

# CS228 Logic for Computer Science 2020

## Lecture 13: Using SAT solver

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# Topic 13.1

## Z3 solver

# Solver basic interface

- ▶ Input : formula
- ▶ Output: sat/unsat

If satisfiable, we may ask for a satisfying assignment.

## Exercise 13.1

*What can we ask from a solver in case of unsatisfiability?*

## Z3: SMT solver

- ▶ Written in C++
- ▶ Provides API in C++ and Python

# Locally Installing Z3 (Linux)

- ▶ Download

```
https://github.com/Z3Prover/z3/releases/download/z3-{4.7.1}/z3-4.7.1-x64-ubuntu-16.04.zip
```

- ▶ Unzip the file in some folder. Say

```
/path/z3-4.7.1-x64-ubuntu-16.04/
```

- ▶ Update the following environment variables

```
$export LD_LIBRARY_PATH=$LD_LIBRARY_PATH:/path/z3-4.7.1-x64-ubuntu-16.04/bin
```

- ▶ After the setup the following call should throw no error

```
$python3 /path/z3-4.7.1-x64-ubuntu-16.04/bin/python/example.py
```

## Steps of using Z3 via the python interface

```
from z3 import *           # load z3 library

p1 = Bool("p1")            # declare a Boolean variable
p2 = Bool("p2")
phi = Or( p1, p2 )         # construct the formula

print(phi)                 # printing the formula

s = Solver()               # allocate solver
s.add( phi )               # add formula to the solver
r = s.check()              # check satisfiability
if r == sat:
    print("sat")
else:
    print("unsat")

# save the script test.py
```

Commentary: In python 3.0 onward one may have to write `print("sat")` instead of `print "sat"`

## Get a model/assignment

In z3, assignments are called models

```
r = s.check()
if r == sat:
    m = s.model()           # read assignment/model
    print(m)               # print model
else:
    print("unsat")
```

## Solve and print model

```
from z3 import *

# packaging solving and model printing
def solve( phi ):
    s = Solver()
    s.add( phi )
    r = s.check()
    if r == sat:
        m = s.model()
        print(m)
    else:
        print("unsat")

# we will use this function in later slides
```

# Pointer and variable

There is a distinction between

the Python variable name and the propositional variable it holds.

```
from z3 import *      # load z3 library

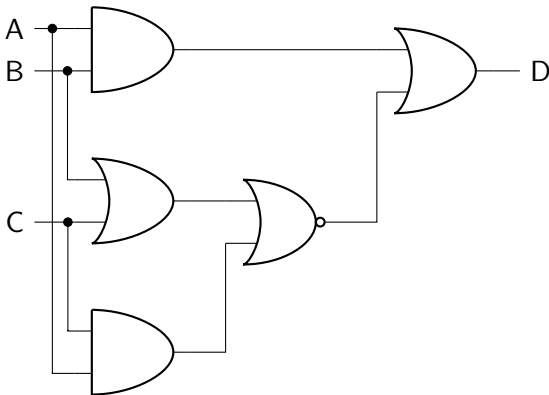
x = Bool("y") #creates Propositional variable y

z = x  # python pointer z also holds variable y
```

## Exercise: encoding Boolean circuit

### Exercise 13.2

Using Z3, find the input values of  $A$ ,  $B$ , and  $C$  such that output  $D$  is 1.



We know you can do it! Please do not shout the answer. Please make counter find it.

## Design of solvers: context vs. solver

Any complex software usually has a context object.

The context consists of a **formula store** containing the constructed formulas.

Z3 Python interface **instantiates a default context**. Therefore, we do not see it explicitly.

A **Solver** is a solving instance. There can be **multiple solvers** in a context.

The **Solver** solves only the added formula.

# Attendance quiz

Which of the following hold about the Z3 API?

`x = Bool("x")` creates Boolean variable `x`  
`phi = Or( p1, p2 )` creates a conjunction of two formulas  
`phi = And( p1, p2 )` creates a disjunction of two formulas  
`s = Solver()` create an instantiation of a Z3 solver  
For solver `s`, `s.add(F)` adds formula `F` in solver `s`  
For solver `s`, `s.check()` runs CDCL on the conjunction of formulas added to `s` so far  
For solver `s`, `s.model()` returns an assignment if `s.check()` has returned "sat"  
For solver `s`, `s.model()` throws an exception if `s.check()` has returned "unsat"  
For solver `s`, `s.model()` throws an exception if `s.check()` has not been called  
`y = Bool("x")` creates Boolean variable `y`  
`phi = Or( p1, p2 )` creates a disjunction of two formulas  
`phi = And( p1, p2 )` creates a conjunction of two formula  
`s = Solver()` create an instantiation of a shopping mall  
For solver `s`, `s.add(F)` removes formula `F` in solver `s`  
For solver `s`, `s.check()` runs DPLL on the conjunction of formulas added to `s` so far  
For solver `s`, `s.model()` returns an assignment if `s.check()` has returned "unsat"  
For solver `s`, `s.model()` throws an exception if `s.check()` has returned "sat"  
For solver `s`, `s.model()` returns an assignment if `s.check()` has not been called

