CS 747, Autumn 2023: Lecture 11

Shivaram Kalyanakrishnan

Department of Computer Science and Engineering Indian Institute of Technology Bombay

Autumn 2023

Markov Decision Problems

- 1. Policy iteration: variants and complexity bounds
- 2. Analysis of bounds
 - Basic tools
 - Howard's PI with k=2
 - BSPI with k=2
 - Open problems
- 3. Review of MDP planning

Markov Decision Problems

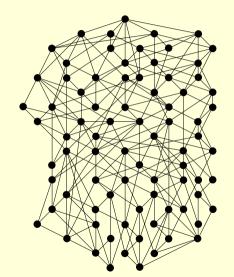
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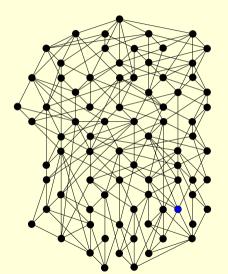
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\pi' \leftarrow \text{PolicyImprovement}(\pi).
\pi \leftarrow \pi'.

Return \pi.
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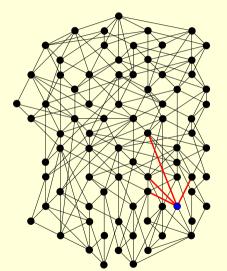
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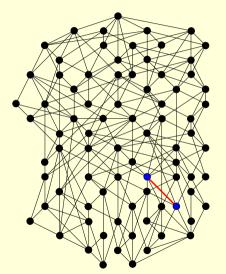
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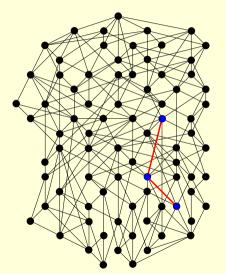
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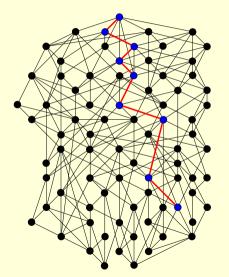
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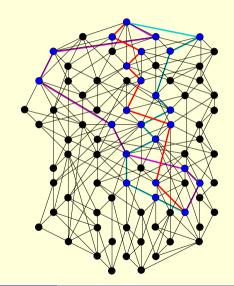
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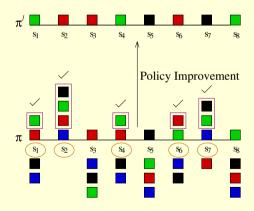
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Path taken (and hence the number of iterations) in general depends on the switching strategy.



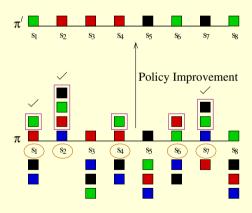
Howard's Policy Iteration

- Reference: Howard (1960).
- Greedy; switch all improvable states.



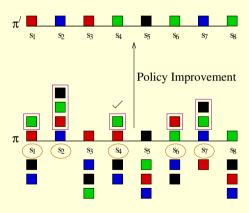
Random Policy Iteration

- Reference: Mansour and Singh (1999).
- Switch a non-empty subset of improvable states chosen uniformly at random.



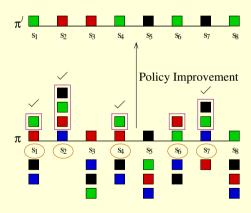
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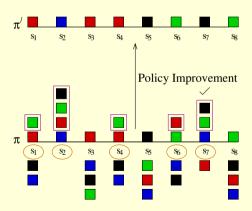
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Simple Policy Iteration

- Reference: Melekopoglou and Condon (1994).
- Assume a fixed indexing of states.
- Switch the improvable state with the highest index.



Upper and Lower Bounds

U(n, k) is an upper bound applicable to a set of PI variants \mathcal{L} if

- for each *n*-state, *k*-action MDP $M = (S, A, T, R, \gamma)$,
- for each policy $\pi: S \to A$,
- for each algorithm $L \in \mathcal{L}$,

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X(n, k) is a lower bound applicable to a set of PI variants \mathcal{L} if

- there exists an *n*-state, *k*-action MDP $M = (S, A, T, R, \gamma)$,
- there exists a policy $\pi: \mathcal{S} \to \mathcal{A}$,
- there exists an algorithm $L \in \mathcal{L}$, such that the expected number of policy evaluations performed by L on M if

initialised at π is at least X(n, k).

Switching Strategies and Bounds

Upper bounds on number of iterations

PI Variant	Type	k = 2	General k
Howard's (Greedy) PI [H60, MS99]	Deterministic	$O\left(\frac{2^n}{n}\right)$	$O\left(\frac{k^n}{n}\right)$
Mansour and Singh's Random PI [MS99]	Randomised	1.7172 ⁿ	$pprox O\left(\frac{k}{2}\right)^n$
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Lower bounds on number of iterations

- $\Omega(n)$ Howard's PI on *n*-state, 2-action MDPs [HZ10].
- $\Omega(2^n)$ Simple PI on *n*-state, 2-action MDPs [MC94].

