CS 747, Autumn 2022: Lecture 11

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Autumn 2022

Markov Decision Problems

1. Policy iteration: variants and complexity bounds

- 2. Analysis of bounds
 - Basic tools
 - Howard's PI with k = 2
 - BSPI with k = 2
 - Open problems
- 3. Review of MDP planning

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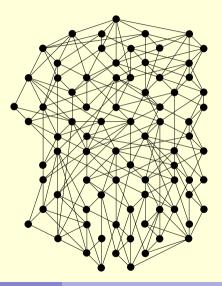
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\pi \leftarrow Arbitrary policy.

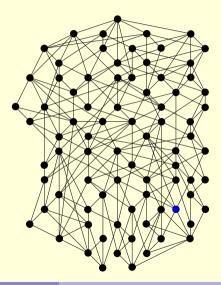
While \pi has improvable states:

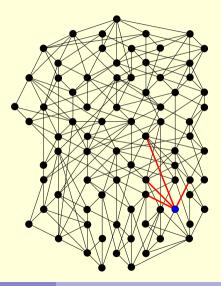
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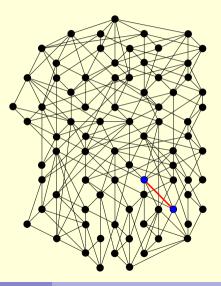
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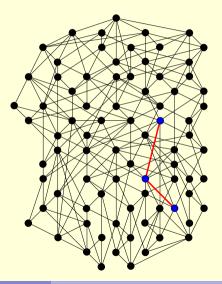
Return \pi.
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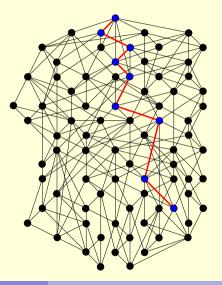






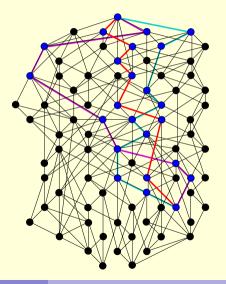






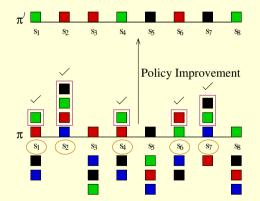
 $\pi \leftarrow$ Arbitrary policy. **While** π has improvable states: $\pi' \leftarrow$ PolicyImprovement(π). $\pi \leftarrow \pi'$. **Return** π .

Path taken (and hence the number of iterations) in general depends on the switching strategy.



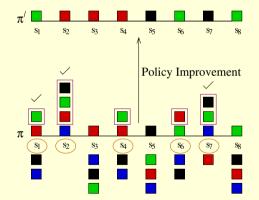
Howard's Policy Iteration

- Reference: Howard (1960).
- Greedy; switch all improvable states.



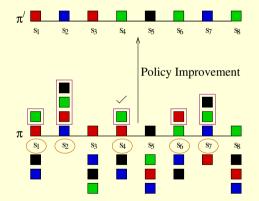
Random Policy Iteration

- Reference: Mansour and Singh (1999).
- Switch a non-empty subset of improvable states chosen uniformly at random.



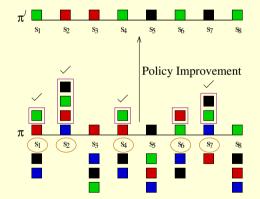
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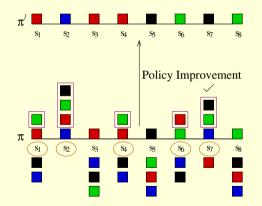
Random Policy Iteration

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Simple Policy Iteration

- Reference: Melekopoglou and Condon (1994).
- Assume a fixed indexing of states.
- Switch the improvable state with the highest index.



Upper and Lower Bounds

U(n, k) is an upper bound applicable to a set of PI variants \mathcal{L} if

- for each *n*-state, *k*-action MDP $M = (S, A, T, R, \gamma)$,
- for each policy $\pi: S \to A$,
- for each algorithm $L \in \mathcal{L}$,

the expected number of policy evaluations performed by *L* on *M* if initialised at π is at most U(n, k).

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the expected number of policy evaluations performed by *L* on *M* if initialised at π is at most U(n, k). L(n, k) is a lower bound applicable to a set of PI variants \mathcal{L} if

- there exists an *n*-state, *k*-action
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- there exists a policy $\pi : S \rightarrow A$,
- there exists an algorithm $L \in \mathcal{L}$,

such that the expected number of policy evaluations performed by *L* on *M* if initialised at π is at least L(n, k).

