

CS 105: DIC on Discrete Structures

Graph theory

Basic terminology, Eulerian walks

Lecture 24

Oct 12 2023

Topic 3: Graph theory

Topics covered in the last three lectures

- ▶ Pigeon-Hole Principle and its extensions.
- ▶ A glimpse of Ramsey theory

Topic 3: Graph theory

Topics covered in the last three lectures

- ▶ Pigeon-Hole Principle and its extensions.
- ▶ A glimpse of Ramsey theory
- ▶ Ramsey numbers $R(k, \ell)$: Existence of **regular sub-structures** in large enough **general structures**.

Topic 3: Graph theory

Topics covered in the last three lectures

- ▶ Pigeon-Hole Principle and its extensions.
- ▶ A glimpse of Ramsey theory
- ▶ Ramsey numbers $R(k, \ell)$: Existence of **regular sub-structures** in large enough **general structures**.
- ▶ **Structures = Graphs**,

Topic 3: Graph theory

Topics covered in the last three lectures

- ▶ Pigeon-Hole Principle and its extensions.
- ▶ A glimpse of Ramsey theory
- ▶ Ramsey numbers $R(k, \ell)$: Existence of **regular sub-structures** in large enough **general structures**.
- ▶ **Structures = Graphs, regular = monochromatic/complete**

Topic 3: Graph theory

Topics covered in the last three lectures

- ▶ Pigeon-Hole Principle and its extensions.
- ▶ A glimpse of Ramsey theory
- ▶ Ramsey numbers $R(k, \ell)$: Existence of **regular sub-structures** in large enough **general structures**.
- ▶ **Structures = Graphs, regular = monochromatic/complete**

Next topic

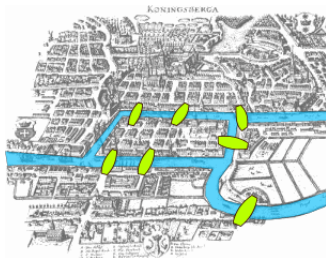
Graphs and their properties!

Topic 3: Graph theory

Textbook Reference

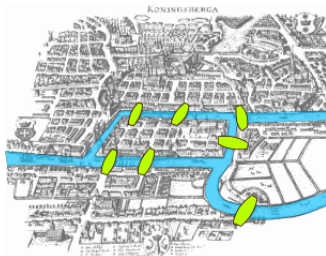
- ▶ **Introduction to Graph Theory, 2nd Ed., by Douglas West.**
- ▶ Low cost Indian edition available, published by PHI Learning Private Ltd.

Königsberg Bridge problem



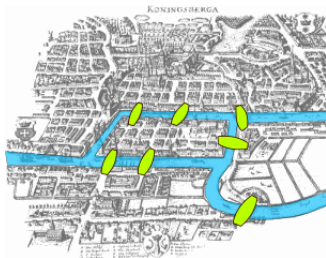
- ▶ In 18th century Prussia, the city on river Pregel...
- ▶ Find a walk from home, crossing every bridge exactly once and returning home.

Königsberg Bridge problem



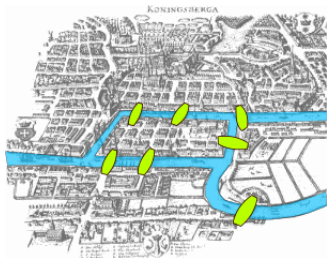
- ▶ In 18th century Prussia, the city on river Pregel...
- ▶ Find a walk from home, crossing every bridge exactly once and returning home.
- ▶ They couldn't find the answer, so in 1735 they asked Leonard Euler, mathematician in St. Petersburg.

Königsberg Bridge problem



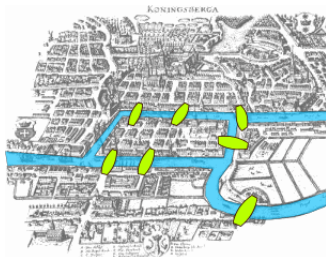
- ▶ In 18th century Prussia, the city on river Pregel...
- ▶ Find a walk from home, crossing every bridge exactly once and returning home.
- ▶ They couldn't find the answer, so in 1735 they asked Leonard Euler, mathematician in St. Petersburg.

