

CS101 Computer Programming and Utilization

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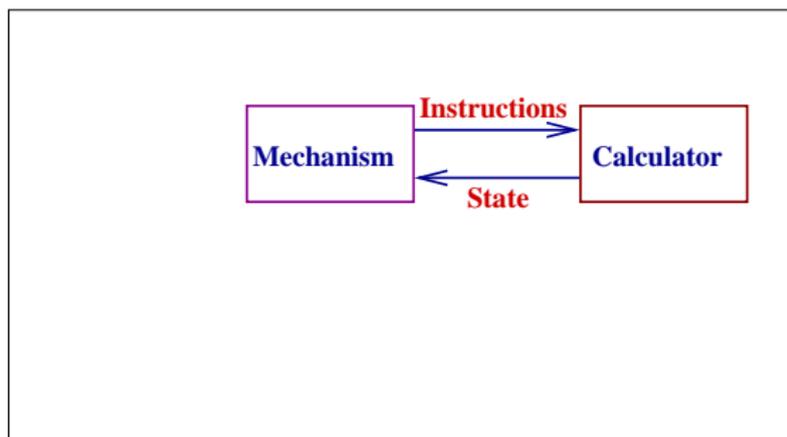
May 10, 2006

1 The Basic Computer

2 Programming Language

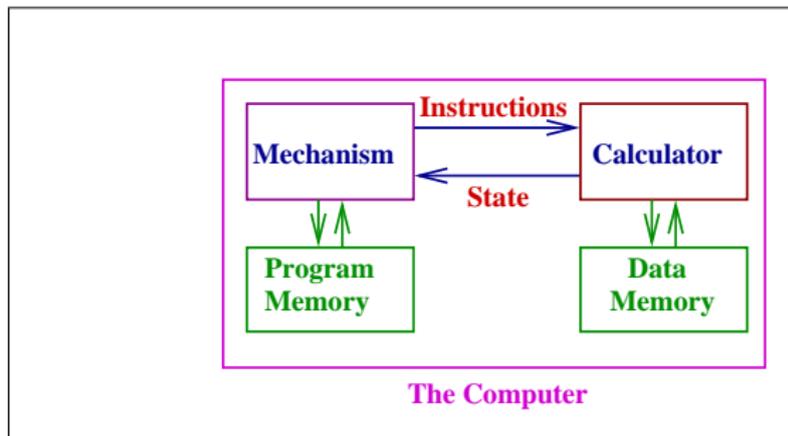
- READIN and assignments
- The IF-ENDIF instruction
- The DO-WHILE instruction

In Summary



- We started off with the basic calculator and the BUM who executed our programs.
- Next we introduced more **memory** in the calculator so that programs became simpler.
- Finally, we replace the BUM by a cleverer mechanism:
 - ▶ who stored the program that we gave him.
 - ▶ could execute the **TEST nos** instruction and re-use the program code.
- Then we saw how to write some programs in for such a composite machine.

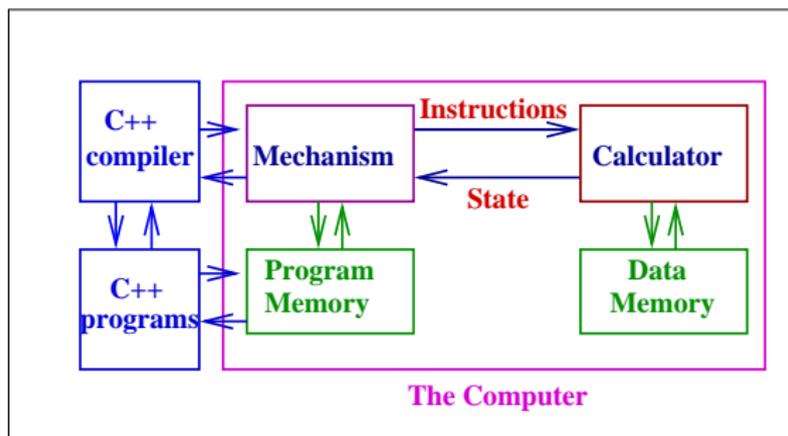
The Basic Computer



The **basic computer** is exactly this machine:

- It is an enhanced calculating machine with a richer instruction set for specific calculations.
- It has enhanced data memory (registers) which can store 10^9 items.
- It has a mechanism which passes instructions to the calculator.
- It has a program memory, wherein the program to be executed is stored.