Switching Strategies and Bounds

Upper bounds on number of iterations

PI Variant	Туре	<i>k</i> = 2	General k
Howard's (Greedy) PI [H60, MS99]	Deterministic	$O\left(\frac{2^n}{n}\right)$	$O\left(\frac{k^n}{n}\right)$
Mansour and Singh's Random PI [MS99]	Randomised	1.7172 ⁿ	$pprox O\left(rac{k}{2} ight)^n$
Mansour and Singh's Random PI [HPZ14]	Randomised	$poly(n) \cdot 1.5^n$	_

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Lower bounds on number of iterations

Ω(n) Howard's PI on *n*-state, 2-action MDPs [HZ10]. $Ω(2^n)$ Simple PI on *n*-state, 2-action MDPs [MC94].

(Polynomial factors ignored. Authors with names underlined once took CS 747!)

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- <u>Ashutosh</u>, <u>Consul</u>, <u>Dedhia</u>, <u>Khirwadkar</u>, <u>Shah</u>, and Kalyanakrishnan (2020) show a *lower bound* of \sqrt{k}^n iterations for a deterministic variant of PI.

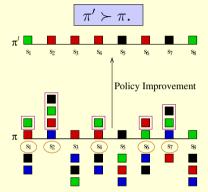
Markov Decision Problems

1. Policy iteration: variants and complexity bounds

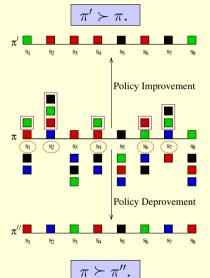
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1. Policy Improvement and Policy "Deprovement"



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Non-optimal policies $\pi, \pi' \in \Pi$ cannot have the same set of improvable states.

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2. Improvement sets in 2-action MDPs

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Switch actions in every improvable state.

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π 0 0 0 0 0 0 0 0 0 0 0 0 0

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Switch actions in every improvable state.

π′ 0 0 0 0 0 0 0 0 1 1 1 1 1

π 0 0 0 0 0 0 0 0 0 0 0 0

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Switch actions in every improvable state.

 Possible?

 π' 0
 0
 0
 0
 1
 1
 1
 1

π 0 0 0 0 0 0 0 0 0 0 0 0

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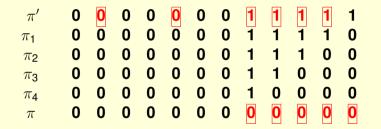
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If π has *m* improvable states and $\pi \xrightarrow{\text{Howard's PI}} \pi'$, then there exist *m* policies π'' such that $\pi' \succeq \pi'' \succ \pi$.

• Take $m^* = \frac{n}{3}$.

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$$\leq \binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \cdots + \binom{n}{m^{\star} - 1}$$

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Number of iterations taken by Howard's PI: $O\left(\frac{2^n}{n}\right)$ [MS99, HGDJ14].

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Batch-Switching Policy Iteration (BSPI) (2-action MDPs)

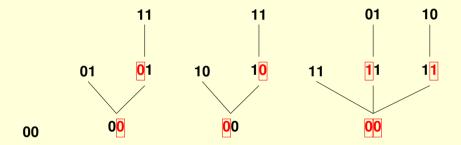
Howard's Policy Iteration takes at most _____ iterations on a 2-state MDP!

Batch-Switching Policy Iteration (BSPI) (2-action MDPs)

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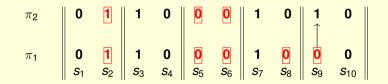
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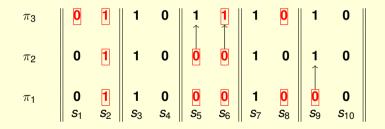
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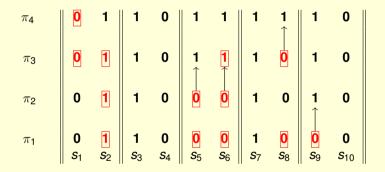
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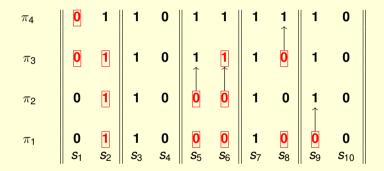
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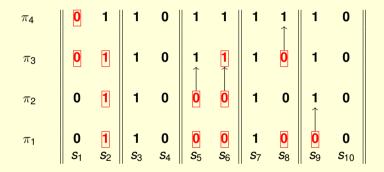


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• Left-most batch can change only when all other columns are non-improvable.

