# CS 433 Automated Reasoning 2024

## Lecture 5: Encoding into reasoning problems

Instructor: Ashutosh Gupta

IITB India

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Topic 5.1

# Z3 solver

# Solver basic interface

- Input : formula
- Output: sat/unsat

If satisfiable, we may ask for a satisfying assignment.

Exercise 5.1

*What can we ask from a solver in case of unsatisfiability?*

# Z3: SMT solver

- Written in C++
- Provides API in C++ and Python
- We will initially use python interface for quick ramp up
- Later classes we will switch to C++ interface

# Installing Z3 (Ubuntu-22.04)

```
$sudo apt install z3 python3-z3
```

Not tested on 20.04

**Commentary:** You may also try $pip3 install python3-z3

# Locally Installing a version of Z3 (Linux)

Let us install `z3-4.7.1`. You may choose another version.

- Download

  `https://github.com/Z3Prover/z3/releases/download/z3-4.7.1/z3-4.7.1-x64-ubuntu-16.04.zip`

- Unzip the file in some folder. Say

  /path/z3-4.7.1-x64-ubuntu-16.04/

- Update the following environment variables

```
$export LD_LIBRARY_PATH=$LD_LIBRARY_PATH:/path/z3-4.7.1-x64-ubuntu-16.04/bin
$export PYTHONPATH=$PYTHONPATH:/path/z3-4.7.1-x64-ubuntu-16.04/bin/python
```

- After the setup the following call should throw no error

  $python3 /path/z3-4.7.1-x64-ubuntu-16.04/bin/python/example.py

Topic 5.2

# Using solver

# Steps of using Z3 via python interface

```
from z3 import *          # load z3 library

p1 = Bool("p1")           # declare a Boolean variable
p2 = Bool("p2")

phi = Or( p1, p2 )        # construct the formula
print(phi)                # printing the formula

s = Solver()              # allocate solver
s.add( phi )              # add formula to the solver
r = s.check()             # check satisfiability
if r == sat:
    print("sat")
else:
    print("unsat")        # save the script test.py
                          # run \$python3 test.py
```

# Get a model

```
r = s.check()
if r == sat:
    m = s.model()      # read model
    print(m)           # print model
else:
    print("unsat")
```

## Exercise 5.2

What happens if we run `m = s.model()` in the unsat case?

# Solve and print model

```
from z3 import *

# packaging solving and model printing
def solve( phi ):
  s = Solver()
  s.add( phi )
  r = s.check()
  if r == sat:
      m = s.model()
      print(m)
  else:
      print("unsat")

 # we will use this function in later slides
```

# Pointer and variable

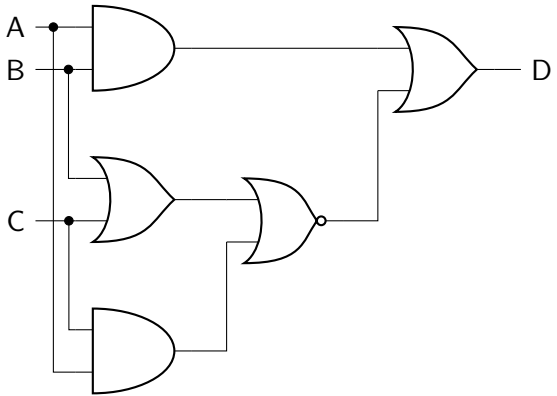
There is a distinction between the Python variable name and the propositional variable it holds.

```
from z3 import *      # load z3 library

x = Bool("y")         # creates Propositional variable y

z = x                 # python pointer z also holds variable y
```

# Exercise: encoding Boolean circuit

## Exercise 5.3

*Using Z3, find the input values of A, B, and C such that output D is 1.*




We know you can do it! Please do not shout the answer. Please make computer find it.

Topic 5.3

# Solver engineering

# Design of solvers: context vs. solver

Any complex software usually has a context object.