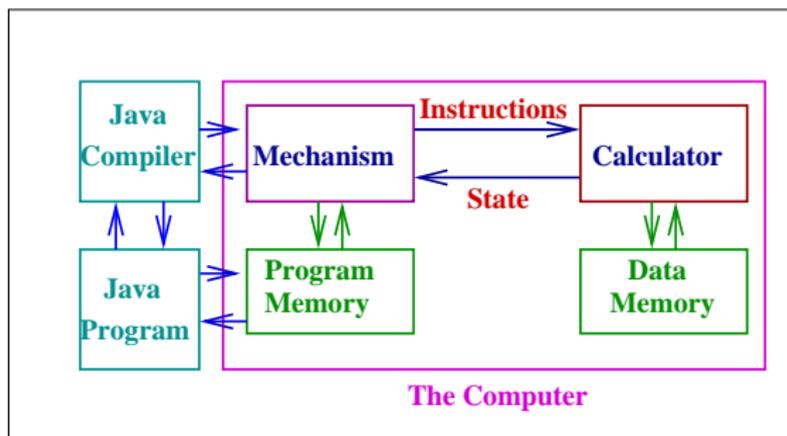
Programming Languages



Different programming languages such as **C++**, **Java** are **front ends** to the basic computer. These languages

- Allow the user to write programs in a more conceptual language.
- Translate this into the calculator language that we know.
- Store this translation into the program memory.

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- Allow the user to write programs in a more conceptual language.
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A Simple Programming Language

- A simple instruction:

```
M3=READIN 78
```

unfolds into

```
78 % put into display
```

```
STO 3 % put it into M1
```

- the instruction:

```
M3=READIN
```

prompts the user to input a
number **nos**

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This instructions puts user values
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This instructions puts user values
into memory locations.

- Another instruction: The
ASSIGNMENT:

```
M1= M1 + 5 * M3
```

unfolds into

```
RCL 1
+
5
*
RCL 3
=
STO 1
```

This instruction allows quick
programming of arithmetic
operations.

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In short, the new instructions saves us from writing long programs for
conceptually easy steps.

The Quadratic Equation $x^2 + 3x + 2$ Revisited

```
M1=READIN 1 % A read in M1
M2=READIN 3 % B read in M2
M3=READIN 2 % C read in M3
```

This finishes the initialization. M6 and M7 contain the constants 2 and 4.

```
M4= M2*M2-4*M1*M3 % the discriminant
M4= M4 SQRT          % completed
```

This computes the discriminant.

```
M5= M2 MINUS          % M5=-B
M5= M5+M4 DIV 2 DIV M1 % root 1
```

Finally the root. Note that READIN statements are easy but **ASSIGNMENT** statements need some care.

Let us analyse the first two **ASSIGNMENT** statements:

```
M4= M2*M2-4*M1*M3      % the discriminant
M4= M4 SQRT            % completed
```

The first statement expands to:

```
RCL 2
*
RCL 2
-
4
*
RCL 1
*
RCL 3
=
STO 4
```

Given the **current values** of the registers, M4 contains $B^2 - 4AC$.

Let us analyse the first two **ASSIGNMENT** statements:

```
M4= M2*M2-4*M1*M3      % the discriminant
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*
RCL 2
-
4
*
RCL 1
*
RCL 3
=
STO 4
```

Given the **current values** of the registers, M4 contains $B^2 - 4AC$.

The next assignment statement is peculiar:

```
M4= M4 SQRT
```

This translates to:

```
RCL 4
SQRT
=
STO 4
```

The **current value** of M4 is used to obtain the **next value** of M4 which is $\sqrt{B^2 - 4AC}$.