(Polynomial factors ignored. Authors with names underlined once took CS 747!)

• Kalyanakrishnan, Mall, and Goyal (2016) devise the Batch-switching PI algorithm (deterministic), and show an upper bound of 1.6479^n for k = 2.

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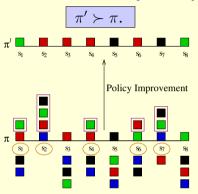
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- Ashutosh, Consul, Dedhia, Khirwadkar, Shah, and Kalyanakrishnan (2020) show a *lower bound* of \sqrt{k}^n iterations for a deterministic variant of PI.

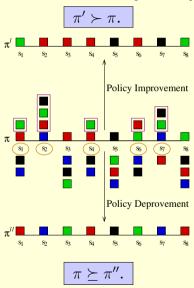
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1. Policy Improvement and Policy "Deprovement"



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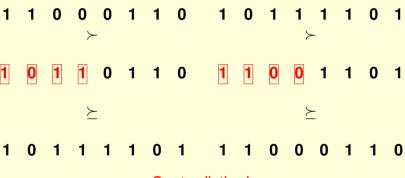


- 1 1 0 0 0 1 1 0
- 1 0 1 1 0 1 1 0 0 1 1 0 0 1 1 0 1



2. Improvement sets in 2-action MDPs

Non-optimal policies $\pi,\pi'\in\Pi$ cannot have the same set of improvable states.



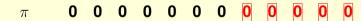
Contradiction!

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Switch actions in every improvable state.

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Switch actions in every improvable state.

 π' 0 0 0 0 0 0 0 1 1 1 1 1

 π 0 0 0 0 0 0 0 0 0 0 0 0 0

Switch actions in every improvable state.

Possible?

 π' 0 0 0 0 0 0 1 1 1 1 1

 π 0 0 0 0 0 0 0 0 0 0 0 0

Switch actions in every improvable state.

 π' 0 0 0 0 0 0 1 1 1 1 1

 π 0 0 0 0 0 0 0 0 0 0 0 0

Switch actions in every improvable state.

$$\pi'$$
 0 0 0 0 0 0 0 1 1 1 1 1 π_1 0 0 0 0 0 0 0 1 1 1 1 0

$$\pi$$
 0 0 0 0 0 0 0 0 0 0 0 0

Switch actions in every improvable state.

 π 0 0 0 0 0 0 0 0 0 0 0

Switch actions in every improvable state.

π'	0	0	0	0	0	0	0	1	1	1	1	1
π_{1}	0	0	0	0	0	0	0	1	1	1	1	0
π_2	0	0	0	0	0	0	0	1	1	1	0	0
π_{3}	0	0	0	0	0	0	0	1	1	0	0	0
π	0	0	0	0	0	0	0	0	0	0	0	0

Switch actions in every improvable state.

π'	0	0	0	0	0	0	0	1	1	1	1	1
π_{1}	0	0	0	0	0	0	0	1	1	1	1	0
π_2	0	0	0	0	0	0	0	1	1	1	0	0
π_{3}	0	0	0	0	0	0	0	1	1	0	0	0
π_{4}	0	0	0	0	0	0	0	1	0	0	0	0
π	0	0	0	0	0	0	0	0	0	0	0	0

Switch actions in every improvable state.

If π has \underline{m} improvable states and $\pi \xrightarrow{\text{Howard's PI}} \pi'$, then there exist \underline{m} policies π'' such that $\pi' \succeq \pi'' \succ \pi$.

• Take $m^* = \frac{n}{3}$.

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$$\leq \binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \cdots + \binom{n}{m^{\star} - 1}$$

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Number of iterations taken by Howard's PI: $O(\frac{2^n}{n})$ [MS99, HGDJ14].

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Batch-Switching Policy Iteration (BSPI) (2-action MDPs)

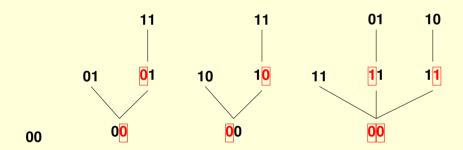
Howard's Policy Iteration takes at most iterations on a 2-state MDP!

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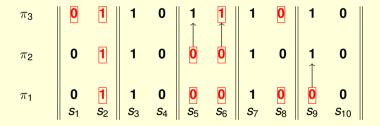
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Batch-Switching Policy Iteration (BSPI) (2-action MDPs)

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π_2	0	1	1	0	0	0	1	0	1 ↑	0
π_1	0 s ₁	1	1 s ₃	0 S ₄	0	0 S ₆	1 s ₇	0 S ₈	0 S 9	0 0 s ₁₀



π_{4}	0	1	1	0	1	1	1	1 ↑	1	0
π_3	0	1	1	0	1 ↑	1	1	0	1	0
π_2	0	1	1	0	0	0	1	0	1	0
π_1	0 s ₁	1	1 s ₃	0	0 S ₅	0	1 s ₇	0	0 S 9	0 s ₁₀

Partition the states into 2-sized batches; arranged from left to right. Given a policy, improve the rightmost set containing an improvable state.