The context consists of a formula store containing the constructed formulas.

Z3 Python interface instantiates a default context. Therefore, we do not see it explicitly.

A `Solver` is a solving instance. There can be multiple solvers in a context.

The `Solver` solves only the added formula.

# Formula handling

```
a = Bool('a')
b = Bool('b')
ab = And( a, b )

# accessing sub-formulas
print(ab.arg(0))
print(ab.arg(1))

# accessing the symbol at the head
ab_decl = ab.decl()
name = ab_decl.name()
if name == "and":
    print("Found an And")
```

# Topic 5.4

## Theory formulas

# Solving rational(real) arithmetic

```
x = Real('x')
y = Real('y')
phi = And(x + y > 5, x > 1, y > 1)
solve( phi )
```

For linear arithmetic
Real == rational

# Solving integer arithmetic

```
x = Int('x')
y = Int('y')
phi = And(x + y > 5, x > 1, y > 1)
solve( phi )
```

# Exercise: bounded model checking

## Exercise 5.4

*Using Z3, find the inputs x and y such that the assert fails.*

```
int foo( int x, int y ) {
  int z = 3*x + 2*y - 3;
  if( y > 0 )
    assert( z != 0 );
}
```

# Solving bit precise

```
x = BitVec('x',32) # declare name and bit length
y = BitVec('y',32)
phi = And(x + y > 5, x > 1, y > 1)
solve( phi )
```

- Bit lengths must match in an operation
- Largely solved by bit blasting
- Far more expensive to solve!

converting Bit-vector formulas into Boolean formulas by replacing vectors by bits and operation by circuits.

# Exercise : observe overflow behavior

Exercise 5.5

*Give a bit-vector formula that is satisfiable due to overflow of addition, but in infinite precision it is unsatisfiable.*

# Uninterpreted functions

```
x = Int('x')
y = Int('y')

# declaring Int -> Int function
h =  Function('h', IntSort(), IntSort() )

phi = And( h( x ) > 5, h( y ) < 2 )

solve( phi )
```

# Exercise:

## Exercise 5.6

*Give a satisfying model of the following formula*

$$g(x, y) < 0 \wedge g(y, x) > 0 \wedge y = x$$

# Uninterpreted sorts

```
u = DeclareSort('U') # declaring new sort

c =  Const('c', u )  # declaring a constant of the sort

f =  Function('f', u, u ) # declaring a function of the sort

# declaring a predicate of the sort
P =  Function('P', u, BoolSort() )

phi = And( f(c) == c,  P( f(c) ),  Not( P(c) ) )

solve( phi )
```

## Exercise 5.7

*Get model after dropping the third atom. Interpret the model.*

**Commentary:** Hint: the solver also chooses domains for the uninterpreted sorts, and the models of the functions are presented in terms of the domains.

Topic 5.5

# Quantified formulas

# Quantifiers

```
u = DeclareSort('U')
H = Function('Human', u, BoolSort() )
M = Function('Mortal', u, BoolSort() )

# Humans are mortals
x = Const('x', u )
all_mort = ForAll( x, Implies( H(x), M(x) ) )

s = Const('Socrates', u )
thm = Implies( And( H(s), all_mort ), M(s) )

solve( Not(thm) )  # negation of a valid theorem
                   # is unsatisfiable
```

# Exercise: solving quantified formulas

Exercise 5.8

*Prove/disprove if the following statement is valid.*

*There is someone such that if the one drinks, then everyone drinks*

Exercise 5.9

*Write a formula that only accepts infinite models. Encode the formula in Z3 and get model.*

# Quantified formula handling

```
u = DeclareSort('U')
H = Function('Human', u, BoolSort() )
M = Function('Mortal', u, BoolSort() )
x = Const( 'x', u )
y = Const( 'y', u )

all_mort = ForAll( x, Implies( H(x), M(x) ) )
print(all_mort.body())
# Output: Implies(Human(Var(0)), Mortal(Var(0)))
# Var(0) is FOL variable
# Naming quantified variables using DeBruijn index

alt = ForAll( x, Exists( y, Implies( H(x), M(y) ) ))
print(alt.body().body())
# Output: Implies(Human(Var(1)), Mortal(Var(0)))
```

Topic 5.6

# SMT2 format

# API vs Input language

- Each solver has their own API
- We need a common input format for
  - interoperability and
  - database of problems

# Standard format for SMT solvers

SMT2 is a standard input format for SMT solvers.

