CS 747, Autumn 2023: Lecture 5

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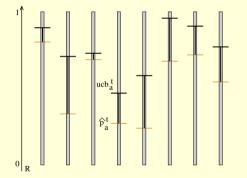
Multi-armed Bandits

- 1. Analysis of UCB
- 2. Other bandit problems

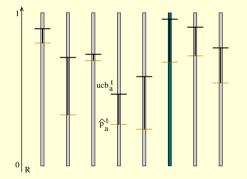
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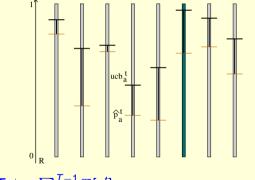
- Pull each arm once.
- For $t \in \{n, n+1, ...\}$, for $a \in A$, $\operatorname{ucb}_a^t \stackrel{\text{def}}{=} \hat{p}_a^t + \sqrt{\frac{2\ln(t)}{u_a^t}}$; pull $\operatorname{argmax}_{a \in A} \operatorname{ucb}_a^t$.



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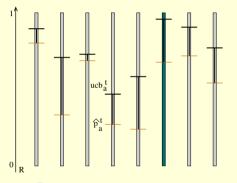


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- Recall that $R_T = Tp^* \sum_{t=0}^{T-1} \mathbb{E}[r^t]$.
- We shall show that UCB achieves $R_T = O\left(\sum_{a:p_a \neq p^*} \frac{1}{p^* p_a} \log(T)\right)$.

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Observe that $\mathbb{E}[z_a^t] = \mathbb{P}\{Z_a^t\}(1) + (1 - \mathbb{P}\{Z_a^t\})(0) = \mathbb{P}\{Z_a^t\}.$

• As in the algorithm, u_a^t is a random variable that denotes the number of pulls arm *a* has received up to time *t*:

$$u_a^t = \sum_{i=0}^{t-1} z_a^i$$

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 - Observe that $\mathbb{E}[z_a^t] = \mathbb{P}\{Z_a^t\}(1) + (1 \mathbb{P}\{Z_a^t\})(0) = \mathbb{P}\{Z_a^t\}.$
- As in the algorithm, u_a^t is a random variable that denotes the number of pulls arm *a* has received up to time *t*:

$$u_a^t = \sum_{i=0}^{t-1} z_a^i$$

• We define an instance-specific constant $\bar{u}_a^T \stackrel{\text{def}}{=} \left\lceil \frac{8}{(\Delta_a)^2} \ln(T) \right\rceil$ that will serve in our proof as a "sufficient" number of pulls of arm *a* for horizon *T*.

Proof Sketch

- To upper-bound R_T , upper-bound the number of pulls of each sub-optimal arm *a*.
- Give each such arm $a \bar{u}_a^T$ pulls for free.
- Beyond \bar{u}_a^T pulls, arm *a*'s UCB will have width at most $\Delta_a/2$.
- If *a* continues to be pulled beyond \bar{u}_a^T pulls, either its empirical mean has deviated by more than $\Delta_a/2$ from its true mean, or \star 's UCB has fallen below its true mean.
- Both events above have a low probability—in aggregate at most a constant even if summed over an infinite horizon.

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- Both events above have a low probability—in aggregate at most a constant even if summed over an infinite horizon.
- KL-UCB uses the KL inequality, and slightly more sophisticated analysis.