π_4	0	1	1	0	1	1	1	1 ↑	1	0
π_3	0	1	1	0	1 ↑	1	1	0	1	0
π_2	0	1	1	0	0	0	1	0	1 ↑	0
π_1	0 s ₁	1	1 s ₃	0 s ₄	0	0	1 s ₇	0	0 S 9	0 s ₁₀

• Left-most batch can change only when all other columns are non-improvable.

π_4	0	1	1	0	1	1	1	1 ↑	1	0
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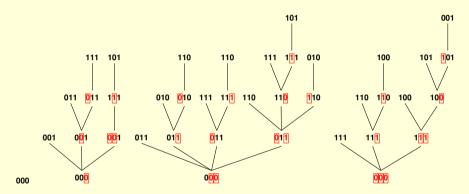
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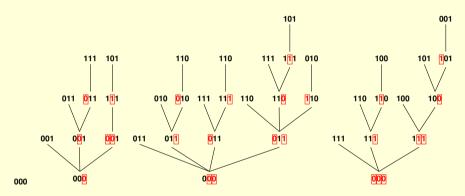
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- $T(n) \leq 3 \times T(n-2) \leq \sqrt{3}^n$.

Howard's Policy Iteration takes at most 5 iterations on a 3-state MDP!

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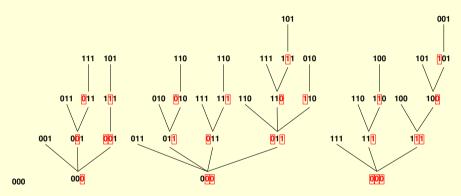


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The structures above are called Trajectory-bounding Trees (TBTs) [KMG16a] (and correspond to the Order Regularity Problem [H12, GHDJ15]).

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The structures above are called Trajectory-bounding Trees (TBTs) [KMG16a] (and correspond to the Order Regularity Problem [H12, GHDJ15]).

BSPI with 3-sized batches gives $T(n) \le 5 \times T(n-3) \le 1.71^n$.

Batch size	Depth of TBT	Bound on number of iterations
1	2	2 ⁿ
2	3	1.7321 ⁿ
3	5	1.7100 ⁿ
4	8	1.6818 ⁿ
5	13	1.6703 ⁿ
6	21	1.6611 ⁿ
7	33	1.6479 ⁿ

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Depth of TBT for batch size 7 due to Gerencsér et al. [GHDJ15].

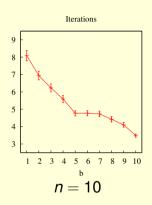
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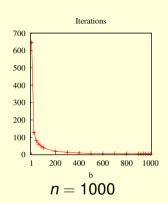
Depth of TBT for batch size 7 due to Gerencsér *et al.* [GHDJ15]. Will the bound continue to be non-increasing in the batch size?

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Depth of TBT for batch size 7 due to Gerencsér *et al.* [GHDJ15]. Will the bound continue to be non-increasing in the batch size? If so, 1.6479ⁿ would be an upper bound for Howard's Policy Iteration!

BSPI: Effect of Batch Size b





Averaged over *n*-state, 2-action MDPs with randomly generated transition and reward functions. Each point is an average over 100 randomly-generated MDP instances and initial policies [KMG16a].

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Open Problems

- Is the complexity of Howard's PI on 2-action MDPs upper-bounded by the Fibonacci sequence (≈ 1.6181ⁿ)?
- Is Howard's PI the most efficient among deterministic PI algorithms (worst case over all MDPs)?
- Is there a super-linear lower bound on the number of iterations taken by Howard's PI on 2-action MDPs?
- Is Howard's PI strongly polynomial on deterministic MDPs?
- Is there a variant of PI that can visit all k^n policies in some n-state, k-action MDP—implying an $\Omega(k^n)$ lower bound?
- Is there a strongly polynomial algorithm for MDP planning?

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Summary of MDP Planning

- MDPs are an abstraction of sequential decision making.
- Many applications; many different formulations.
- Essential solution concept: optimal policy (known to exist).
- Three main families of planning algorithms: value iteration, linear programming, policy iteration.
- Have strengths and weaknesses in theory and in practice. Can combine.
- We showed correctness of all three methods.
- Used Banach's fixed-point theorem, Bellman (optimality) operator.
- What if T, R were not given, but have to be learned from interaction? Can we still learn to act optimally?
- Yes: that's the reinforcement learning problem. Next class!