CS 344 Artificial Intelligence By Prof: Pushpak Bhattacharya Class on 02/Feb/2007

Formalization of propositional logic (review) <u>Axioms :</u>

$$(A \to (B \to A)) \qquad A1$$

$$((A \to (B \to C)) \to ((A \to B) \to (A \to C))) \qquad A2$$

$$(((A \to F) \to F) \to A) \qquad A3$$

<u>Inference rule:</u> Given $(A \rightarrow B)$ and *A*, write *B*

<u>A Proof is:</u> A sequence of i) Hypotheses ii) Axioms iii) Results of MP <u>A Theorem is an</u> Expression proved from axioms and inference rules Example: To prove $(P \rightarrow P)$ A1 : P for A and Bi) $P \rightarrow (P \rightarrow P)$ A1 : P for A and $(P \rightarrow P)$ for Bii) $P \rightarrow ((P \rightarrow P) \rightarrow P)$ A1: P for A and $(P \rightarrow P)$ for Biii) $[(P \rightarrow ((P \rightarrow P) \rightarrow P)) \rightarrow ((P \rightarrow (P \rightarrow P))) \rightarrow ((P \rightarrow P)))]$ A2: with P for $A, (P \rightarrow P)$ for B and P for Civ) $(P \rightarrow (P \rightarrow P) \rightarrow (P \rightarrow P))$ MP, (ii), (iii)v) $(P \rightarrow P)$ MP, (i), (iv)

Shorthand

- 1. $\neg P$ is written as $P \rightarrow F$ and called '*NOT P*'
- 2. $((P \rightarrow F) \rightarrow Q)$ is written as $(P \lor Q)$ and called

'P OR Q'

- 3. $((P \rightarrow (Q \rightarrow F)) \rightarrow F)$ is written as $(P \land Q)$ and called 'P AND Q'
- Exercise: (Challenge)
 - Prove that $A \rightarrow \neg(\neg(A))$

A very useful theorem (Actually a meta theorem, called deduction theorem)

Statement If $A_1, A_2, A_3, \dots, A_n \models B$ then $A_1, A_2, A_3, \dots, A_{n-1} \models A_n \rightarrow B$ \models is read as 'derives'

Given



Use of Deduction Theorem Prove $A \rightarrow \neg(\neg(A))$

i.e., $A \to ((A \to F) \to F)$

- $A, A \to F \models F \tag{M.P}$
- $A \models (A \to F) \to F \tag{D.T}$
 - $\vdash A \to ((A \to F) \to F) \tag{D.T}$

Very difficult to prove from first principles, *i.e.*, using axioms and inference rules only

Prove $P \rightarrow (P \lor Q)$ i.e. $P \rightarrow ((P \rightarrow F) \rightarrow Q)$ $P, P \rightarrow F, Q \rightarrow F \vdash F$ $P, P \to F \vdash (Q \to F) \to F \text{ (D.T)}$ (M.P with A3) $\vdash Q$ $\mathbf{P} \vdash (P \to F) \to Q$ $\vdash P \to ((P \to F) \to Q)$