$$R_{T} = T p^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[r^{t}]$$

$$R_{T} = Tp^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[r^{t}] = Tp^{\star} - \sum_{t=0}^{T-1} \sum_{a \in A} \mathbb{P}\{Z_{a}^{t}\}\mathbb{E}[r^{t}|Z_{a}^{t}]$$

$$R_{T} = Tp^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[r^{t}] = Tp^{\star} - \sum_{t=0}^{T-1} \sum_{a \in A} \mathbb{P}\{Z_{a}^{t}\}\mathbb{E}[r^{t}|Z_{a}^{t}]$$
$$= Tp^{\star} - \sum_{t=0}^{T-1} \sum_{a \in A} \mathbb{E}[z_{a}^{t}]p_{a}$$

$$egin{aligned} \mathcal{R}_{\mathcal{T}} &= \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[r^t] = \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{P}\{Z_a^t\} \mathbb{E}[r^t | Z_a^t] \ &= \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{E}[z_a^t] \mathcal{p}_a = \left(\sum_{a \in \mathcal{A}} \mathbb{E}[u_a^T]
ight) \mathcal{p}^{\star} - \sum_{a \in \mathcal{A}} \mathbb{E}[u_a^T] \mathcal{p}_a \end{aligned}$$

$$\begin{aligned} \mathcal{R}_{T} &= \mathcal{T}\boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[\boldsymbol{r}^{t}] = \mathcal{T}\boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{P}\{\boldsymbol{Z}_{a}^{t}\} \mathbb{E}[\boldsymbol{r}^{t} | \boldsymbol{Z}_{a}^{t}] \\ &= \mathcal{T}\boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{z}_{a}^{t}]\boldsymbol{p}_{a} = \left(\sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}]\right) \boldsymbol{p}^{\star} - \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}]\boldsymbol{p}_{a} \\ &= \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}](\boldsymbol{p}^{\star} - \boldsymbol{p}_{a}) \end{aligned}$$

$$egin{aligned} &\mathcal{R}_{\mathcal{T}} = \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[r^t] = \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{P}\{Z_a^t\} \mathbb{E}[r^t|Z_a^t] \ &= \mathcal{T} \mathcal{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{E}[z_a^t] \mathcal{p}_a = \left(\sum_{a \in \mathcal{A}} \mathbb{E}[u_a^T]\right) \mathcal{p}^{\star} - \sum_{a \in \mathcal{A}} \mathbb{E}[u_a^T] \mathcal{p}_a \ &= \sum_{a \in \mathcal{A}} \mathbb{E}[u_a^T] (\mathcal{p}^{\star} - \mathcal{p}_a) = \sum_{a: \mathcal{p}_a
eq \mathcal{p}^{\star}} \mathbb{E}[u_a^T] \Delta_a. \end{aligned}$$

$$\begin{aligned} \mathcal{R}_{T} &= \mathcal{T} \boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \mathbb{E}[\boldsymbol{r}^{t}] = \mathcal{T} \boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{P}\{\boldsymbol{Z}_{a}^{t}\} \mathbb{E}[\boldsymbol{r}^{t} | \boldsymbol{Z}_{a}^{t}] \\ &= \mathcal{T} \boldsymbol{p}^{\star} - \sum_{t=0}^{T-1} \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{Z}_{a}^{t}] \boldsymbol{p}_{a} = \left(\sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}]\right) \boldsymbol{p}^{\star} - \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}] \boldsymbol{p}_{a} \\ &= \sum_{a \in \mathcal{A}} \mathbb{E}[\boldsymbol{u}_{a}^{T}](\boldsymbol{p}^{\star} - \boldsymbol{p}_{a}) = \sum_{a: \boldsymbol{p}_{a} \neq \boldsymbol{p}^{\star}} \mathbb{E}[\boldsymbol{u}_{a}^{T}] \Delta_{a}. \end{aligned}$$

To show the regret bound, we shall show for each sub-optimal arm *a* that

$$\mathbb{E}[u_a^T] = O\left(\frac{1}{(\Delta_a)^2}\log(T)\right).$$

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$$\mathbb{E}[u_a^T] = \sum_{t=0}^{T-1} \mathbb{E}[z_a^t]$$

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$$= \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\} + \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t \ge \bar{u}_a^T)\}$$

$$\mathbb{E}[u_a^T] = \sum_{t=0}^{T-1} \mathbb{E}[z_a^t] = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t\}$$
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$$= A + B.$$

