

## SUPPLEMENT MATERIAL

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### 1. PERFORMANCE BOUNDS FOR Q1 AND Q2

**Theorem 1:** - Bounds for Q1

Abiding by the notations in [1], let  $B(\cdot)$  be an  $(\delta, k)$  RIP obeying linear operator [1], i.e.  $\forall X \neq 0$  and  $\|X\|_0 \leq k$

$$\left| \frac{\|B(X)\|_2^2}{\|X\|_2^2} - 1 \right| < \delta \quad (1)$$

Let  $X^*$  be the true rank 1 matrix that satisfies the constraints in Q1, and  $\tilde{X}$  be the solution to Q1 for an appropriate choice of parameters. We define  $\Delta = \tilde{X} - X^*$ . For a matrix  $X \in \mathbb{R}^{d \times d}$ ,  $X_T$  denotes a matrix with all values zero except the indices in  $T$ , which are set to the corresponding values of  $X$ .  $\forall \lambda \geq \frac{2d}{1-\rho} + \frac{d}{k^{\frac{1}{2}}}$  and  $\delta \leq \sqrt{2} - 1$  we get

$$\|\Delta\|_2 \leq \frac{2\alpha\varepsilon(1-\rho)^{-1} + 2(1+\rho)(1-\rho)^{-1}k^{-\frac{1}{2}}\|\tilde{X} - X_k^*\|_1}{1 - \left(\frac{2d}{1-\rho} + \frac{d}{k^{\frac{1}{2}}}\right)\frac{1}{\lambda}} \quad (2)$$

where  $\rho = \frac{\sqrt{2}\delta}{1-\delta} \leq 1$ ,  $X_k^*$  is the matrix with the  $k$  largest elements in  $X^*$  at the corresponding indices and the others elements to be 0,  $\varepsilon \sim \mathcal{O}(\sqrt{n})$  and  $\alpha = \frac{4\sqrt{(I+c)(1+\delta)}}{1-\delta}$ .  $I$  is the true intensity of the underlying measurement, i.e.  $\|B(X^*)\|_1$ , which is assumed to be known for naturally acquired measurements.

**Theorem 2:** - Bounds for Q2

Given that the assumptions in section 2.2 hold, and  $\Psi$  has sufficiently small RIP constant  $\delta$ , then, there exist positive absolute constants  $C_1, C_2$ , and  $C_3$  such that if  $m \geq c_1 k \log \frac{d}{k}$ , and  $n \geq C_1 m$  then any estimate  $\hat{X}$  of the Algorithm obeys

$$\left\| \hat{X} - X^* \right\|_F \leq C_2,$$

for all rank-one and  $k \times k$ -sparse matrices  $X^* \succcurlyeq 0$  with probability exceeding  $1 - e^{-C_3 n}$ . The constant  $C_2$  depends on  $\sqrt{I}$  and since  $\varepsilon \sim \mathcal{O}(\sqrt{n})$ , we don't get a dependence on  $\frac{1}{\sqrt{n}}$  like that in [2]. The proofs of these theorems are inspired by the performance bounds for compressive phase retrieval presented in [3] by Ohlsson et. al for the standard CPRL problem (P1), and in [2] by Bahamani et. al for the 2-stage sparse recovery method (P2). We adapted these proofs for Poisson corrupted signals using the techniques in [4] and [5].

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### 2. PROVING PERFORMANCE BOUNDS

We briefly present here the proof outline for performance bounds of Q1 and Q2.

#### 2.1. Proof for Q1

As defined in theorem 1, let  $B(\cdot)$  be an  $(\delta, k)$  RIP obeying linear operator [1], i.e.  $\forall X \neq 0$  and  $\|X\|_0 \leq k$

$$\left| \frac{\|B(X)\|_2^2}{\|X\|_2^2} - 1 \right| < \delta \quad (3)$$

Let  $X^*$  be the true rank 1 matrix that satisfies the constraints in Q1, and  $\tilde{X}$  be the solution to Q1 for an appropriate choice of parameters. We define  $\Delta = \tilde{X} - X^*$ . For a matrix  $X \in \mathbb{R}^{d \times d}$ ,  $X_T$  denotes a matrix with all values zero except the indices in  $T$ , which are set to the corresponding values of  $X$ .  $T_0$  is the index set of the  $k$  largest elements of  $X^*$  in absolute value, and  $T_0^c$  is the corresponding complement set. Let  $T_1$  be the index set associated with the  $k$  largest elements in absolute value of  $\Delta_{T_0^c}$  and  $T_{0,1} \triangleq T_0 \cup T_1$  be the union. Let  $T_2$  be the index set associated with the  $k$  largest elements in absolute value of  $\Delta_{T_{0,1}^c}$ , and so on. Using the fact that -:

$$\|\Delta\|_2 = \|\Delta_{T_{0,1}} + \Delta_{T_{0,1}^c}\|_2 \leq \|\Delta_{T_{0,1}}\|_2 + \|\Delta_{T_{0,1}^c}\|_2. \quad (4)$$

We independently derive bounds for both the terms on the RHS, and then combine the result to get a bound on  $\|\Delta\|_2$ .

