
CS206 Tutorial No. #1

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1. Using the \wedge_i , \wedge_e , \vee_i , \vee_e , \rightarrow_i , \rightarrow_e , \perp_i (also called \neg_e), \perp_e , \neg_i , $\neg\neg_e$, $\neg\neg_i$ rules, Modus Tollens and **zero applications** of LEM, prove the following sequents.

(a) $\phi_1 \rightarrow (\phi_2 \rightarrow \neg\phi_1) \vdash \phi_2 \rightarrow \neg\phi_1$

(b) Distributive Laws:

- i. $(\phi_1 \wedge \phi_2) \vee (\phi_1 \wedge \phi_3) \vdash \phi_1 \wedge (\phi_2 \vee \phi_3)$
- ii. $\phi_1 \wedge (\phi_2 \vee \phi_3) \vdash (\phi_1 \wedge \phi_2) \vee (\phi_1 \wedge \phi_3)$
- iii. $(\phi_1 \vee \phi_2) \wedge (\phi_1 \vee \phi_3) \vdash \phi_1 \vee (\phi_2 \wedge \phi_3)$
- iv. $\phi_1 \vee (\phi_2 \wedge \phi_3) \vdash (\phi_1 \vee \phi_2) \wedge (\phi_1 \vee \phi_3)$

(c) De Morgan's Laws:

- i. $\neg(\phi_1 \vee \phi_2) \vdash \neg\phi_1 \wedge \neg\phi_2$
- ii. $\neg\phi_1 \wedge \neg\phi_2 \vdash \neg(\phi_1 \vee \phi_2)$
- iii. $\neg(\phi_1 \wedge \phi_2) \vdash \neg\phi_1 \vee \neg\phi_2$
- iv. $\neg\phi_1 \vee \neg\phi_2 \vdash \neg(\phi_1 \wedge \phi_2)$

(d) $\vdash \phi_1 \rightarrow (\phi_2 \rightarrow (\neg\phi_3 \vee (((\phi_4 \vee \phi_5 \rightarrow \phi_6) \rightarrow (\phi_4 \rightarrow \phi_6) \wedge (\phi_5 \rightarrow \phi_6))) \vee ((\phi_7 \rightarrow (\phi_8 \rightarrow \phi_9)) \rightarrow ((\phi_7 \rightarrow \phi_8) \rightarrow \phi_9))))$

2. Let p and q be atomic propositions that take values from the set $\{True, False\}$. Consider the following two formulae: $\phi_1 = (p \rightarrow \neg\phi_2)$, and $\phi_2 = (q \rightarrow \neg\phi_1)$.

(a) Show using natural deduction that $\vdash \phi_1 \vee \phi_2$.

Note: You may use LEM *almost once* in the proof. Other than this, you must use only the basic introduction and elimination rules of natural deduction.

(b) Show that there are exactly *two pairs* of propositional logic formulae (ϕ_1, ϕ_2) that satisfy the above definitions. Also, give justification for your answer.