To prove $\mathbb{E}[u_a^T] = O\left(\frac{1}{\Delta_a^2}\log(T)\right)$, we show $\mathbb{E}[u_a^T] \leq \bar{u}_a^T + C$ for constant *C*.

$$\mathbb{E}[u_a^T] = \sum_{t=0}^{T-1} \mathbb{E}[z_a^t] = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t\}$$
$$= \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\} + \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t \ge \bar{u}_a^T)\}$$
$$= A + B.$$

We show A is upper-bounded by \bar{u}_a^T and B is upper-bounded by a constant.

$$A = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\}$$

8/16

$$A = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\}$$
$$= \sum_{t=0}^{T-1} \sum_{m=0}^{\bar{u}_a^T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t = m)\}$$

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= $\sum_{m=0}^{\bar{u}_a^T - 1} \mathbb{P}\{Z_a^0, (u_a^0 = m) \text{ or } Z_a^1, (u_a^1 = m) \text{ or } \dots \text{ or } Z_a^{T-1}, (u_a^{T-1} = m)\}$

$$A = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\}$$

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= $\sum_{m=0}^{\bar{u}_a^T - 1} \mathbb{P}\{Z_a^0, (u_a^0 = m) \text{ or } Z_a^1, (u_a^1 = m) \text{ or } \dots \text{ or } Z_a^{T-1}, (u_a^{T-1} = m)\}$
 $\leq \sum_{m=0}^{\bar{u}_a^T - 1} 1$

Step 3: Bounding A

$$A = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t < \bar{u}_a^T)\}$$

= $\sum_{t=0}^{T-1} \sum_{m=0}^{\bar{u}_a^T - 1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t = m)\} = \sum_{m=0}^{\bar{u}_a^T - 1} \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t = m)\}$
= $\sum_{m=0}^{\bar{u}_a^T - 1} \mathbb{P}\{Z_a^0, (u_a^0 = m) \text{ or } Z_a^1, (u_a^1 = m) \text{ or } \dots \text{ or } Z_a^{T-1}, (u_a^{T-1} = m)\}$
 $\leq \sum_{m=0}^{\bar{u}_a^T - 1} 1 = \bar{u}_a^T.$

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 $\leq \sum_{m=0}^{\bar{u}_a^T - 1} 1 = \bar{u}_a^T.$

We have used the fact that for $0 \le i < j \le t - 1$, $(Z_a^i, (u_a^i = m))$ and $(Z_a^j, (u_a^j = m))$ are mutually exclusive.

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$$B = \sum_{t=0}^{T-1} \mathbb{P}\{Z_a^t \text{ and } (u_a^t \geq \bar{u}_a^T)\}$$

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$$\leq \sum_{t=n}^{T-1} \mathbb{P}\left\{\left(\hat{p}_a^t + \sqrt{\frac{2}{u_a^t}\ln(t)} \ge \hat{p}_\star^t + \sqrt{\frac{2}{u_\star^t}\ln(t)}\right) \text{ and } (u_a^t \ge \bar{u}_a^T)\right\}$$

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 $\le \sum_{t=n}^{T-1} \sum_{x=\bar{u}_a^T} \sum_{y=1}^t \mathbb{P}\left\{\hat{p}_a(x) + \sqrt{\frac{2}{x}\ln(t)} \ge \hat{p}_\star(y) + \sqrt{\frac{2}{y}\ln(t)}\right\} \text{ where}$

 $\hat{p}_a(x)$ is the empirical mean of the first *x* pulls of arm *a*, and $\hat{p}_{\star}(y)$ is the empirical mean of the first *y* pulls of arm \star .

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Step 4.2: Bounding *B* • Fix $x \in {\overline{u}_a^T, \overline{u}_a^T + 1, \dots, t}$ and $y \in {1, 2, \dots, t}$.

