

CS 735: Formal Models for Concurrent and Asynchronous Systems

– Introduction

Instructor: S. Akshay

Jan-Apr 2024

Course hours: Slot09,
Mondays and Thursdays 3:30-5:00pm
Office hours: To be announced

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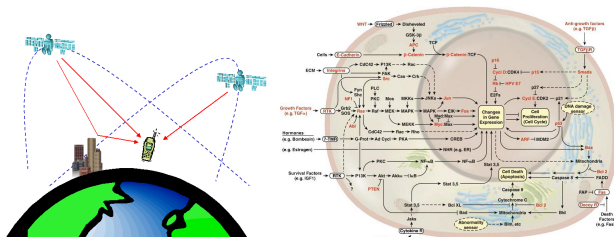
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Queries: Email me with [CS735-2024] in subject line
akshayss@cse.iitb.ac.in

Goal

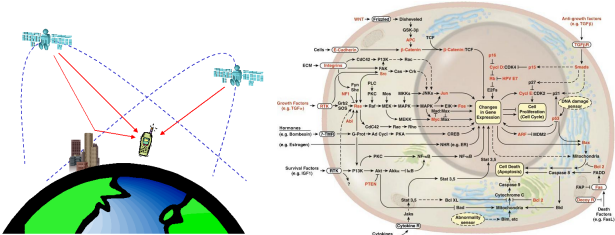
Formal Models for distributed and infinite-state systems



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Formal Models for distributed and infinite-state systems

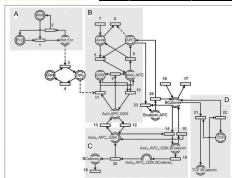
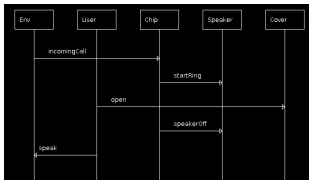
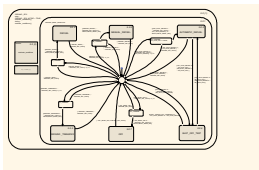
- **Distributed:** Concurrent, asynchronous, communicating,...



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- ▶ **Distributed:** Concurrent, asynchronous, communicating,...
- ▶ **Formal models:** Mathematical description, graphical notations, Automata models



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Questions that we will tackle

- ▶ **Analysis** of such models
- ▶ **Characterization**, relations
- ▶ **Underlying** properties, generalizations

Course contents

Topics and models that we will cover in this course:

1. Petri nets
2. Well-structured transition systems
3. Distributed automata models and their behaviors
4. Advanced topics, extensions and applications

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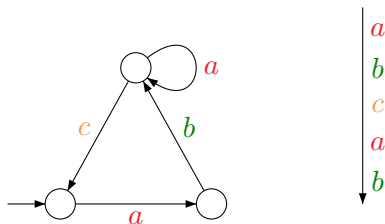
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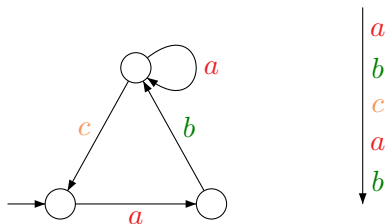
- ▶ Concurrency in Programs
- ▶ Concurrency and Quantities (Time/Probabilities)

Automata



- ▶ Behaviours are words, i.e., sequences of actions over a finite alphabet $\Sigma = \{a, b, c\}$.

Automata

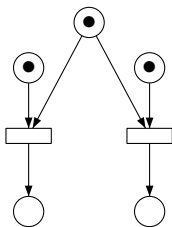


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Questions

- ▶ How shall we distribute it?
- ▶ How shall we add concurrent behaviors?

Petri Nets

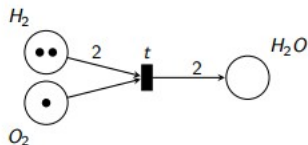


- ▶ An old model for distributed systems
 - ▶ invented by Carl Petri (-at the age of 13- in 1939? or '62)
 - ▶ to model resource consumption and so on...

