Lecture 3: Process API

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What API does the OS provide to user programs?

- API = Application Programming Interface
 - = functions available to write user programs
- API provided by OS is a set of "system calls"
 - System call is a function call into OS code that runs at a higher privilege level of the CPU
 - Sensitive operations (e.g., access to hardware) are allowed only at a higher privilege level
 - Some "blocking" system calls cause the process to be blocked and descheduled (e.g., read from disk)

So, should we rewrite programs for each OS?

- POSIX API: a standard set of system calls that an OS must implement
 - Programs written to the POSIX API can run on any POSIX compliant OS
 - Most modern OSes are POSIX compliant
 - Ensures program portability
- Program language libraries hide the details of invoking system calls
 - The printf function in the C library calls the write system call to write to screen
 - User programs usually do not need to worry about invoking system calls
 Jubc --> Syscall

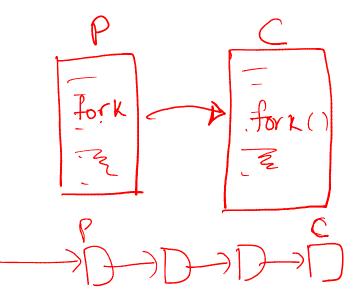
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Process related system calls (in Unix)

- fork() creates a new child process
 All processes are created by forking from a parent
 The init process is ancestor of all processes
- <u>exec()</u> makes a process execute a given executable
- exit() terminates a process
- wait() causes a parent to block until child terminates
- Many variants exist of the above system calls with different arguments

What happens during a fork?

- A new process is created by making a copy of parent's memory image
- The new process is added to the OS process list and scheduled
- Parent and child start execution just after fork (with different return values)
- Parent and child execute and modify the memory data independently



```
#include <stdio.h>
1
   #include <stdlib.h>
2
                                                YC
                                                              ſĊ.
   #include <unistd.h>
3
4
5
   int
                                               PIC
6
   main(int argc, char *argv[])
7
       printf("hello world (pid:%d)\n", (int) getpid());
8
       int rc = fork();
9
       if (rc < 0) {
                           // fork failed; exit
10
           fprintf(stderr, "fork failed\n");
11
          exit(1);
12
       } else if (rc == 0) { // child (new process)
13
14
       } else {
15
16
17
18
       return 0;
19
20
```

```
Figure 5.1: Calling fork() (p1.c)
```

When you run this program (called p1.c), you'll see the following:

```
prompt> ./p1
hello world (pid:29146)
hello, I am parent of 29147 (pid:29146)
hello, I am child (pid:29147)
prompt>
```

Waiting for children to die...

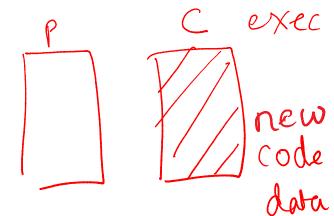
- Process termination scenarios
 - By calling exit() (exit is called automatically when end of main is reached)
 - OS terminates a misbehaving process
- Terminated process exists as a zombie
- When a parent calls wait(), zombie child is cleaned up or "reaped"
- wait() blocks in parent until child terminates (non-blocking ways to invoke wait exist)
- What if parent terminates before child? init process adopts orphans and reaps them

```
#include <stdio.h>
1
    #include <stdlib.h>
2
    #include <unistd.h>
3
    #include <sys/wait.h>
4
5
6
    int
    main(int argc, char *argv[])
7
8
        printf("hello world (pid:%d) \n", (int) getpid());
9
        int rc = fork();
10
        if (rc < 0) { // fork failed; exit
11
            fprintf(stderr, "fork failed\n");
12
            exit(1);
13
        } else if (rc == 0) { // child (new process)
14
            printf("hello, I am child (pid:%d) \n", (int) getpid());
15
        } else {
                              // parent goes down this path (main)
16
            int wc = wait(NULL);
17
            printf("hello, I am parent of %d (wc:%d) (pid:%d) \n",
18
                     rc, wc, (int) getpid());
19
20
        return 0;
21
22
    }
                Figure 5.2: Calling fork() And wait() (p2.c)
```

What happens during exec?

