

CS 105: DIC on Discrete Structures

Graph theory Graph Isomorphism

Lecture 31
Oct 15 2024

Part 4: Graph theory

Recap of last **four** lectures:

1. Basics: graphs, paths, cycles, walks, trails; connected graphs.
2. **Eulerian graphs** and a characterization in terms of degrees of vertices.
3. **Bipartite graphs** and a characterization in terms of odd length cycles.

Reference: Sections 1.1-1.3 of Chapter 1 from **Douglas West**.

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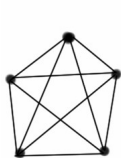
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Today

Graph representations and naming

Some special graphs and notations



K_5



$K_{2,3}$



P_5



C_6

- ▶ Complete graphs K_n
- ▶ Complete bipartite graphs $K_{i,j}$
- ▶ Paths P_n
- ▶ Cycles C_n

Some special graphs and notations

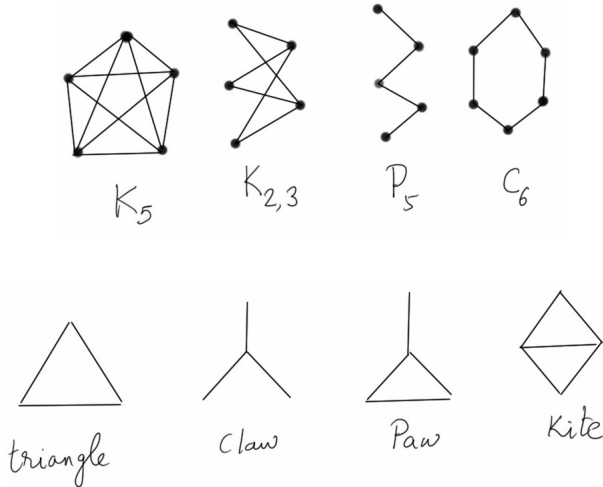
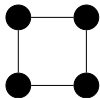
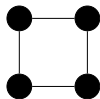


Figure: A whole graph zoo!

Are these graphs the same?



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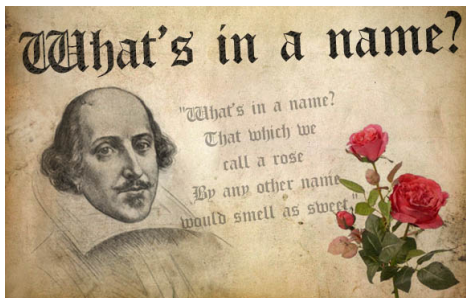


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Representing and comparing graphs

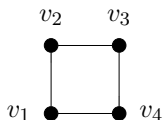
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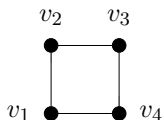
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► As an adjacency list:

v_1	v_2, v_4
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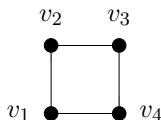
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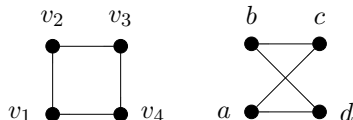
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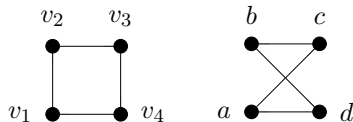
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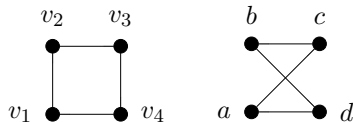
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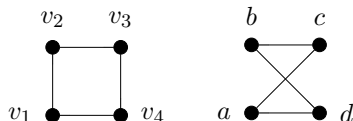
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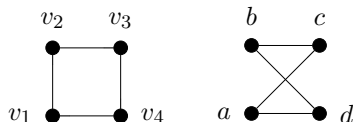
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- ▶ How do we formalize this?

Isomorphism

Definition

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- ▶ When we talked about an “unlabeled” graph till now, we actually meant the isomorphism class of that graph!

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- Are C_5 and $P_5 \cup \{e\}$ isomorphic?

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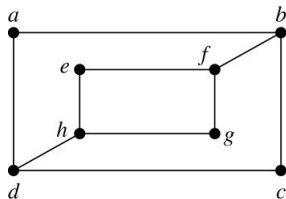
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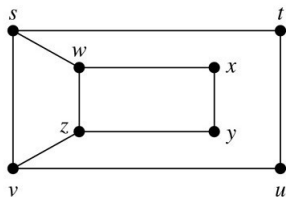
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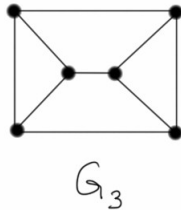
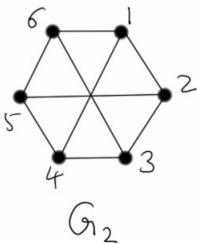
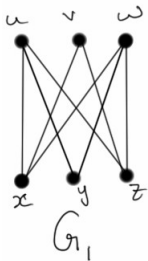
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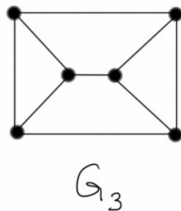
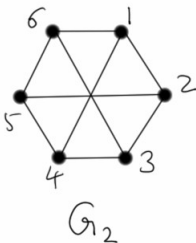
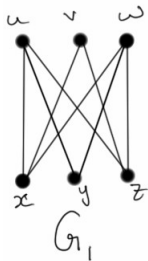
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5. G is bipartite iff H is bipartite.
6. ...

