CS 105: DIC on Discrete Structures

Graph theory Connectedness in graphs

> Lecture 32 Oct 17 2024

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Last topic: Graph theory

Recap

- 1. Basic definitions: graphs, paths, cycles, walks, trails; connected graphs.
- 2. Eulerian graphs and a characterization in terms of degrees of vertices.
- 3. Bipartite graphs and a characterization in terms of odd length cycles.
- 4. Graph representation (as matrices, lists, etc.)
- 5. Graph isomorphisms and automorphisms

Reference: Sections 1.1-1.3 of Chapter 1 from Douglas West.

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Automorphisms are a measure of symmetry.

Practical applications in graph drawing, visualization, molecular symmetry, structured boolean satisfiability, formal verification





- Consider a large social network graph where friends are linked by an edge.
- ▶ What is the largest clique of friends?
- ► If we want to spread a youtube video, how many people should we send it to so that we are guaranteed everyone will see it (assuming friends forward to each other)?



Cliques and independent sets

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▶ Thus, a clique in a graph G is a complete subgraph of G.

Subgraphs of a graph ${\cal G}$

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 $V(H) \subseteq V(G)$ and

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& assignment of endpoints to edges in H is same as in G.



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An independent set in G is a complete subgraph of G
, where G
 is the complement of G obtained by making all adjacent vertices non-adjacent and vice versa.

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Ramsey's theorem - restated

In any graph with $R(k, \ell)$ vertices, there exists either a clique of size k or an independent set of size ℓ .

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Definition

A (connected) component of G is a maximal connected subgraph, i.e., a subgraph that is connected and is not contained in any other connected subgraph of G.

Thus, equivalence classes of P are the vertex sets of the components of G.

Recall: Difference between maximal and maximum

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Exercises!

- 1. Give a path which is maximal but not maximum.
- 2. Give a subgraph of a graph which is maximally connected, but not maximum (i.e., does not have maximum # edges).
- 3. How many maximal/maximum independent sets does $K_{r,s}$ have?

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H.W. Exercise!

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Theorem: Characterize cut-edges using cycles

H.W. Exercise! An edge is a cut-edge iff it belongs to no cycle.