- After fork, parent and child are running same code

 Not too useful!
- A process can run exec() to load another executable to its memory image
 - So, a child can run a different program from parent
- Variants of exec(), e.g., to pass commandline arguments to new executable



```
#include <stdio.h>
1
    #include <stdlib.h>
2
   #include <unistd.h>
3
    #include <string.h>
4
    #include <sys/wait.h>
5
6
    int
7
    main(int argc, char *argv[])
8
9
        printf("hello world (pid:%d)\n", (int) getpid());
10
        int rc = fork();
11
        if (rc < 0) {
                              // fork failed; exit
12
            fprintf(stderr, "fork failed\n");
13
            exit(1);
14
        } else if (rc == 0) { // child (new process)
15
            printf("hello, I am child (pid:%d)\n", (int) getpid());
16
            char *myargs[3];
17
            myargs[0] = strdup("wc"); // program: "wc" (word count)
18
            myargs[1] = strdup("p3.c"); // argument: file to count
19
            myargs[2] = NULL;
                                         // marks end of array
20
            execvp(myargs[0], myargs); // runs word count
21
           printf("this shouldn't print out");
22
                               // parent goes down this path (main)
        } else {
23
            int wc = wait(NULL);
24
            printf("hello, I am parent of %d (wc:%d) (pid:%d) \n",
25
                    rc, wc, (int) getpid());
26
27
        return 0;
28
29
```

Figure 5.3: Calling fork(), wait(), And exec() (p3.c)

Case study: How does a shell work?

- In a basic OS, the init process is created after initialization of hardware
- The init process spawns a shell like bash
- Shell reads user command, forks a child, execs the command executable, waits for it to finish, and reads next command
- Common commands like ls are all executables that are simply exec'ed by the shell

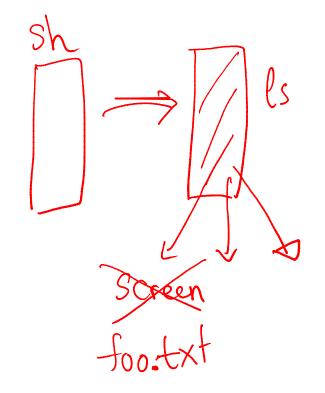
prompt>ls

a.txt b.txt c.txt

Fork exec (ls)

More funky things about the shell

- Shell can manipulate the child in strange ways
- Suppose you want to redirect output from a command to a file
- prompt>ls > foo.txt
- Shell spawns a child, rewires its standard output to a file, then calls exec on the child



```
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <unistd.h>
4 #include <string.h>
5 #include <fcntl.h>
  #include <sys/wait.h>
6
7
   int
8
    main(int argc, char *argv[])
9
10
    {
11
        int rc = fork();
        if (rc < 0) {
                               // fork failed; exit
12
            fprintf(stderr, "fork failed\n");
13
            exit(1);
14
        } else if (rc == 0) { // child: redirect standard output to a file
15
            close (STDOUT_FILENO);
16
            open("./p4.output", O_CREAT|O_WRONLY|O_TRUNC, S_IRWXU);
17
                                                                                  ychild
18
            // now exec "wc"...
19
            char *myargs[3];
20
            myargs[0] = strdup("wc"); // program: "wc" (word count)
21
            myargs[1] = strdup("p4.c"); // argument: file to count
22
            myargs[2] = NULL;
                                        // marks end of array
23
            execvp(myargs[0], myargs); // runs word count
24
                              // parent goes down this path (main)
        } else {
25
            int wc = wait (NULL);
26
27
        return 0;
28
29
   ł
```

Figure 5.4: All Of The Above With Redirection (p4.c)

Here is the output of running the p4.c program:

```
prompt> ./p4
prompt> cat p4.output
32 109 846 p4.c
prompt>
```