

CS 105: DIC on Discrete Structures

Graph theory Connectedness in graphs

Lecture 29
Oct 26 2023

Topic 3: Graph theory

Recap

1. Basic definitions: graphs, paths, cycles, walks, trails; connected graphs.
2. Eulerian graphs and a characterization in terms of degrees of vertices.
3. Bipartite graphs and a characterization in terms of odd length cycles.
4. Graph representation (as matrices, lists, etc.)
5. Graph isomorphisms and automorphisms
6. Subgraphs:
 - ▶ Cliques and independent sets,
 - ▶ A zoo of graphs.

Reference: Sections 1.1-1.3 of Chapter 1 from **Douglas West**.

Recall: Graph Automorphisms

Definition

An **isomorphism** from simple graph G to H is a **bijection** $f : V(G) \rightarrow V(H)$ such that $uv \in E(G)$ iff $f(u)f(v) \in E(H)$.

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- ▶ How many automorphisms does $K_{r,s}$ have?

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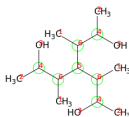
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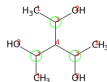
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Automorphisms are a measure of symmetry.

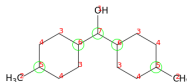
Practical applications in graph drawing, visualization, molecular symmetry, structured boolean satisfiability, formal verification



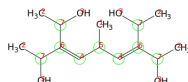
Symmetry Classes
Stereogenic Atoms
(rule 2b')



Symmetry Classes
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(rule 2b', only 1 true stereocenter)



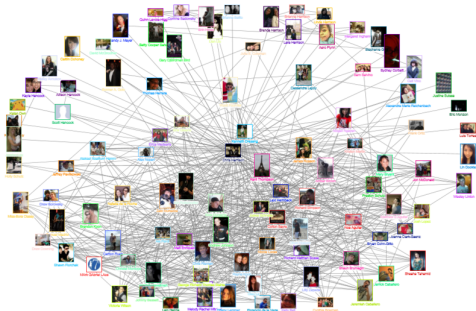
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Symmetry Classes
Stereogenic Atoms
(rule 3)

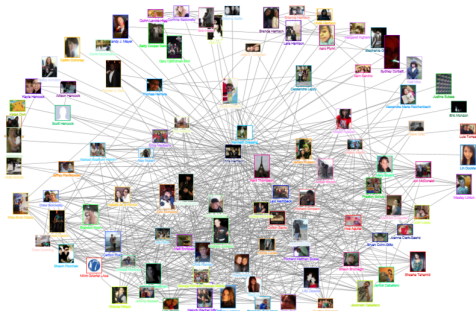


Cliques and independent sets



- ▶ Consider a large social network graph where friends are linked by an edge.
- ▶ What is the largest clique of friends?
- ▶ If we want to spread a youtube video, how many people should we send it to so that we are guaranteed everyone will see it (assuming friends forward to each other)?

Cliques and independent sets



Cliques and independent sets

- ▶ A **clique** in a graph is a set of pairwise adjacent vertices.
- ▶ An **independent set** in a graph is a set of pairwise non-adjacent vertices.

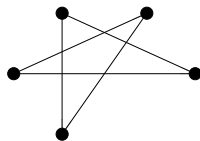
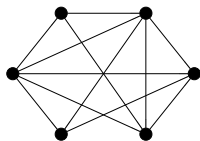
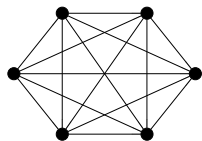
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Cliques and independent sets

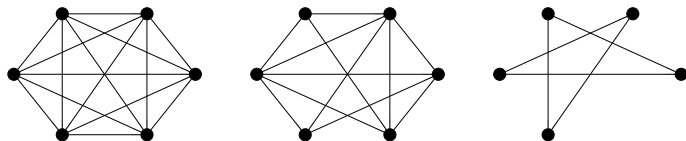
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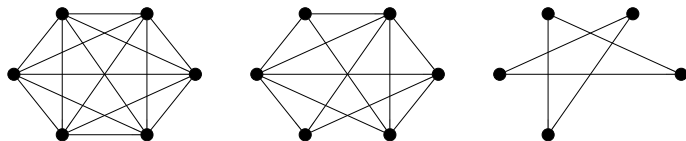


- ▶ Thus, a **clique** in a graph G is a complete subgraph of G .

