

# CS 105: DIC on Discrete Structures

Instructor : S. Akshay

Sept 07, 2023

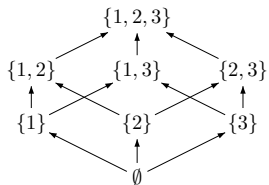
Lecture 14 – a little bit on lattices and on to Counting

## Recap: Partial order relations

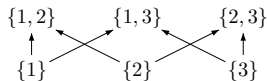
### Last two classes we saw

- ▶ Partial orders: definition and examples
- ▶ Posets, chains and anti-chains
- ▶ Graphical representation as Directed Acyclic Graphs
- ▶ Topological sorting (application to task scheduling)
- ▶ Mirsky's theorem (application to parallel task scheduling)

# Minimal and maximal elements

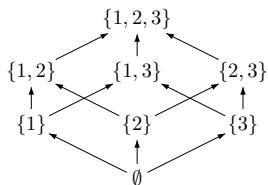


Poset  $P_1 = (S_1, \subseteq)$   
where  $S_1 = \mathcal{P}(\{1, 2, 3\})$

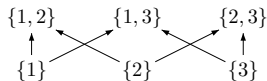


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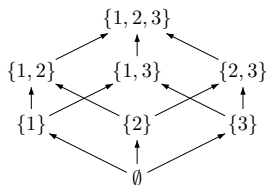


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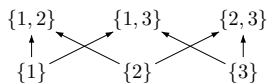
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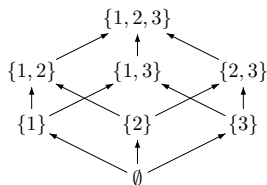


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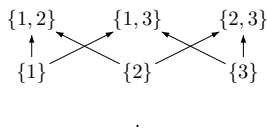
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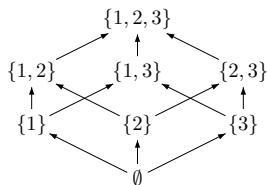
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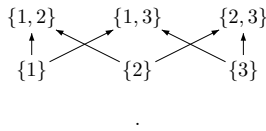
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**Exercise:** What are the maximal/maximum elements in  $P_1, P_2$ ?

## Least upper bounds and greatest lower bounds

Let  $(S, \preceq)$  be a poset and  $A \subseteq S$ .

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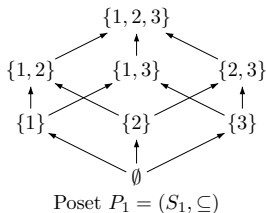
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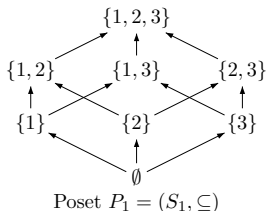
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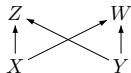


- ▶ Let  $A = \{\{1\}, \{2\}\}$ . Then  $\{1, 2\}, \{1, 2, 3\}$  are upper bounds of  $A$  in  $P_1$  and  $\{1, 2\}$  is the lub of  $A$ .

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- ▶ Consider  $P_3 = (S_3, \preceq)$  where  $S_3 = \{X, Y, Z, W\}$  and the  $\preceq$  is as given by the arrows. Let  $B = \{X, Y\}$ . Then  $Z, W$  are both upper bounds of  $B$  in  $P_3$ , but  $B$  has no lub in  $P_3$ .

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Some Observations (**Exercise**: Prove it!)

- ▶ The lub/glb of a subset  $A$  in  $S$ , if it exists, is unique.
- ▶ If the lub/glb of  $A \subseteq S$  belongs to  $A$ , then it is the greatest/least element of  $A$ .

# Lattices

## Definition

- ▶ A **lattice** is a poset in which every pair of elements has both a lub and a glb (in the set), i.e.,  $\forall x, y \in S$ , there exists  $l, u \in S$  such that  $l$  is the glb and  $u$  is the lub of  $\{x, y\}$ .

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- ▶ Finite lattices have a strong link with Boolean Algebra
- ▶ Several other applications in many domains of mathematics and CS, including **formal semantics of programming languages, program verification.**

## Summary till now

### Course Outline

1. Proofs and structures
2. Counting and combinatorics
3. Introduction to graph theory

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- ▶ Relations: equivalence relations, partial orders, lattices
- ▶ Some applications

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  - ▶ Functions: To compare infinite sets
  - ▶ Using diagonalization to prove impossibility results.
  - ▶ Equivalences: Defining “like” partitions.
  - ▶ Posets: Topological sort, (parallel) task scheduling, lattices

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# Pop Quiz



## Pop Quiz

Fill the feedback form at

<https://forms.gle/uP39XHMMmqmx63qUTA>

## Next chapter: Counting and Combinatorics

### Topics to be covered

- ▶ Basics of counting
- ▶ Subsets, partitions, Permutations and combinations
- ▶ Pigeonhole Principle and its extensions
- ▶ Recurrence relations and generating functions