gram counting order
 - Gram location order
- Experiments
- Conclusion

Approximate String Search



Information Retrieval

Web search query with string "Posgre SQL" instead of "Postgre SQL"

Data Cleaning

"13 Computing Road" is the same as "#13 Computing Rd"?

Bioinformatics

Find out all protein sequences similar to "ACBCEEACCDECAAB"

Edit Distance





Existing Solution

Observation: If ED(s_1, s_2)=d, they agree on at least min($|s_1|, |s_2|$)+Q-1-d*(Q+1) grams

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Existing Solution

Inverted List

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Inverted List Method

Limited queries supported

	Range Query	Join Query	Top-K Query	Top-K Join
Edit Distance	Y	Y	Ν	Ν
Normalized ED	Ν	Ν	Ν	Ν

Uncontrollable memory consumptionConcurrency protocol

Our Contributions

■ B^{ed}-Tree

□Wide support on different queries and distances

	Range Query	Join Query	Top-K Query	Top-K Join
Edit Distance	Y	Y	Y	Y
Normalized ED	Y	Y	Y	Y

□Adjustable buffer size and low I/O cost

- Highly concurrent
- Easy to implement
- Competitive performance

Basic Index Framework

Bed-Tree Framework

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String Order Properties

P1: Comparability

Given two string s1 and s2, we know the order of s1 and s2 under the specified string order

P2: Lower Bounding

Given an interval [L,U] on the string order, we know a lower bound on edit distance to the query string

String Order Properties

■ P3: Pairwise Lower Bounding

□ Given two intervals [L,U] and [L',U'], we know the lower bound of edit distance between s1 from [L,U] and s2 from [L',U']

P4: Length Bounding

Given an interval [L,U] on the string order, we know the minimal length of the strings in the interval

String Order Properties

Properties v.s. supported queries and distances

	Range Query	Join Query	Top-K Query	Top-K Join
Edit Distance	P1, P2	P1, P3	P1, P2	P1, P3
Normalized ED	P1, P2, P4	P1, P3, P4	P1, P2, P4	P1, P3, P4

	Description	
P1	Comparability	
P2	Lower Bounding	
P3	Pair-wise Lower Bounding	
P4	Length Bounding	

All strings are ordered alphabetically, satisfying P1, P2 and P3

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Gram Counting Order

Jim Gray

Gram Counting Order

Transform the count vector to a bit string with z-order

Order the strings with this signature

Gram Counting Order

Extension of Gram Counting Order

Include positional information of the grams

Allow better estimation of mismatch gramsHarder to encode

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Experiment Settings

Dataset	# of Strings	Max. Length	Avg. Length
Author	2,948,929	56	22
Title	$1,\!158,\!648$	675	74
Actor	1,213,391	80	23
Movie	1,568,893	247	26
Protein	$508,\!038$	$1,\!999$	347

Five Index Schemes

Bed-Tree: BD, BGC, BGL

Inverted List: Flamingo, Mismatch

Default Setting

□Q=2, Bucket=4, Page Size=4KB

Data

Empirical Observations

■ How good is B^{ed}-Tree?

- □With small threshold, Inverted Lists are better
- □When threshold increases, B^{ed}-Tree is not worse

Empirical Observations

Which string order is better?

- Gram counting order is generally better
- Gram Location order: tradeoff between gram content information and position information

- A new B+ tree index scheme
 - □All similarity queries supported
 - □Both edit distance and normalized distance
 - General transaction and concurrency protocol
 - **Competitive efficiencies**

Results

Range Query

■Top-K Query

Normalized Edit Distance & Join Query

