CS228 Logic for Computer Science 2022

Lecture 18: FOL Resolution

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Topic 18.1

Refutation proof systems

Recall: derivations starting from CNF

We have a set of formulas in the lhs, which is viewed as the conjunction of the formulas.

$$\Sigma \vdash F$$

The conjunction of CNF formulas is also a CNF formula.

If all formulas are in CNF, we may assume Σ as a set of clauses.

Recall: refutation proof system

Let us suppose we are asked to derive $\Sigma \vdash F$.

We assume Σ is finite. We can relax this due to compactness of FOL.

We will convert $\bigwedge \Sigma \land \neg F$ into a set of FOL clauses Σ'.

We apply the a refutation proof method on Σ'.

If we derive $\perp$ clause, $\Sigma \vdash F$ is derivable.

Topic 18.2

Unification and resolution

Applying resolution in FOL

We apply resolution when an atom and its negation are in two clauses.

$$\textsc{Resolution}\frac{F \lor C \quad \neg F \lor D}{C \lor D}$$

A complication: we may have terms in the FOL atoms with variables.

We can make two terms equal by substitutions.

Example 18.1

Consider two clauses $P(x, f(y)) \lor C$ *and* $\neg P(z, z) \lor D$

We may be able to make $P(x, f(y))$ *and* $P(z, z)$ *equal by unification.*

Three issues with unification

Before looking at the proof rules, we need a clear understanding of the following three issues.

1. Did we learn about unifying atoms?
2. Is substitution a valid operation for derivations?
3. How do we handle variables across clauses?

Issue 1: unification of atoms

We can lift the idea of unifying terms to atoms.

Simply, treat a predicate as a function.

Example 18.2

Consider atoms $P(x, f(y))$ and $P(z, z)$.

We can unify them using mgu $\sigma = \{x \mapsto f(y), z \mapsto f(y)\}$.

We obtain

- $P(x, f(y))\sigma = P(f(y), f(y))$
- $P(z, z)\sigma = P(f(y), f(y))$

Issue 2: deriving from substitution?

Commentary: The following simultaneous application of ∀-Elim needs proving as derived rule. We presented it as an exercise in the formal proofs lecture.

We know that the following derivation is valid if Σ is a set of sentences.

1. $\Sigma \vdash \forall x, y. F(x, y)$ Premise
2. $\Sigma \vdash F(t_1(x, y), t_2(x, y))$ ∀-Elim
3. $\Sigma \vdash \forall x, y.\ F(t_1(x, y), t_2(x, y))$ ∀-Intro

Therefore the following derivations in our clauses are sound

$$\frac{C}{C\sigma}\sigma \text{ is a substitution.}$$

Example 18.3

The following derivation is a valid derivation

$$\frac{P(x) \lor Q(y)}{P(x) \lor Q(x)}\sigma = \{y \mapsto x\}$$

Issue 3: variables across clauses are not the same

Recall: universal quantifiers distribute over conjunction.

So we can easily distribute the quantifiers and scope only each clauses.

Example 18.4

Consider $\forall w.\forall y.(R(f(w), y) \land \neg R(w, c))$.
After the distribution the formula appears as follows,

$$\forall w.\forall y.R(f(w), y) \land \forall w.\forall y.\neg R(w, c)$$

Therefore, we may view the variables occurring in different clauses as different variables. Even if we use the same name.

Source of confusion. Pay attention!

Topic 18.3

Resolution theorem proving

Resolution theorem proving

Commentary: Please note that the consequences are also statements in the proof system. We simply do not write "$F\ \vdash$" repeatedly, because the left hand side does not change over the course of a resolution proof.

Input: a set of FOL clauses F

Inference rules:

$$\textsc{Assumption}\frac{}{C}\; C \in F$$

$$\textsc{Resolution}\frac{\neg A \lor C \qquad B \lor D}{(C \lor D)\sigma}\; \sigma = mgu(A, B)$$

Example: resolution proof

Example 18.5

Consider statement $\emptyset \vdash (\exists x.\forall y R(x,y) \Rightarrow \forall y.\exists x R(x,y))$.

We translate $\neg(\exists x.\forall y R(x,y) \Rightarrow \forall y.\exists x R(x,y))$ *into the following FOL CNF*

$$R(f(w), y) \land \neg R(w, c)$$

Note that w *and* w *in both clauses are different variables.*

We apply resolution.

$$\textsc{Resolution}\frac{R(f(w), y) \quad \neg R(w, c)}{\bot\sigma}\sigma = \{w \mapsto f(w), y \mapsto c\}$$

Therefore, $\neg(\exists x.\forall y R(x,y) \Rightarrow \forall y.\exists x R(x,y))$ *is unsat, i.e,* $\emptyset \vdash (\exists x.\forall y R(x,y) \Rightarrow \forall y.\exists x R(x,y))$ *holds.*

Example: resolution with unification

Example 18.6

Consider two clauses $P(x, y) \vee Q(y)$ *and* $\neg P(x, x) \vee R(f(x))$.

xs within a clause should be treated as same variable.