Königsberg Bridge problem



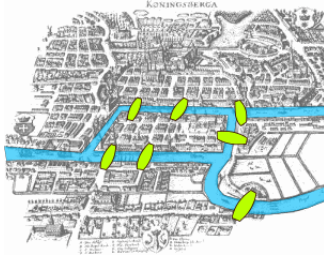
- ▶ In 18th century Prussia, the city on river Pregel...
- ▶ Find a walk from home, crossing every bridge exactly once and returning home.
- ▶ *“This question is so banal, but seemed to me worthy of attention in that [neither] geometry, nor algebra, nor even the art of counting was sufficient to solve it.”*

Königsberg Bridge problem

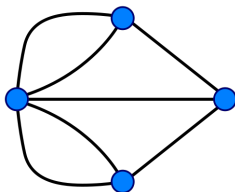
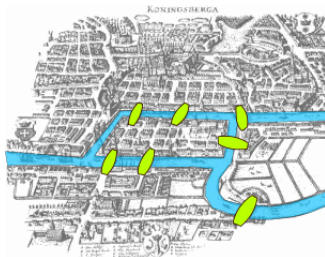


- ▶ In 18th century Prussia, the city on river Pregel...
- ▶ Find a walk from home, crossing every bridge exactly once and returning home.
- ▶ *“This question is so banal, but seemed to me worthy of attention in that [neither] geometry, nor algebra, nor even the art of counting was sufficient to solve it.”*
- ▶ Still, he wrote a paper showing that this is impossible!
- ▶ Thus, he “gave birth” to the area of graph theory.

Königsberg Bridge problem

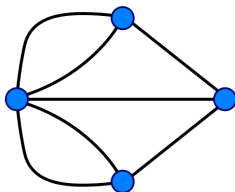
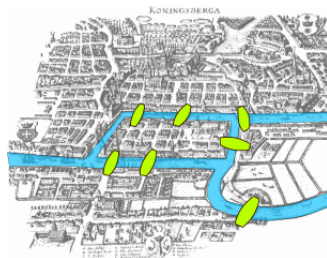


Königsberg Bridge problem



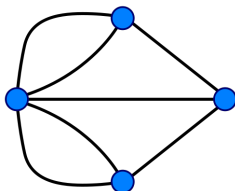
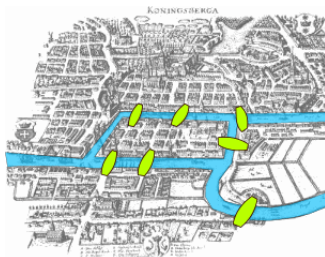
- Qn: Find a walk from home, which crosses every bridge exactly once and returns home.

Königsberg Bridge problem



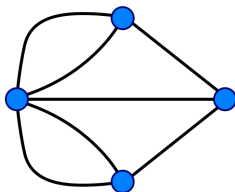
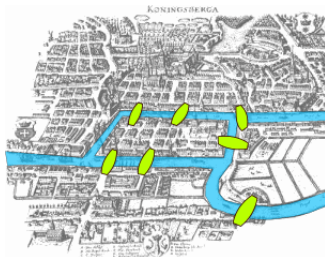
- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:

Königsberg Bridge problem



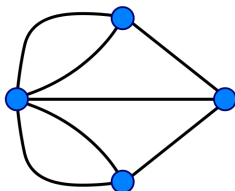
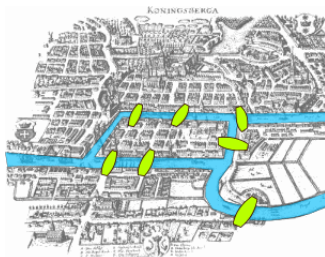
- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.

Königsberg Bridge problem



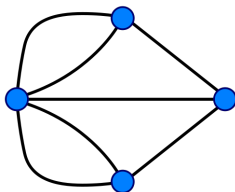
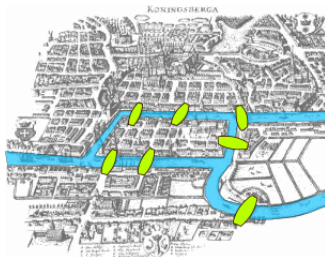
- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.
 - ▶ So each vertex must be connected to an even no. of vertices.