Examples of Petri nets

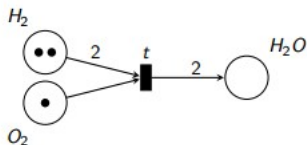
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- ▶ A library
- ▶ A producer-consumer example
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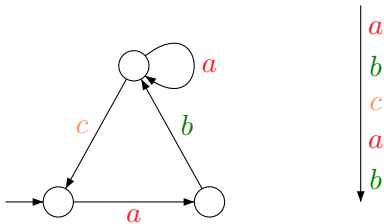
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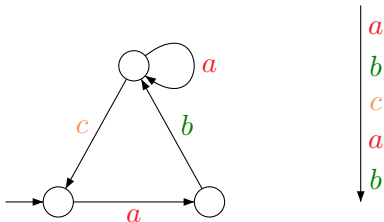
Applications: Business process models, stochastic processes, biological networks and so on

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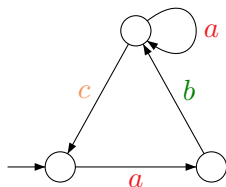


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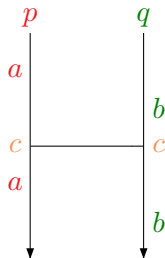
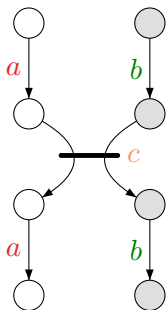
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Asynchronous Automata



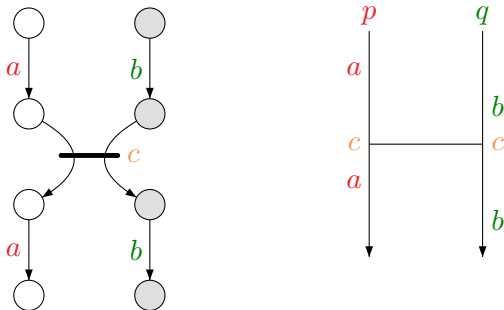
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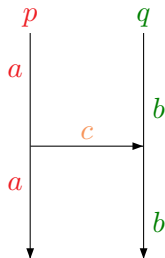
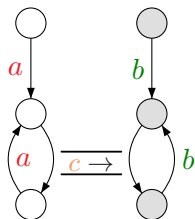
- ▶ Actions are distributed across processes (with sharing!)
- ▶ Some actions are shared, e.g., c is allowed only if both p and q move on c .

Asynchronous Automata



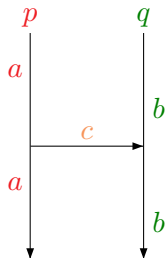
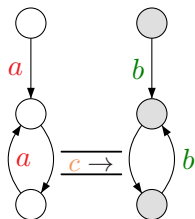
- ▶ What are the properties of languages accepted by such automata? E.g. above accepts $\{\underline{abcab}, bacab, bacba, abcba\}$.
- ▶ Given a language L , (when) can it be accepted by such an asynchronous automaton?

Message Passing Automata



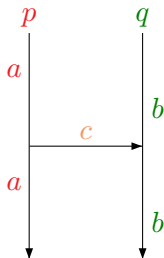
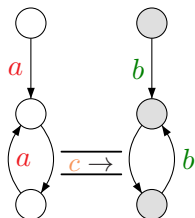
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- ▶ (Surprising fact: If you are allowed to lose messages randomly then it is decidable!) These are called Lossy channel systems.

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What are good formal models for concurrent programs?

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Does this capture reality of programs in today's world?

Modeling concurrent programs

Two issues

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Modeling concurrent programs

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- ▶ In the multi-processor world: memory access is no longer atomic!
- ▶ There is no non-determinism! How to avoid exploring runs

Leads to:

1. Weak memory models
2. Partial order reduction techniques

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Another face of concurrency

- ▶ For decomposability of large systems!
- ▶ Reasoning and composing systems

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- ▶ Under-approximate verification
- ▶ Fixed-point approaches

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Pictures and Mathematics

- ▶ How do you write these objects mathematically?
- ▶ Why write them mathematically?

Some take-aways from this course

- ▶ Different formal models for distributed systems
- ▶ Mathematical formalisms that reason about (the infinite) behaviors of such systems.
- ▶ Techniques to automatically analyze such systems.
- ▶ How to use them and where they are applied.

Logistics

Evaluation (flexible/tentative... upto a point)

- ▶ Continuous evaluation - assignments/quizzes : 35%
- ▶ Exam (Midsem/Endsem): 35 %
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Course material, references will be posted at

- ▶ <http://www.cse.iitb.ac.in/~akshayss/teaching.html>
- ▶ Piazza will be set up soon