If we unify $P(x, y)$ *and* $\neg P(x, x)$, *we obtain most general unifier* $\{x \mapsto x, y \mapsto x\}$.

Therefore,

$$\textsc{Resolution}\frac{P(x, y) \vee Q(y) \qquad \neg P(x, x) \vee R(f(x))}{Q(x) \vee R(f(x))}\sigma = \{x \mapsto x, y \mapsto x\}$$

Commentary: In exams, use $x_1, x_2, \ldots$ to indicated the variables in different clauses. The evaluation will be easier for us and you will be less likely to make mistake.

Why MGU? – not just any unifier

Commentary: Most general consequences allow the maximum opportunity of unifications later, thereby allow us to apply resolution in the maximum possible ways. Therefore, we have have maximum opportunity of finding the empty clause.

MGU keeps maximum generality in the consequence

Example 18.7

We may derive the following, using a σ that is not mgu of $P(x, y)$ and $P(x, x)$.

$$\textsc{Resolution}\frac{P(x,y) \lor Q(y) \qquad \neg P(x,x) \lor R(f(x))}{Q(d) \lor R(f(d))}\sigma = \{x \mapsto d, x \mapsto d, y \mapsto d\}$$

The above conclusion can always be derived from the mgu consequence

$$\cfrac{\cfrac{P(x,y) \lor Q(y) \qquad \neg P(x,x) \lor R(f(x))}{Q(x) \lor R(f(x))}\sigma = \{x \mapsto x, y \mapsto x\}}{Q(d) \lor R(f(d))}\sigma = \{x \mapsto d\}$$

Once a clause becomes specific, we can not go back. Why not keep it general?

Resolution theorem proving : factoring

Commentary: Factoring may appear as a superfluous rule. A bad note in an otherwise beautiful symphony of logic. However, factoring is essential for FOL completeness. It captures the idea that we can express same concept in many possible ways. Sometime, a prover needs to recognize the situation and identify the similarities.

A clause may have copies of facts that can be unified.

We need a rule that allows us to simplify clauses.

$$\textsc{Factor}\frac{L_1 \lor .. \lor L_k \lor C}{(L_1 \lor C)\sigma}\sigma = mgu(L_1, .., L_k)$$

Example 18.8

Let suppose we have a clause $P(x) \lor P(y)$. This clause is not economical.

We can derive $P(x)$ using factoring as follows

$$\textsc{Factor}\frac{P(x) \lor P(y)}{P(x)}\sigma = mgu(P(x), P(y)) = \{y \mapsto x\}$$

Commentary: Here is the application of resolution for step 3.

$$\frac{P(x) \lor P(y) \quad \neg P(x) \lor \neg P(y)}{P(x) \lor \neg P(y)} \sigma = \{x \mapsto x, y \mapsto y, y \mapsto x\}$$

Please note that the above σ is not an output of Robinson algorithm, since y and y do not occur in the unifying atoms. How do we understand this?

Example: why FACTOR rule?

Example 18.9

1. $P(x) \lor P(y)$ — *Assumption*
2. $\neg P(x) \lor \neg P(y)$ — *Assumption*
3. $P(x) \lor \neg P(y)$ — RESOLUTION *applied to 1 and 2*
4. $P(x)$ — *Applied* FACTOR *applied to 1*
5. $\neg P(x)$ — FACTOR *applied to 2*
6. $\perp$ — RESOLUTION *applied to 4 and 5*

No progress without FACTOR

In the above, we have written the consequences as a sequence, which is equivalent to the DAGs.

Exercise 18.1

Why do we not need similar rule for propositional logic?

Resolution theorem proving : apply equality over clauses

The rule does not differentiate between $s = t$ and $t = s$

$$\textsc{Paramodulation}\frac{s = t \vee C \quad D(u)}{(C \vee D(t))\sigma}\sigma = mgu(s, u)$$

Example 18.10

Consider clauses $f(x) = d \vee \underbrace{P(x)}_{C}$ *and* $\underbrace{Q(f(y))}_{D}$

$$\frac{f(x) = d \vee P(x) \qquad Q(f(y))}{P(y) \vee Q(d)}\sigma = mgu(f(x), f(y)) = \{x \mapsto y\}$$

Commentary: From some σ we have the following implications $(s = t \vee C) \Rightarrow (s\sigma = t\sigma \vee C\sigma)$, $D \Rightarrow D\sigma$, $(s\sigma = t\sigma \vee C\sigma) \wedge (D\sigma) \Rightarrow (s\sigma = t\sigma \wedge (D\sigma)) \vee C\sigma$, $(s\sigma = t\sigma \wedge (D\sigma)) \Rightarrow D\sigma\{s\sigma \mapsto t\sigma\}$. Therefore, Paramodulation is valid.