Bounds for  $\|\Delta_{T_{0,1}^c}\|_2$  can be derived exactly as in the proof of theorem 8 in [1]. For bounding  $\|\Delta_{T_{0,1}}\|_2$ , we first use a technique from [5] to upper bound  $\|B(\Delta)\|_2$ , which handles Poisson corruption for signals in Q1 and Q2. Some algebraic manipulation gives

$$\begin{aligned} \|B(\Delta)\|_2^2 &= \|B(\tilde{X} - X^*)\|_2^2 \\ &= \sum_{i=1}^N ((\sqrt{(B(\tilde{X}))_i + c} - \sqrt{(B(X^*))_i + c})^2 \\ &\quad (\sqrt{(B(\tilde{X}))_i + c} + \sqrt{(B(X^*))_i + c})^2). \end{aligned}$$

1. Since  $\|\sqrt{y+c} - \sqrt{B(X)+c}\|_2$  is upper bounded by  $\varepsilon$ , using triangular inequality we get  $\|\sqrt{B(\tilde{X})+c} - \sqrt{B(X^*)+c}\|_2 \leq \|\sqrt{y+c} - \sqrt{B(\tilde{X})+c}\|_2 + \|\sqrt{y+c} - \sqrt{B(X^*)+c}\|_2 \leq 2\varepsilon$ .

2. For scalars  $v_1 \geq 0, v_2 \geq 0$ , we have  $(\sqrt{v_1} + \sqrt{v_2})^2 \leq 4\max(v_1, v_2)$ . Likewise we also have  $(B\tilde{X})_i \leq I$ , where  $I$  is the true intensity of the underlying measurement, i.e.  $\|B(X^*)\|_1$ , which is assumed to be known for naturally acquired measurements. Hence we get  $(\sqrt{(B(\tilde{X}))_i + c} + \sqrt{(B(X^*))_i + c})^2 \leq 4(I + c)$ .
3. Using the results of all the steps above, we have  $\|B(\Delta)\|_2 \leq 4\epsilon\sqrt{I + c}$ .

By a derivation similar to that in [4], one can show that

$$\|\Delta_{T_{0,1}}\|_2 \leq \frac{\alpha\epsilon}{1-\rho} + \frac{2\rho}{1-\rho} k^{-1/2} \|\tilde{X} - X^*_k\|_1 - \frac{1}{\lambda} \frac{1}{1-\rho} \text{Tr } \Delta \quad (5)$$

where  $\rho = \frac{\sqrt{2\delta}}{1-\delta} \leq 1$ ,  $X^*_k$  is the matrix with the  $k$  largest elements in  $X^*$  at the corresponding indices and the others elements to be 0, and  $\alpha = \frac{4\sqrt{(I+c)(1+\delta)}}{1-\delta}$ . The  $\text{Tr } \Delta$  term makes the bound a bit unpractical, and can be eliminated using the inequality -  $\text{Tr } \Delta \leq d\|\Delta\|_2$ . This inequality can be obtained by upper bounding trace by  $l_1$  norm and then using the Cauchy Schwartz inequality subsequently as done in [1]. We observe that this inequality differs slightly from that in the proof of the corollary 9 in [1] since  $\Delta$  is not  $k$ -sparse. Using this inequality, and the bounds obtained for  $\|\Delta_{T_{0,1}}\|_2$  and  $\|\Delta_{T_{0,1}^c}\|_2$ ,  $\forall \lambda \geq \frac{2d}{1-\rho} + \frac{d}{k^{\frac{1}{2}}}$  and  $\delta \leq \sqrt{2} - 1$  we get

$$\|\Delta\|_2 \leq \frac{2\alpha\epsilon(1-\rho)^{-1} + 2(1+\rho)(1-\rho)^{-1}k^{-\frac{1}{2}}\|\tilde{X} - X^*_k\|_1}{1 - (\frac{2d}{1-\rho} + \frac{d}{k^{\frac{1}{2}}})\frac{1}{\lambda}} \quad (6)$$

## 2.2. Proof for Q2

First we find an upper bound on  $\|\hat{A} - A^*\|_2$ .

$\hat{A}$  is the solution to Q2 and  $A^*$  is the true low rank matrix. The performance bounds for estimation of  $A$  can be derived exactly as done in [2], with a modification to the proof of theorem 1.2 presented [6] to account for the VST based constraint to handle Poisson noise. In section 6 of [6], on taking into account our constraint, and following the steps to bound  $\|B(\Delta)\|_2$  as done for Q1 above, equation 6.2 in section 6 of [6] changes to  $\|\mathcal{A}(H)\|_2 \leq 4\epsilon\sqrt{I + c}$ , where  $H = \Delta$  and  $\mathcal{A}$  corresponds to the linear operator  $B(\cdot)$  in Q1. Since the upper bound for  $\|\mathcal{A}(H)\|_2$  is changed only by a multiplicative constant, the rest of the proof proceeds in a similar manner to [6], giving us a bound of the form similar to that in [2] with a different constant on the R.H.S, i.e.

$$\|\hat{A} - A^*\|_F \leq \frac{C_2''\epsilon}{\sqrt{n}}, \quad (7)$$

for all valid  $A^*$ , thereby for all valid  $X^*$ , with probability exceeding  $1 - e^{-C_3n}$  for some constant  $C_3$ .  $C_2''$  is a constant that depends on  $\sqrt{I + c}$  since the  $\epsilon$  gets scaled by that factor

due to the VST constraint. We see that the bound we obtained has a dependence on  $\sqrt{n}$  as opposed to the theorem 1.2 in [6]. This difference arises due to the fact that the noise vector in [6] is assumed to be a sphere of radius  $\sqrt{n}$ , and hence we see a dependence on  $\sqrt{n}$  in the 2-stage sparse recovery algorithm as no such assumption is taken for the noise vector in this case. The  $\hat{A}$  estimated from above step now can be used to get the sparse estimate  $\hat{X}$ . The upper bound on  $\|\hat{X} - X^*\|_2$  can now be derived exactly as done in [2]. This completes the proof sketch for Q1 and Q2.

## 3. REFERENCES

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