`http://smtlib.cs.uiowa.edu/language.shtml`

- Formulas are written in prefix notation (Why?)

```
(>= (* 2 x) (+ y z))
```

- There is a simple type system. Similar to Z3 API.
- Solver interacts like a stack

# File format

An SMT2 file has five parts

1. Preamble declarations
2. Sort declarations
3. Variable declarations
4. Asserting formulas
5. Solving commands

# Preamble declaration

- Set configurations of the solvers

```
(set-logic QF_UFLIA)                ;setting Theory/Logic

(set-option :produce-proofs true)   ;enable proof generation if input is unsat
```

# Sort declarations

- Declare new sorts of the variables

```
(declare-sort symbol numeral)
```

## Example 5.1

```
(declare-sort U 0)     ; new sort with no parameters
(declare-sort Arr 2)   ; new sort with two parameters
```

# Variable declarations

- Declare variables and functions that may be used in the formulas

```
(declare-fun symbol (sort*) sort)
```

## Example 5.2

```
(declare-fun x () Int)                  ;declare variable
(declare-fun f (Int) Int)               ;declare a function with one argument
(declare-fun g (Int Int) Int)           ;declare a function with two arguments
(declare-fun h ((Arr U Int) Int) Int);declare a function with two argument
```

Sorts with parameters

# Asserting formulas

- Formulas are asserted in a sequence

## Example 5.3

```
(assert (>= (* 2 x) (+ y z)))
(assert (< (f x) (g x x)))
(assert (> (f y) (g x x)))
```

# Commands

- Commands are the actions that solver needs to do

Example 5.4

```
(check-sat) ; checks if the conjunction of asserted formula is sat
(get-model) ; returns a model if the formulas are sat
```

# Stack interaction

The standard is designed to be interactive

- Asserted formulas are pushed in the stack of the solver
- `(push)` command places marker on the stack
- `(pop)` removes the formulas upto the last marker

## Example 5.5

```
(push)
(assert (= x y))
(check-sat)
(pop)
```

*After the pop the solver state goes back to the last push. Useful in interactive use of solver.*

# Full example

```
(set-logic QF_UFLIA)
(set-option :produce-proofs true)
(declare-fun x () Int)
(declare-fun y () Int)
(declare-fun z () Int)
(declare-fun f (Int) Int)
(declare-fun g (Int Int) Int)
(assert (>= (* 2 x) (+ y z)))
(assert (< (f x) (g x x)))
(assert (> (f y) (g x x)))
(check-sat)
(get-model)
(push)
(assert (= x y))
(check-sat)
(pop)
(exit)
```

# Demo

http://rise4fun.com/z3

Topic 5.7

# Problems

# Exercise : Python programming

## Exercise 5.10

Write a Python program that generates a random graph in a file `edges.txt` for $n$ nodes and $m$ edges, which are given as command line options.
Please store edges in `edges.txt` as the following sequence of tuples

```
10,12
30,50
....
```

## Exercise 5.11

Write a program that reads a directed graph from `edges.txt` and finds the number of strongly connected components in the graph

## Exercise 5.12

Write a program that reads a directed graph from `edges.txt` and finds the cliques of size $k$, which is given as a command line option.

