

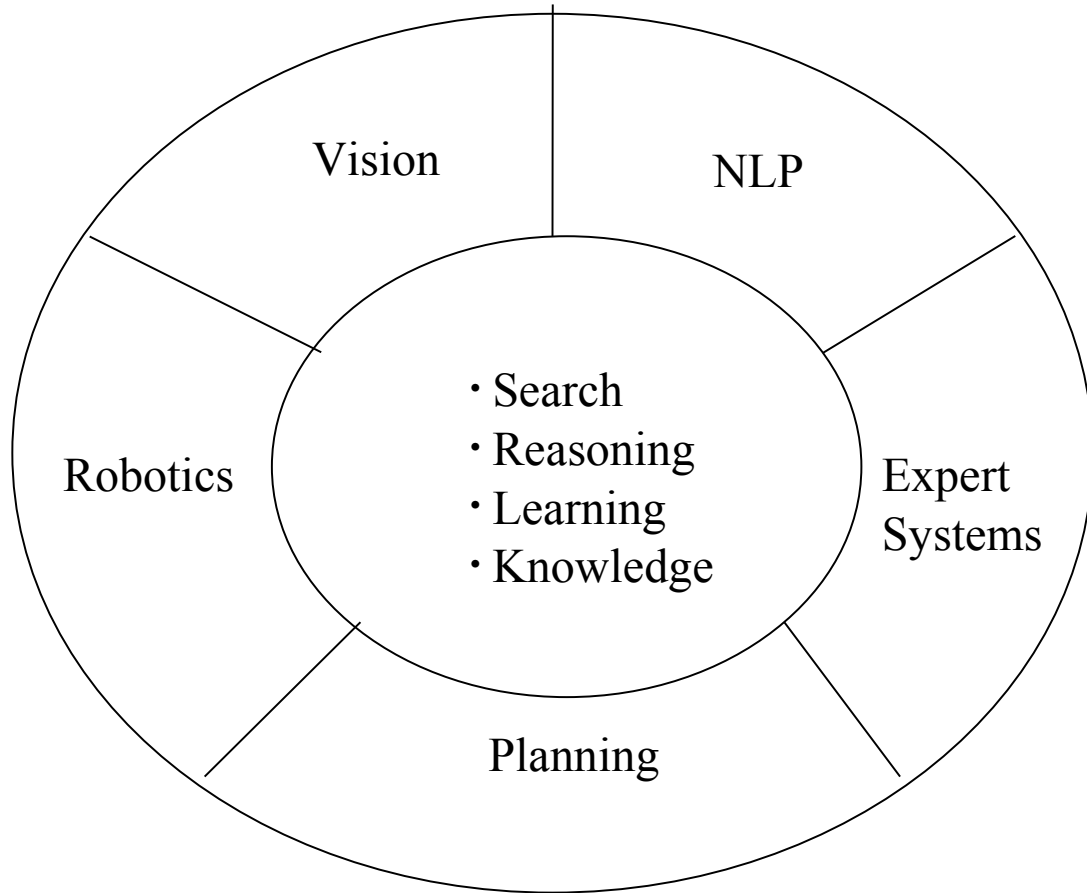
CS 344

Artificial Intelligence

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# Logic and inferencing



Obtaining implication of given facts and rules -- Hallmark of intelligence

# Inferencing through

- Deduction (General to specific)
- Induction (Specific to General)
- Abduction (Conclusion to hypothesis in absence of any other evidence to contrary)

## Deduction

Given:           All men are mortal (rule)  
                  Shakespeare is a man (fact)

To prove:       Shakespeare is mortal (inference)

## Induction

Given:           Shakespeare is mortal  
                  Newton is mortal                   (Observation)  
                  Dijkstra is mortal

To prove:       All men are mortal (Generalization)

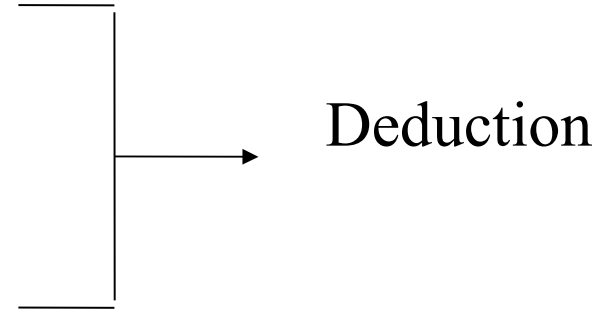
If there is rain, then there will be no picnic

Fact1: There was rain

Conclude: There was no picnic

Fact2: There was no picnic

Conclude: There was no rain (?)



