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## CS206 Tutorial No. #6

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1. Use Natural Deduction to prove following first order predicate formulae

(a)

$$\begin{array}{l} \forall(x)[P(x) \rightarrow (x = b)], \quad P(b), \\ \forall(x)[(x = b) \rightarrow P(x)] \vdash \forall(x)[\forall(y)[P(x) \wedge P(y) \rightarrow (x = y)]] \end{array}$$

(b)

$$\begin{array}{l} \forall(x)[P(x) \vee Q(F(x))], \\ \forall(x)[R(F(x)) \rightarrow \neg P(x)], \vdash \exists(x)[Q(F(x))] \\ \forall(x)[R(F(x))] \end{array}$$

(c)

$$\begin{array}{l} \forall(x)[\forall(y)[(P(x, y) \wedge P(y, x)) \rightarrow (x = y)], \\ \forall(x)[P(x, F(x))], \vdash \exists(x)[u = F(u)] \\ \exists(z)[\forall(x)[P(x, z)]] \end{array}$$

2. Write the Prenex normal form of the formula:  $\exists x(P(x) \wedge \forall y(Q(x, y) \wedge \forall x \forall w S(x, y, w) \wedge \exists y R(x, y)))$   
Argue that “it is not the case that only infinite models exist for this formula”.