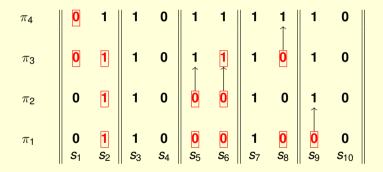
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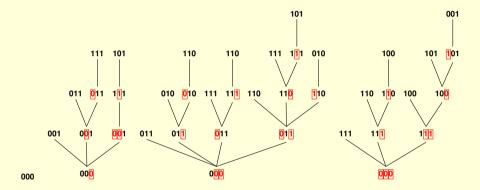
•
$$T(n) \leq 3 \times T(n-2) \leq v$$

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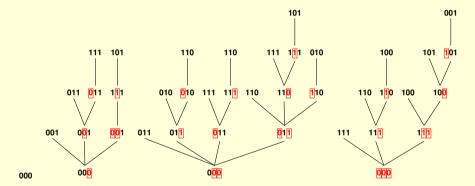
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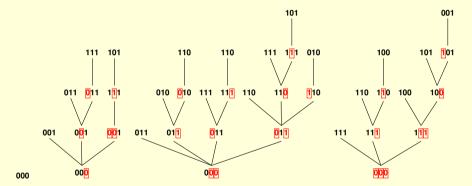


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The structures above are called Trajectory-bounding Trees (TBTs) [KMG16a] (and correspond to the Order Regularity Problem [H12, GHDJ15]).

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The structures above are called Trajectory-bounding Trees (TBTs) [KMG16a] (and correspond to the Order Regularity Problem [H12, GHDJ15]). BSPI with 3-sized batches gives $T(n) \le 5 \times T(n-3) \le 1.71^n$.

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Batch size	Depth of TBT	Bound on number of iterations
1	2	2 ⁿ
2	3	1.7321 ^{<i>n</i>}
3	5	1.7100 ⁿ
4	8	1.6818 ⁿ
5	13	1.6703 ^{<i>n</i>}
6	21	1.6611 ^{<i>n</i>}
7	33	1.6479 ^{<i>n</i>}

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Depth of TBT for batch size 7 due to Gerencsér et al. [GHDJ15].

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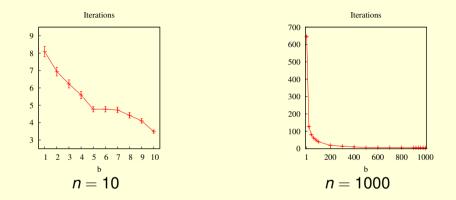
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Depth of TBT for batch size 7 due to Gerencsér *et al.* [GHDJ15]. Will the bound continue to be non-increasing in the batch size? If so, 1.6479^{*n*} would be a bound for Howard's Policy Iteration!

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BSPI: Effect of Batch Size b



Averaged over *n*-state, 2-action MDPs with randomly generated transition and reward functions. Each point is an average over 100 randomly-generated MDP instances and initial policies [KMG16a].

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Open Problems

- Is the complexity of Howard's PI on 2-action MDPs upper-bounded by the Fibonacci sequence (≈ 1.6181ⁿ)?
- Is Howard's PI the most efficient among deterministic PI algorithms (worst case over all MDPs)?
- Is there a super-linear lower bound on the number of iterations taken by Howard's PI on 2-action MDPs?
- Is Howard's PI strongly polynomial on deterministic MDPs?
- Is there a variant of PI that can visit all kⁿ policies in some n-state, k-action MDP—implying an Ω(kⁿ) lower bound?
- Is there a strongly polynomial algorithm for MDP planning?

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Summary of MDP Planning

- MDPs are an abstraction of sequential decision making.
- Many applications; many different formulations.
- Essential solution concept: optimal policy (known to exist).
- Three main families of planning algorithms: value iteration, linear programming, policy iteration.
- Have strengths and weaknesses in theory and in practice. Can combine.
- We showed correctness of all three methods.
- Used Banach's fixed-point theorem, Bellman (optimality) operator.
- What if *T*, *R* were not given, but have to be *learned* from interaction? Can we still learn to act optimally?
- Yes: that's the reinforcement learning problem. Next week!