Graph isomorphism

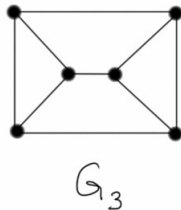
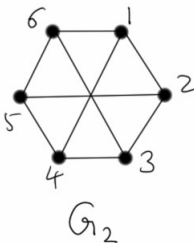
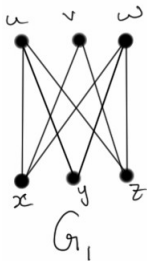


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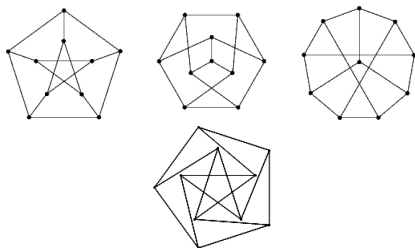
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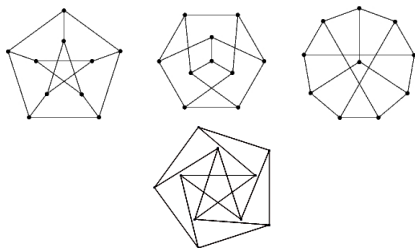
- ▶ To show that two graphs are isomorphic, you have to
 1. give names to vertices
 2. specify a bijection
 3. check that it preserves the adjacency relation
- ▶ To show that two graphs are **non-isomorphic**, find a structural property that is different.

Is checking graph isomorphism easy?



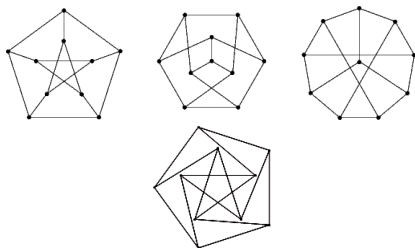
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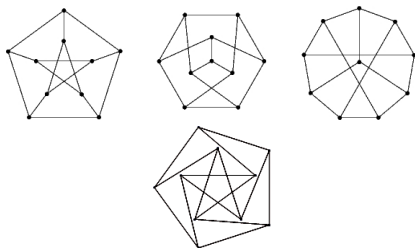
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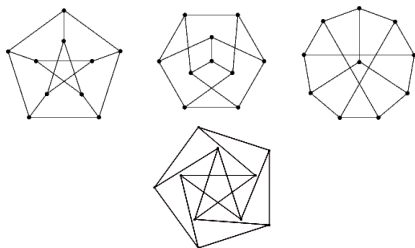
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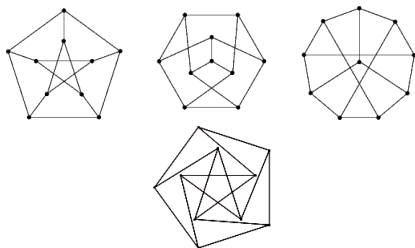
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Further reading: Graph and sub-graph isomorphism problems.

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- ▶ How many automorphisms does $K_{r,s}$ have?

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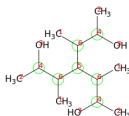
Definition

An **isomorphism** from simple graph G to H is a **bijection** $f : V(G) \rightarrow V(H)$ such that $uv \in E(G)$ iff $f(u)f(v) \in E(H)$.

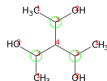
An **automorphism** of G is an isomorphism from G to itself, i.e. a **bijection** $f : V(G) \rightarrow V(G)$ s.t. $uv \in E(G)$ iff $f(u)f(v) \in E(G)$.

Automorphisms are a measure of symmetry.

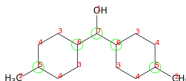
Practical applications in graph drawing, visualization, molecular symmetry, structured boolean satisfiability, formal verification



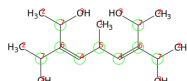
Symmetry Classes
Stereogenic Atoms
(rule $2b'$)



Symmetry Classes
Stereogenic Atoms
(rule $2b'$, only 1 true stereocenter)



Symmetry Classes
Stereogenic Atoms
(rule $2a''$)



Symmetry Classes
Stereogenic Atoms
(rule 3)

