CS 105: DIC on Discrete Structures

Graph theory Matchings! Guest Lecture: Rohit Gurjar

> Lecture 33 Oct 21 2024

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- 2. Eulerian graphs: characterization using degrees of vertices.
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A General Recipe

- 1. Consider an "interesting" problem.
- 2. Model it as a "special" class of graphs.
- 3. Characterize them using properties on vertices, cycles etc.
- 4. Not done in this course: Build algorithms based on the characterization, analyze and implement them!

This lecture, we will consider a new problem:

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- ▶ What are the properties of such an assignment?
- ► Another practical example: the dating scene!

- ▶ A matching in a graph G is a set of (non-loop) edges with no shared end-points.
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