Induction and abduction are fallible forms of reasoning. Their conclusions are susceptible to retraction

Two systems of logic

- 1) Propositional calculus
- 2) Predicate calculus

## Propositions

- Stand for facts/assertions
- Declarative statements
  - As opposed to interrogative statements (questions) or imperative statements (request, order)

## Operators

*AND* ( $\wedge$ ), *OR* ( $\vee$ ), *NOT* ( $\neg$ ), *IMPLICATION* ( $\Rightarrow$ )

$\Rightarrow$  and  $\neg$  form a minimal set (can express other operations)  
- Prove it.

Tautologies are formulae whose truth value is always T, whatever the assignment is

## Model

In propositional calculus any formula with  $n$  propositions has  $2^n$  models (assignments)

- Tautologies evaluate to T in all models.

Examples:

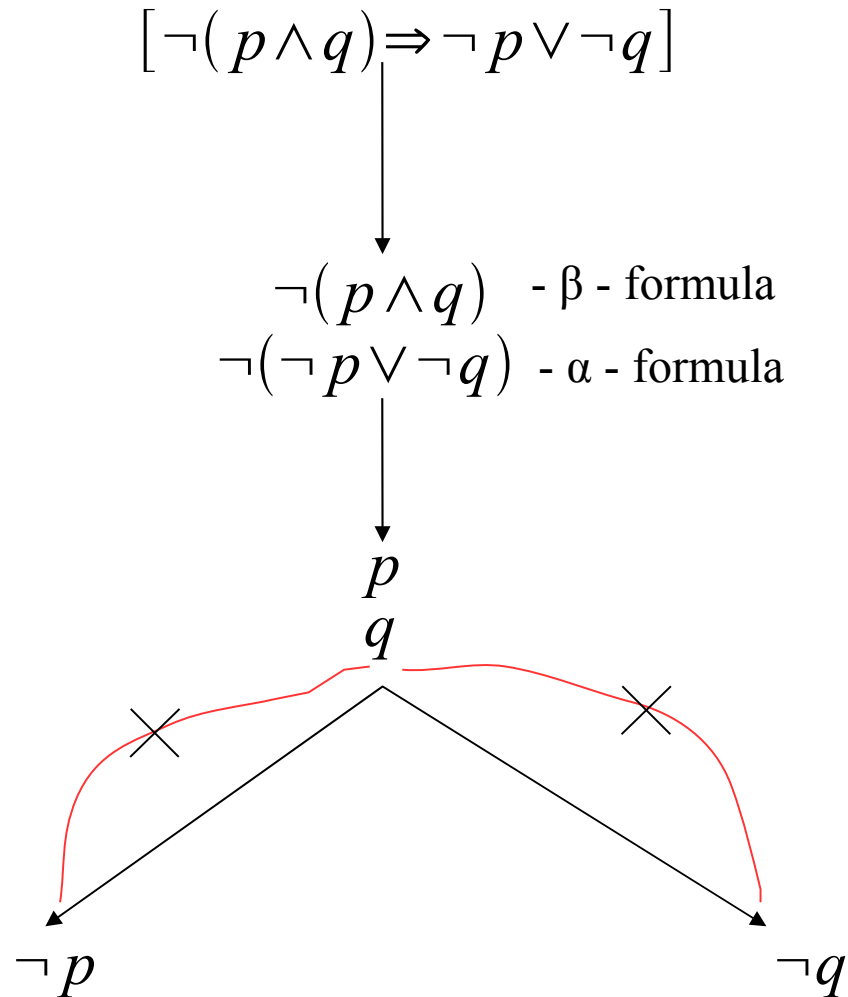
1)  $p \wedge \neg p$

2)  $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$  - De Morgan with AND