The IF-ENDIF instructions

The **IF** instructions is used as follows:

```
IF M4
```

unfolds into:

```
RCL 4  
TEST nos
```

The argument **nos** is captured by the **ENDIF** instructions as follows:

```
ENDIF
```

This records the line number of the next instruction.

The IF-ENDIF instructions

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IF M4
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unfolds into:

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RCL 4  
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The argument **nos** is captured by the **ENDIF** instructions as follows:

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ENDIF
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This records the line number of the next instruction.

```
M1=READIN 1 % A read in M1  
M2=READIN 3 % B read in M2  
M3=READIN 2 % C read in M3
```

```
M4= M2*M2-4*M1*M3 % the discriminant
```

```
IF M4          %M4>0 then go to nos  
STOP  
ENDIF          %this is nos
```

```
M4= M4 SQRT          % completed  
M5= M2 MINUS          % M5=-B  
M5= M5+M4 DIV 2 DIV M1 % root 1
```

In other words:

```
CODE BLOCK 1
```

```
IF M4
```

```
CODE BLOCK 2
```

```
ENDIF
```

```
CODE BLOCK 3
```

In other words:

```
CODE BLOCK 1  
IF M4  
CODE BLOCK 2  
ENDIF  
CODE BLOCK 3
```

causes the following two possibilities:

- if $M4 > 0 \Rightarrow$
CodeBlock1;CodeBlock3.
- if $M4 \leq 0 \Rightarrow$ Code-
Block1;CodeBlock2;CodeBlock3.

Warning

The ENDIF of the IF **must follow** the IF.

In other words:

```
CODE BLOCK 1
IF M4
CODE BLOCK 2
ENDIF
CODE BLOCK 3
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causes the following two possibilities:

- if $M4 > 0 \Rightarrow$ CodeBlock1;CodeBlock3.
- if $M4 \leq 0 \Rightarrow$ CodeBlock1;CodeBlock2;CodeBlock3.

Warning

The ENDIF of the IF **must follow** the IF.

Assignment

- Write PL-code for computing the other root.
- Expand the last two ASSIGNMENT statements into CAL-code.
- Modify the quadratic programming code to take care of $a \neq 0$.
- Write PL-code for computing 2^n .

The DO-WHILE instruction

Here is another useful instruction:

```
DO
```

merely **records** the line number of the **next** instruction say **nos** as it scans the program.

The **DO** instruction must be coupled with the **WHILE** instruction:

```
WHILE M5
```

Let M10 be an **unused** register, The above instruction causes the following output:

```
M10=M5;  
RCL 10  
TEST nos
```

summary...

The **DO** records the line number of the **next** instruction. Thus, the presence of a **WHILE** causes the execution to go to **nos** if $M5 > 0$. Otherwise the next statement is executed.

Here is the **log** example again:

```
M1=READIN 178 % the value of n
M2=0 % this stores log
M3=1 % this stores 2^log
M4=M1-M3
```

```
D0 * nos=5
```

```
M2=M2+1 % add 1
M3=M3*10 % multiply by 10
M4=M1-M3
```

```
WHILE M4 * M10=M4
RCL 10
TEST nos
```

```
STOP
```

Here is the **log** example again:

```
M1=READIN 178 % the value of n
M2=0 % this stores log
M3=1 % this stores 2^log
M4=M1-M3
```

```
DO * nos=5
```

```
M2=M2+1 % add 1
M3=M3*10 % multiply by 10
M4=M1-M3
```

```
WHILE M4 * M10=M4
RCL 10
TEST nos
```

```
STOP
```

Let us see what happens:

- The first time the **DO** instruction is encountered, the line number is noted of the **next** instruction, which is **5**.
- Next:

	M1	M2	M3	M4
do 1	178	0	1	177
while 1	178	1	10	168
do 2	178	1	10	168
while 1	178	2	100	78
do 3	178	2	100	78
while 1	178	3	1000	-822
STOP				

DO-WHILE abstracted

The following code