Königsberg Bridge problem



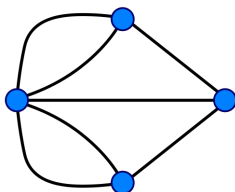
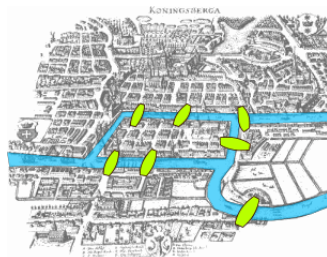
- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.
 - ▶ So each vertex must be connected to an even no. of vertices.
 - ▶ Which is not the case here, hence it is impossible.

Königsberg Bridge problem



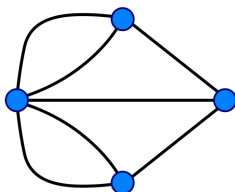
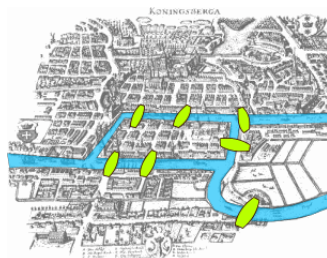
- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.
 - ▶ So each vertex must be connected to an even no. of vertices.
 - ▶ Which is not the case here, hence it is impossible.
- ▶ Clearly, this is a sufficient condition, but is it necessary?

Königsberg Bridge problem



- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.
 - ▶ So each vertex must be connected to an even no. of vertices.
 - ▶ Which is not the case here, hence it is impossible.
- ▶ Clearly, this is a sufficient condition, but is it necessary?
- ▶ If every vertex is connected to an even no. of vertices in a graph, is there such a walk?

Königsberg Bridge problem



- ▶ Qn: Find a walk from home, which crosses every bridge exactly once and returns home.
- ▶ Leonard Euler showed in 1735 that this is impossible. The argument is as follows:
 - ▶ Each time you enter or leave a vertex, you use an edge.
 - ▶ So each vertex must be connected to an even no. of vertices.
 - ▶ Which is not the case here, hence it is impossible.
- ▶ Clearly, this is a sufficient condition, but is it necessary?
- ▶ If every vertex is connected to an even no. of vertices in a graph, is there such a walk? This is called **Eulerian walk**.

What are graphs

What are graphs

Definition

A simple graph G is a pair (V, E) of a set of vertices/nodes V and edges E between unordered pairs of vertices called end-points: $e = vu$ means that e is an edge between v and u ($u \neq v$).

What are graphs

Definition

A simple graph G is a pair (V, E) of a set of vertices/nodes V and edges E between unordered pairs of vertices called end-points: $e = vu$ means that e is an edge between v and u ($u \neq v$).

- ▶ What about loops?
- ▶ What about directed edges?
- ▶ What about multiple edges?

What are graphs

Definition

A simple graph G is a pair (V, E) of a set of vertices/nodes V and edges E between unordered pairs of vertices called end-points: $e = vu$ means that e is an edge between v and u ($u \neq v$).

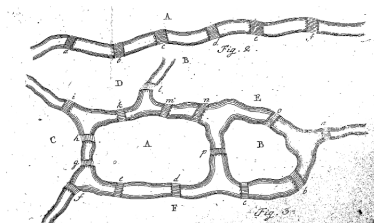
- ▶ What about loops?
- ▶ What about directed edges?
- ▶ What about multiple edges?

General Definition

A graph G is a triple V, E, R where V is a set of vertices, E is a set of edges and $R \subseteq E \times V \times V$ is a relation that associates each edge with two vertices called its end-points.

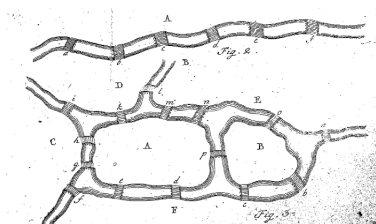
We will consider **only** finite graphs (i.e., $|V|, |E|$ are finite) and **often** deal with simple graphs.

Basic terminology



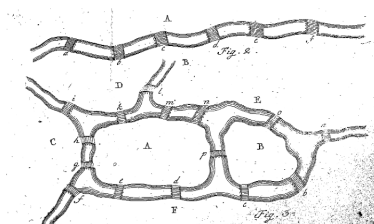
Ex. Draw the graphs!