• Fix $x \in {\bar{u}_a^T, \bar{u}_a^T + 1, ..., t}$ and $y \in {1, 2, ..., t}$. 1. We have:

$$\hat{p}_a(x) + \sqrt{rac{2}{x}\ln(t)} \ge \hat{p}_\star(y) + \sqrt{rac{2}{y}\ln(t)}$$
 $\implies \left(\hat{p}_a(x) + \sqrt{rac{2}{x}\ln(t)} \ge p_\star\right) ext{ or } \left(\hat{p}_\star(y) + \sqrt{rac{2}{y}\ln(t)} < p_\star
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• Fix $x \in {\bar{u}_a^T, \bar{u}_a^T + 1, ..., t}$ and $y \in {1, 2, ..., t}$. 1. We have:

$$\begin{split} \hat{p}_{a}(x) + \sqrt{\frac{2}{x}\ln(t)} &\geq \hat{p}_{\star}(y) + \sqrt{\frac{2}{y}\ln(t)} \\ \implies \left(\hat{p}_{a}(x) + \sqrt{\frac{2}{x}\ln(t)} \geq p_{\star}\right) \text{ or } \left(\hat{p}_{\star}(y) + \sqrt{\frac{2}{y}\ln(t)} < p_{\star}\right). \end{split}$$

Fact: If $\alpha > \beta$, then $\alpha \ge \gamma$ or $\beta < \gamma$. Holds for arbitrary α, β, γ !

• Fix $x \in {\bar{u}_a^T, \bar{u}_a^T + 1, ..., t}$ and $y \in {1, 2, ..., t}$. 1. We have:

$$\begin{split} \hat{p}_{a}(x) + \sqrt{\frac{2}{x}\ln(t)} &\geq \hat{p}_{\star}(y) + \sqrt{\frac{2}{y}\ln(t)} \\ \implies \left(\hat{p}_{a}(x) + \sqrt{\frac{2}{x}\ln(t)} \geq p_{\star}\right) \text{ or } \left(\hat{p}_{\star}(y) + \sqrt{\frac{2}{y}\ln(t)} < p_{\star}\right). \end{split}$$

Fact: If $\alpha > \beta$, then $\alpha \ge \gamma$ or $\beta < \gamma$. Holds for arbitrary α, β, γ !

2. Since
$$x \ge \bar{u}_a^T$$
, we have $\sqrt{\frac{2}{x}\ln(t)} \le \sqrt{\frac{2}{\bar{u}_a^T}\ln(t)} \le \frac{\Delta_a}{2}$, and so
 $\hat{p}_a(x) + \sqrt{\frac{2}{x}\ln(t)} \ge p_\star \implies \hat{p}_a(x) \ge p_a + \frac{\Delta_a}{2}$.

Continuing from Step 4.1, using the two results from Step 4.2, and invoking Hoeffding's Inequality:

$$B \leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_a^T}^t \sum_{y=1}^t \mathbb{P}\left\{\hat{p}_a(x) + \sqrt{\frac{2}{x}\ln(t)} \geq \hat{p}_\star(y) + \sqrt{\frac{2}{y}\ln(t)}\right\}$$

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$$\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_a^T} \sum_{y=1}^t \left(\mathbb{P}\left\{\hat{p}_a(x) \geq p_a + \frac{\Delta_a}{2}\right\} + \mathbb{P}\left\{\hat{p}_\star(y) < p_\star - \sqrt{\frac{2}{y}\ln(t)}\right\}\right)$$

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$$\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_a^T} \sum_{y=1}^t \left(\mathbb{P}\left\{\hat{p}_a(x) \geq p_a + \frac{\Delta_a}{2}\right\} + \mathbb{P}\left\{\hat{p}_\star(y) < p_\star - \sqrt{\frac{2}{y}\ln(t)}\right\}\right)$$
$$\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_a^T} \sum_{y=1}^t \left(e^{-2x\left(\frac{\Delta_a}{2}\right)^2} + e^{-2y\left(\sqrt{\frac{2}{y}\ln(t)}\right)^2}\right)$$