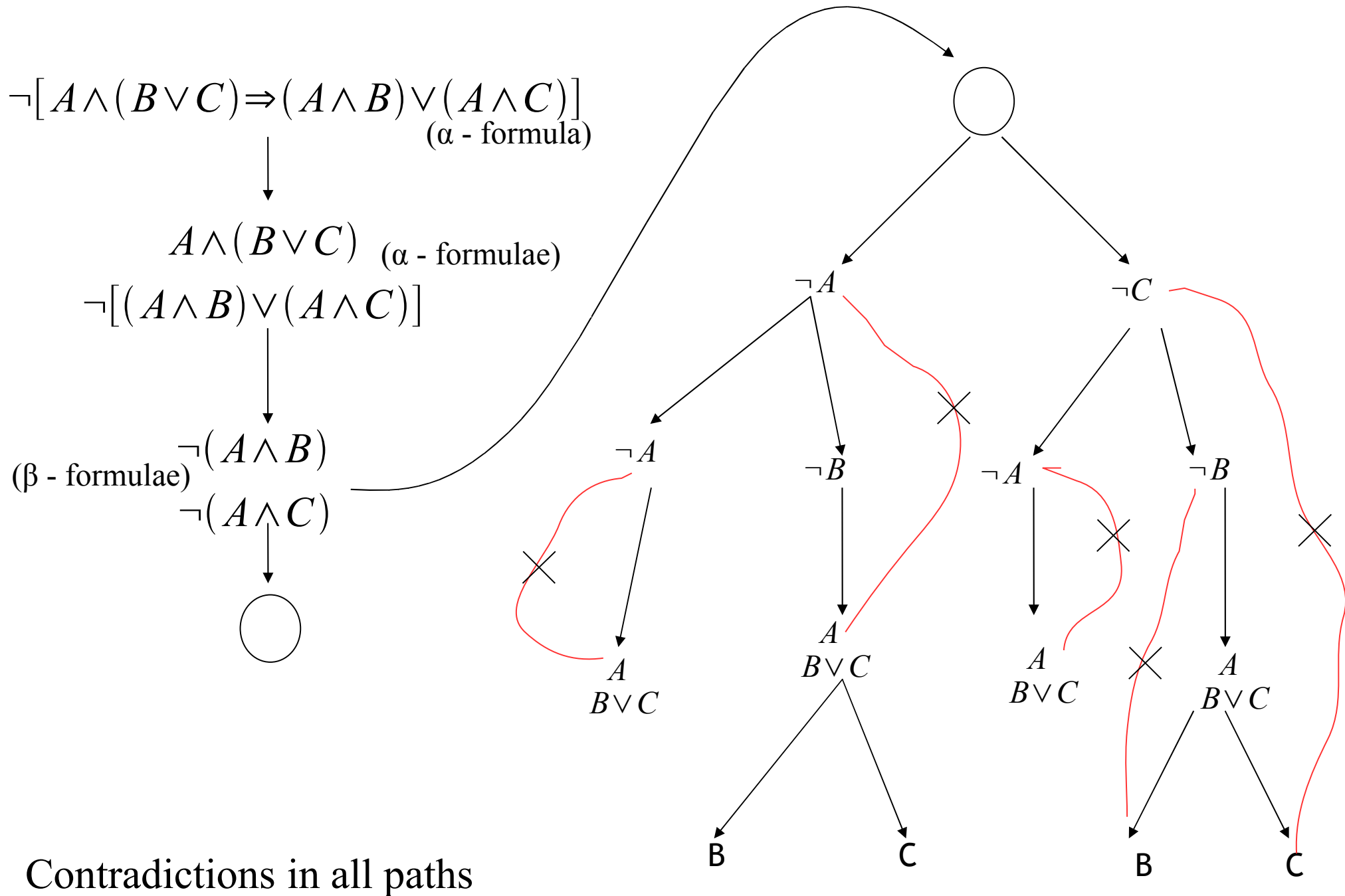
# Semantic Tree/Tableau method of proving tautology

To prove:  $\neg(p \wedge q) \Rightarrow \neg p \vee \neg q$

Start with the negation of the formula to be proved a tautology



Example 2:  $A \wedge (B \vee C) \Rightarrow (A \wedge B) \vee (A \wedge C)$





Exercise:

Prove the backward implication in the previous example