Subgraphs of a graph G

A subgraph H of a graph G is a graph H such that $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$ (and the assignment of endpoints to edges in H is same as in G).

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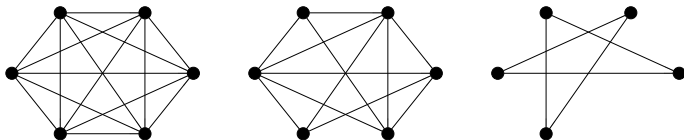


- ▶ Thus, a **clique** in a graph G is a complete subgraph of G .
- ▶ An **independent set** in G is a complete subgraph of \overline{G} , where \overline{G} is the **complement of G** obtained by making all adjacent vertices non-adjacent and vice versa.

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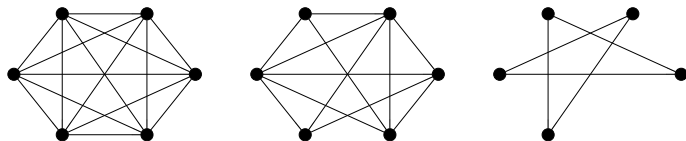
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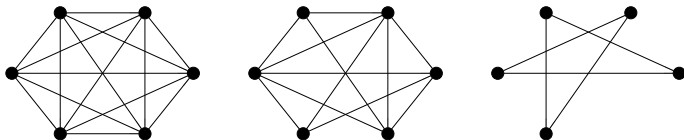
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- ▶ What is the size of the largest clique/independent set in each of the above graphs? In any complete graph?
- ▶ Given graph G , integer k , does G have a clique of size k ?

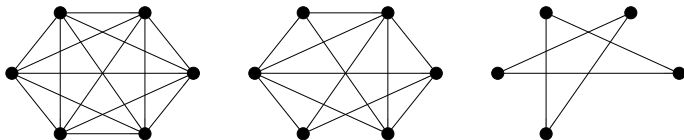
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- ▶ In a graph with 6 vertices, can you always find a clique or an independent set of size 3?

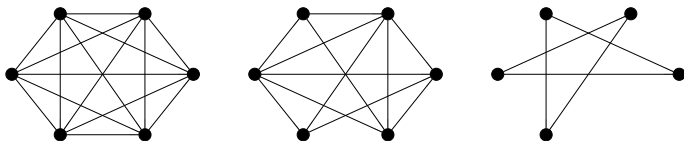
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Ramsey's theorem - restated

In any graph with $R(k, \ell)$ vertices, there exists either a clique of size k or an independent set of size ℓ .

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Definition

A **(connected) component** of G is a **maximal connected subgraph**, i.e., a subgraph that is connected and is not contained in any other connected subgraph of G .

Thus, equivalence classes of P are the vertex sets of the components of G .

Recall: Difference between maximal and maximum

- ▶ Is every maximal path maximum, i.e., have maximum length?
- ▶ A maximal structure is a structure that is **not contained** in a larger structure, i.e., increasing the structure will violate some property.
- ▶ Maximum just means that size is the greatest among all possible.

Exercises!

1. Give a path which is maximal but not maximum.
2. Give a subgraph of a graph which is maximally connected, but not maximum (i.e., does not have maximum # edges).
3. How many maximal/maximum independent sets does $K_{r,s}$ have?

Components and cut-edges

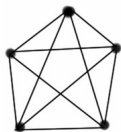
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K_5



$K_{2,3}$



P_5

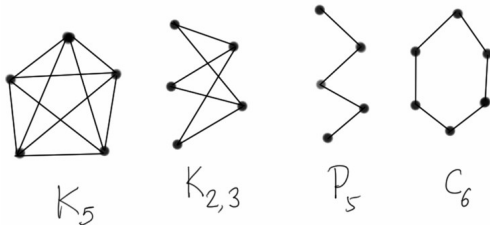


C_6

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Theorem: Characterize cut-edges using cycles

Exercise! An edge is a cut-edge iff it belongs to no cycle.