## Topic 13.2

### Problems

## Exercise : Python programming

### Exercise 13.3

Write a Python program that generates a random graph in a file `edges.txt` for  $n$  nodes and  $m$  edges, which are given as command line options.

Please store edges in `edges.txt` as the following sequence of tuples

10,12

30,50

...

### Exercise 13.4

Write a program that reads a directed graph from `edges.txt` and finds the number of **strongly connected components** in the graph

### Exercise 13.5

Write a program that reads a directed graph from `edges.txt` and finds the cliques of size  $k$ , which is given as a command line option.

# Proving theorems

## Exercise 13.6

*Prove/disprove the following theorems*

- ▶ *Sky is blue. Space is black. Therefore sky and space are blue or black.*
- ▶ *Hammer and chainsaw are professional tools. Professional tools and vehicles are rugged. Therefore, hammers are rugged.*

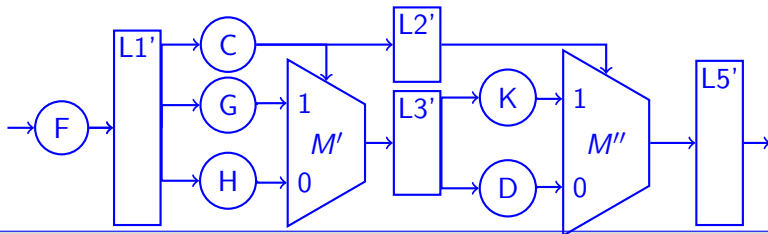
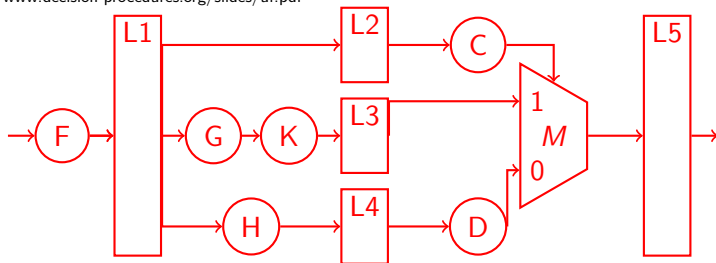
# Exercise: translation validation

## Exercise 13.7

Show that the following two circuits are equivalent.

*Ls are latches, circles are Boolean circuits, and Ms are multiplexers.*

Source: <http://www.decision-procedures.org/slides/uf.pdf>



## Write a function: find positive variables

### Exercise 13.8

*Find the set of Boolean variables that occur only positively in a propositional logic formula.*

*An occurrence of a variable is positive if there are even number of negations from the occurrence to the root of the formula.*

*Examples:*

*Only  $q$  occurs positively in  $p \wedge \neg(\neg q \wedge p)$ .*

*$p$  occurs positively in  $\neg\neg p$ .*

*$p$  does not occur positively in  $\neg p$ .*

*$p$  and  $q$  occur positively in  $(p \vee \neg r) \wedge (r \vee q)$ .*

# Write a function: find unrelated constraints

## Exercise 13.9

*Consider a CNF formula  $F$ . Find the subsets of  $F$  that have disjoint set of symbols.*

*Examples:*

*$(x \vee y) \wedge (x \vee z) \wedge u$  has two unrelated subsets  $\{(x \vee y), (x \vee z)\}$  and  $\{u\}$*

Write a function: find maximum occurring symbol

### Exercise 13.10

*Consider a formula  $F$ . Find the variable in  $F$  that occurs most often.*

*Example:*

*$x$  occurs most often in  $(x \vee \neg x) \wedge (y \vee \neg x)$ .*

Write a function: detect CNF or NNF

### Exercise 13.11

*Consider a formula  $F$ . Write a function that detects if  $F$  is CNF.*

### Exercise 13.12

*Consider a formula  $F$ . Write a function that detects if  $F$  is NNF.*

# End of Lecture 13