# Proving theorems

## Exercise 5.13

*Prove/disprove the following theorems using a solver*

- *Sky is blue. Space is black. Therefore sky and space are blue or black.*
- *Hammer and chainsaw are professional tools. Professional tools and vehicles are rugged. Therefore, hammers are rugged.*

# Write a function: find positive variables

## Exercise 5.14

*Find the set of Boolean variables that occur only positively in a propositional logic formula.*

*An occurrence of a variable is positive if there are even number of negations from the occurrence to the root of the formula.*

*Examples:*

*Only $q$ occurs positively in $p \wedge \neg(\neg q \wedge p)$.*

*$p$ occurs positively in $\neg\neg p$.*

*$p$ does not occur positively in $\neg p$.*

*$p$ and $q$ occur positively in $(p \vee \neg r) \wedge (r \vee q)$.*

# Write a function: compute linear coefficient

Exercise 5.15

*Find coefficient of each variable in a linear term. If the term is non-linear, throw an exception.*

*Examples:*

$x - 2x + y + 4$ *should return* $[4, -1, 1]$ *if variables are ordered* $[x, y]$.

$x - x + 4y - 2(2y)$ *should return* $[0, 0, 0]$ *if variables are ordered* $[x, y]$.

$(x + 1) * y$ *should throw an exception*

# Write a function: find quantifier alternation depth

## Exercise 5.16

*Compute quantifier alternations depth of a sentence.*

*Maximum number of quantifier type switches is any path from an atom to the root.*

*Examples: quantifier alternations depth of*
$\forall x.\ \exists y.\ E(x, y)$ *is 1.*

$\forall x.\ \exists y.\ \forall z.\ E(x, y, z)$ *is 2.*

$\forall x.\ \forall y.\ E(x, y)$ *is 0.*

$\forall x.\ (\ (\ \exists y.\ H(x, y)\ ) \Rightarrow G(x))$ *is 0.* *($\exists$ under negation is $\forall$ and vice-versa)*

$\forall x.\ (\ (\ \exists y.\ H(x, y)\ ) \Rightarrow \exists z.\ G(x, z))$ *is 1.*

# Write a function: find unrelated constraints

## Exercise 5.17

*Consider a formula $F$ consists of only a conjunction of atoms. Find the partitions of $F$ that have disjoint set of uninterpreted symbols.*

*Examples:*

$x = y \wedge x = z \wedge P(u)$ *has two unrelated subsets* $\{x = y, x = z\}$ *and* $\{P(u)\}$

$x + y = 3 \wedge z + u \geq 10$ *has two unrelated subsets* $\{x + y = 3\}$ *and* $\{z + u \geq 10\}$*, while they have a common interpreted symbol* $+$*.*

# Write a function: find maximum occurring symbol

Exercise 5.18

*Consider a formula $F$. Find the uninterpreted symbol in $F$ that occurs most often.*

*Examples:*

*$x$ occurs most often in $g(g(x,x),g(x,x))$.*

*$f$ occurs most often in $f(x,y) = f(x,b) \wedge f(2,3) > 10$.*

*$D$ occurs most often in $\exists x.(D(x) \Rightarrow D(x+1))$.* *quantified variables are not counted.*

# Write a function: find common symbols

Exercise 5.19

*Consider formulas $F_1$ and $F_2$. Find uninterpreted symbols that occur both in $F_1$ and $F_2$.*

*Examples:*

$\{x, f\}$ *occurs* $f(x) > 3$ *and* $f(y) < x$, *but not* $y$

$\{f\}$ *occurs* $f(x) > 3$ *and* $\forall x.f(x) > y$, *but not* $x$ *and* $y$.

$\{f, x\}$ *occurs* $f(x) > 3$ *and* $x > 20 \vee \forall x.f(x) > y$. *quantified variables are not counted.*

# Integer vs. Reals

Exercise 5.20

*Consider the following constraints*

$$3x - y \geq 2 \wedge 3y - z \geq 3 \wedge 3 \geq x + y$$

*Solve the above constraints using SMT solver under the following theories*

- *Reals (QF_LRA)*
- *Int (QF_LIA)*

# End of Lecture 5