Basic terminology



The **degree** $d(v)$ of a vertex v (in an undirected loopless graph) is the number of edges incident to it, i.e., $|\{vw \in E \mid w \in V\}|$.
A vertex of degree 0 is called an **isolated vertex**.

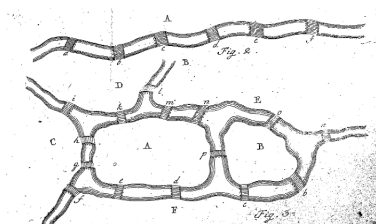
Basic terminology



The **degree** $d(v)$ of a vertex v (in an undirected loopless graph) is the number of edges incident to it, i.e., $|\{vw \in E \mid w \in V\}|$. A vertex of degree 0 is called an **isolated vertex**.

- ▶ A **walk** is a sequence of vertices v_1, \dots, v_k such that $\forall i \in \{1, \dots, k-1\}, (v_i, v_{i+1}) \in E$. The vertices v_1 and v_k are called the **end-points** and others are called **internal vertices**.

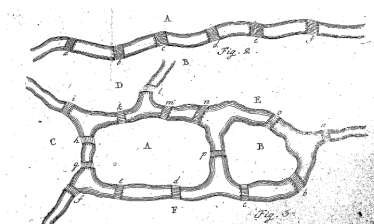
Basic terminology



The **degree** $d(v)$ of a vertex v (in an undirected loopless graph) is the number of edges incident to it, i.e., $|\{vw \in E \mid w \in V\}|$. A vertex of degree 0 is called an **isolated vertex**.

- ▶ A **walk** is a sequence of vertices v_1, \dots, v_k such that $\forall i \in \{1, \dots, k-1\}, (v_i, v_{i+1}) \in E$. The vertices v_1 and v_k are called the **end-points** and others are called **internal vertices**.
- ▶ A walk is called **closed** if it starts and ends with the same vertex, i.e., its endpoints are the same.

Basic terminology



The **degree** $d(v)$ of a vertex v (in an undirected loopless graph) is the number of edges incident to it, i.e., $|\{vw \in E \mid w \in V\}|$. A vertex of degree 0 is called an **isolated vertex**.

- ▶ A **walk** is a sequence of vertices v_1, \dots, v_k such that $\forall i \in \{1, \dots, k-1\}, (v_i, v_{i+1}) \in E$. The vertices v_1 and v_k are called the **end-points** and others are called **internal vertices**.
- ▶ A walk is called **closed** if it starts and ends with the same vertex, i.e., its endpoints are the same.

Graph is **connected** if there is a walk between any two vertices.

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

- ▶ If G is Eulerian, each vertex must have an even degree.
- ▶ This is a necessary condition, but is it sufficient?

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

- ▶ If G is Eulerian, each vertex must have an even degree.
- ▶ This is a necessary condition, but is it sufficient?

Theorem

A graph G with no isolated vertices is Eulerian iff it is connected and all its vertices have even degree.

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

- ▶ If G is Eulerian, each vertex must have an even degree.
- ▶ This is a necessary condition, but is it sufficient?

Theorem

A graph G with no isolated vertices is Eulerian iff it is connected and all its vertices have even degree.

Proof:

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

- ▶ If G is Eulerian, each vertex must have an even degree.
- ▶ This is a necessary condition, but is it sufficient?

Theorem

A graph G with no isolated vertices is Eulerian iff it is connected and all its vertices have even degree.

Proof: (\implies)

- ▶ Suppose G is Eulerian: every vertex has even degree.
 - ▶ each passage through a vertex uses two edges (in and out).
 - ▶ at the first vertex first edge is paired with last.

Eulerian graphs

Definition

A graph is called **Eulerian** if it has a closed walk that contains all edges, and each edge occurs exactly once. Such a walk is called an **Eulerian walk**.

- ▶ If G is Eulerian, each vertex must have an even degree.
- ▶ This is a necessary condition, but is it sufficient?

Theorem

A graph G with no isolated vertices is Eulerian iff it is connected and all its vertices have even degree.

Proof: (\implies)

- ▶ Suppose G is Eulerian: every vertex has even degree.
 - ▶ each passage through a vertex uses two edges (in and out).
 - ▶ at the first vertex first edge is paired with last.
- ▶ Any two edges are in the same walk implies graph is connected (unless it has isolated vertices).