CS310 : Automata Theory 2019 IITB, India

Tutorial sheet 5 Context-free Grammar and Pushdown automata

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Compile date: 2019-02-18

- 1. Give a context-free grammar(CFG) for the following languages.
 - (a) all binary words with both an even number of zeroes and an even number of ones
 - (b) $\{0^n 1^\ell 0^m | n+m=\ell\}$
 - (c) $\{a^m b^n \mid m \ge n\}$
 - (d) $\{a^m b^n \mid m \neq n\}$
 - (e) $\{a^m b^n \mid n > m\}$
 - (f) $\{a^m b^n \mid m < 2n < 3m\}$
 - (g) All propositional logic formulas in conjunctive normal form, where variables are denoted by x followed by a binary number.

Also give unambiguous grammars, if possible.

- 2. For each of the above languages, give left most derivations of some word of length longer than 5.
- 3. Prove that any grammar G = (N, T, P, S) that has only the following types of rules generates a regular language : $A \to aB$, $A \to a$, or $A \to \epsilon$, where $A, B \in N$ and $a \in T$.
- 4. Let $P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$ be an PDA accepting language L. Show that there exists another NPDA $P' = (Q', \Sigma, \Gamma', \delta', q'_0, Z_0, F')$ accepting the same language L such that $|\Gamma'| = 2$. Can we compare the size of P and P'?
- 5. Let us consider a derivation $S \stackrel{*}{\Rightarrow} w$ in grammar G = (N, T, P, S) such that no word is repeated in the derivation. Can we bound the length of the derivations in terms of |G| and |w|? If not, give a counterexample.
- 6. Let P be a PDA with empty-stack language $L = L^{\epsilon}(P)$, and suppose that ϵ is not in L. Modify P such that it accepts $L \cup \{\epsilon\}$ by empty stack.
- 7. Solve Exercise 6.2.5 of the text book.