```
CODE BLOCK 1  
  
DO  
  
CODE BLOCK 2  
  
WHILE M4  
  
CODE BLOCK 3
```

causes the following execution:

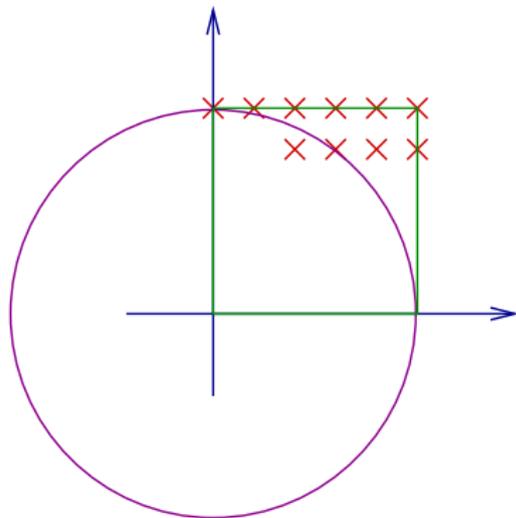
```
CB1  
CB2 first time (always)  
CB2 M4 > 0  
CB2 M4 > 0  
CB2 M4 non-positive  
CB3
```

Caution

The **WHILE** must always come after the **DO**.

Compute $\pi/4$

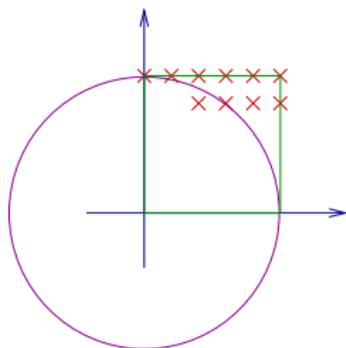
```
M1=READIN 100
M2=1 DIV M1 % the delta
M3=0; % count
M11=1
do
  M10=1
  do
    M4=M10*M10+M11*M11-1
    IF M4
      M3=M3+1
    ENDIF
    M10=M10-M2
  while M10
  M11=M11-M2
while M11
M3=M3 DIV M1 DIV M1
```



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      M3=M3+1
    ENDIF
    M10=M10-M2
  while M10
  M11=M11-M2
while M11

M3=M3 DIV M1 DIV M1
```

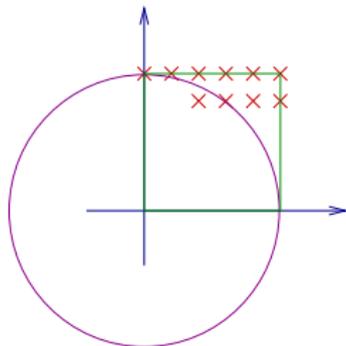


- M11 changes only in the **green loop**. Thus it is constant in the **blue loop** and the **IF-ENDIF**.
- For this fixed value of M11, M10 is initialized to 1. In the **blue loop**, this value goes from M10=1, 0.99,... upto M10=0.01. Thus the **IF-ENDIF** is executed exactly 100 times for each value of M11.

Compute $\pi/4$

```
M1=READLN 100
M2=1 DIV M1 % the delta
M3=0; % count
M11=1
do
M10=1
do
M4=M10*M10+M11*M11-1
IF M4
M3=M3+1
ENDIF
M10=M10-M2
while M10
M11=M11-M2
while M11
```

```
M3=M3 DIV M1 DIV M1
```



- At $M10=0.0$, the blue loop stops and a new value of $M11$ is computed.
- Thus there are 100×100 iterations of the IF-ENDIF which counts the number of points in the circle. Finally, the approximation to $\pi/4$ is computed.

Problem

- Given the following sequence of valid brackets, tell which open-brackets match with which closed bracket.

1	2	3	4	5	6	7	8	9	10	11	12
{	{	{	}	}	{	}	}	{	{	}	}

- Given a sequence a open and close brackets, how will you detect if it is a valid sequence?