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Continuing from Step 4.1, using the two results from Step 4.2, and invoking Hoeffding's Inequality:

$$\begin{split} B &\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_{a}^{T}} \sum_{y=1}^{t} \mathbb{P}\left\{ \hat{p}_{a}(x) + \sqrt{\frac{2}{x}}\ln(t) \geq \hat{p}_{\star}(y) + \sqrt{\frac{2}{y}}\ln(t) \right\} \\ &\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_{a}^{T}} \sum_{y=1}^{t} \left(\mathbb{P}\left\{ \hat{p}_{a}(x) \geq p_{a} + \frac{\Delta_{a}}{2} \right\} + \mathbb{P}\left\{ \hat{p}_{\star}(y) < p_{\star} - \sqrt{\frac{2}{y}}\ln(t) \right\} \right) \\ &\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_{a}^{T}} \sum_{y=1}^{t} \left(e^{-2x\left(\frac{\Delta_{a}}{2}\right)^{2}} + e^{-2y\left(\sqrt{\frac{2}{y}}\ln(t)\right)^{2}} \right) \\ &\leq \sum_{t=n}^{T-1} \sum_{x=\bar{u}_{a}^{T}} \sum_{y=1}^{t} \left(e^{-4\ln(t)} + e^{-4\ln(t)} \right) \leq \sum_{t=n}^{T-1} t^{2}\left(\frac{2}{t^{4}}\right) \leq \sum_{t=1}^{\infty} \frac{2}{t^{2}} = \frac{\pi^{2}}{3}. \end{split}$$

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We are done!

Shivaram Kalyanakrishnan (2023)

Multi-armed Bandits

- 1. Analysis of UCB
- 2. Other bandit problems

- In this course, we have covered
 - stochastic multi-armed bandits,
 - minimisation of expected cumulative regret.

There are many other variations/formulations.

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 - Arm 2 gives rewards 48 and 50, each w.p. 1/2.
 - Which arm would you prefer?

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 - Which arm would you prefer?
- What if the arms' (true) means vary over time?
 - Nonstationary setting, seen for example, in on-line ads.
 - Approach depends on nature of drift/change in rewards.
 - In practice, one might only trust most recent data from arms.
 - In practice, the set of arms can itself change over time!

- Pure exploration.
 - Separate "testing" and "live" phases.
 - In testing phase, rewards don't matter.
 - ▶ PAC formulation: W.p. at least 1δ , must return an ϵ -optimal arm, while incurring a small number of pulls.
 - Simple regret formulation: Given a budget of *T* pulls, must output an arm *a* such that p_a is large, or equivalently, simple regret = $p^* p_a$ is small).

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- Limited number of feedback stages.
 - Suppose you are given budget *T*, but your algorithm can look at history only *s* < *T* times?
 - UCB, Thompson Sampling, etc. are fully sequential (s = T).
 - How to manage with fewer "stages" s?

- What if the number of arms is large (thousands, millions)?
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- What if we are interacting with many bandits simultaneously?
 - Contextual bandits: If the bandits themselves can be described using features (a "context"), data from one can be used to generate estimates about others.
- What if the rewards do not come from a fixed random process?
 - Adversarial bandits make no assumption on the rewards.
 - Possible to show sub-linear regret when compared against playing a single arm for the entire run.
 - Necessary to use a randomised algorithm.

Multi-armed Bandits

- The exploration-exploitation dilemma
- Definitions: Bandit, Algorithm
- *e*-greedy algorithms
- Evaluating algorithms: Regret
- Achieving sub-linear regret
- A lower bound on regret
- UCB, KL-UCB algorithms
- Thompson Sampling algorithm
- Concentration bounds
- Understanding Thompson Sampling
- Other bandit problems

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• Next class: Markov Decision Problems