Resolution theorem proving : finishing disequality

If we have a disequality, we can eliminate it if both sides can be unified.

$$\text{RELEXIVITY}\frac{t \neq u \vee C}{C\sigma}\sigma = mgu(t, u)$$

Example 18.11

The following derivation removes a literal from the clause.

$$\text{RELEXIVITY}\frac{x \neq f(y) \vee P(x)}{P(f(y))}\sigma = mgu(x, f(y)) = \{x \mapsto f(y)\}$$

Example: a resolution proof

Example 18.12

Consider the following set of input clauses

1. $\neg Mother(x, y) \vee husbandOf(y) = fatherOf(x)$
2. $Mother(geoff, maggie)$
3. $bob = husbandOf(maggie)$
4. $fatherOf(geoff) \neq bob$ *Assumption*
5. $husbandOf(maggie) = fatherOf(geoff)$ *Resolution applied to 1 and 2*
6. $bob = fatherOf(geoff)$ *Paramodulation applied to 3 and 5*
7. $\perp$ *Resolution applied to 4 and 6*

Commentary: Example source `http://www.cs.miami.edu/home/geoff/Courses/TPTPSYS/FirstOrder/Paramodulation.shtml`

A resolution theorem prover: 5 rules to rule them all

$$\textsc{Assumption}\ \frac{}{C}\ C \in F$$

$$\textsc{Resolution}\ \frac{\neg A \vee C \quad B \vee D}{(C \vee D)\sigma}\ \sigma = mgu(A, B) \qquad \textsc{Factor}\ \frac{L_1 \vee .. \vee L_k \vee C}{(L_1 \vee C)\sigma}\ \sigma = mgu(L_1, .., L_k)$$

$$\textsc{Paramodulation}\ \frac{s = t \vee C \quad D(u)}{(C \vee D(t))\sigma}\ \sigma = mgu(s, u) \qquad \textsc{Relexivity}\ \frac{t \neq u \vee C}{C\sigma}\ \sigma = mgu(t, u)$$

Next semester

CS433 : automated reasoning

- How to make sat solvers efficient?
- FOL + arithmetic + decision procedures
- Applications to program verification

Topic 18.4

Problems

Exercise: prove the unsatisfiability

Exercise 18.2

Prove that the following set of clauses is unsatisfiable.

1. $Even(sum(twoSquared, b))$
2. $twoSquared = four$
3. $\neg Zero(x) \lor difference(four, x) = sum(four, x)$
4. $Zero(b)$
5. $\neg Even(difference(twoSquared, b))$

Exercise: prove the unsatisfiability

Exercise 18.3

Prove that the following set of clauses is unsatisfiable.

1. $P(f(a))$
2. $a = c$,
3. $\neg Q(x, x) \lor f(c) = f(d)$,
4. $Q(b, b)$
5. $\neg P(f(d))$

Proving transitivity

Exercise 18.4

Prove transitivity of equality using Paramodulation rule.

Exercise: merge factor and resolution

Exercise 18.5

a. Prove the following proof rule using Factor and Resolution.

$$\textsc{ExtendedResolution}\frac{\neg A_1 \vee ... \vee \neg A_m \vee C \quad B_1 \vee \cdots \vee B_n \vee D}{(C \vee D)\sigma}\sigma = mgu(A_1, \ldots, A_m, B_1, \ldots, B_n)$$

b. Show that the above rule subsumes Factor and Resolution.

c. Prove/disprove that we can replace both Resolution rule and Factor rule in the FOL resolution proof system by the following rule without losing completeness of deriving false clause.

$$\textsc{RestrictedExtendedResolution}\frac{\neg A \vee C \quad B_1 \vee \cdots \vee B_n \vee D}{(C \vee D)\sigma}\sigma = mgu(A, B_1, \ldots, B_n)$$

General resolution

Exercise 18.6

Consider the following general resolution rule which is more liberal in application in comparison to the standard FOL resolution. Show that if the following proof rule can prove $\perp$ *then the standard resolution can also do it.*

$$\textsc{GeneralResolution}\frac{\neg A \lor C \quad B \lor D}{(C\sigma_1 \lor D\sigma_2)} A\sigma_1 = B\sigma_2$$

Topic 18.5

Extra slides: optimizations in resolution proof systems

Example: Redundancies due to equality reasoning

Example 18.13

Consider the following clauses

1. $a = c$
2. $b = d$
3. $P(a, b)$
4. $\neg P(c, d)$ — *Assumption*
5. $P(c, b)$ — PARAMODULATION *applied to 1 and 3*
6. $P(a, d)$ Redundant derivation — PARAMODULATION *applied to 2 and 3*
7. $P(c, d)$ — PARAMODULATION *applied to 2 and 5*
8. $\perp$ — RESOLUTION *applied to and 7*

- Many clauses can be derived due to simple permutations
- Often derived clauses do not add new information
- A typical solver restricts application of the rules by imposing order

End of Lecture 18