



# CS230: Digital Logic Design and Computer Architecture Lecture 6: MIPS instructions contd.

https://www.cse.iitb.ac.in/~biswa/courses/CS230/main.html

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# Phones (smart/non-smart) on silence plz, Thanks

111M

## Logistics

When: Quiz-I, January 27, 11:15 AM Where: LA-001 and CC-105 Report by 11.10 AM

Duration: 1 hour

Make sure you clear all your doubts on Monday 1:30 to 5:30 PM

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## Recap

- ISA
- Assembly
- Machine level
- Instructions



## Memory Instructions



#### Stored Program & Von Neumann



## Memory



4GB of Memory (DRAM)

Say, a word: four bytes

How to access instructions: Program Counter (PC) A register that stores the address of the instruction

32-bit processor: addresses are of width 32 bits (devil is in the details ☺ )

So the processor fetches PC, PC+4, PC+8, ..... in a sequential order

## 1946 onwards

#### Since 1946 all computers have had 5 components



## Example (Remember PC for the time being)

PCX: lw PCY: add PCZ: lui

#### PCZ=PCY+4 and PCY = PCX+4



# Why Memory? Why Not Registers?

- Registers are limited. More #registers, higher access time.
- How? we will see sooner than later.
- Let's focus on the data part now. How to access data for our instructions?

## Memory Instructions







## Both instructions and data from memory

g = h + A [8];

PCX: w \$t0, 8(\$3) # A[8]
PCY: add \$s1, \$s2, \$t0 # g = h + t0



PCY = PCX+4



## A quick recap

Von Neumann (stored program) concept

As registers are limited, data can be there in the registers or in the memory

Register accesses are through register names/numbers

Memory accesses are through addresses stored in registers

Let's move on: Decision Making Instructions • Decisions: if, else ....

Two instructions: beq (branch equals to) and bne (branch not equals to)

beq \$t0, \$t1, L1 bne \$t0, \$t1, L1

**Computer Architecture** 

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## Branch Instructions: Conditional branches

beq \$t0, \$t1, L1

goto L1 (statements labeled as L1) if \$t0 equals \$t1

bne \$t0, \$t1, L1

goto L1 (statements labeled as L1) if \$t0 does not equal to \$t1

The slt instruction (Set on less than) if (a < b) // beq and bne won't work here

else c=0

c=1

slt \$t3, \$t1, \$t2 // t1 and t2 contain a and b

We can slti too; one of the operand will be a constant

## Loops: How to deal with it?

while(CS230[i] == k) i+=1;

say i and k are in \$s3 and \$s5, and the base of CS230 in \$s6

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## Loops continued

## while(CS230[i] == k) i+=1;

- 1. LOAD CS230[i], base address of CS230 is in \$s6
- 2. We need to go to CS230[i]
- Assuming CS230 is an integer array, each index is of 4 bytes. We need to go to CS230 [i\*4 bytes]

## Loops contd. (\$s3=i, \$s5=k, \$s6=base address)

Exit: // do nothing

Where is the Loop?

## Loops continued

Loop: sll \$t1, \$s3, 2 // i\*4 while(CS230[i] == k) add \$t1, \$t1, \$s6 // address of CS230[i] lw \$t0, 0(\$t1) // t0 = CS230[i] bne \$t0, \$s5, Exit // go to Exit if CS230[i] not equals to k addi \$s3, \$s3,1 // i=i+1

Exit:

// do nothing

How to jump to the Loop?

### Loops continued

Loop: sll \$t1, \$s3, 2 // i\*4 while(CS230[i] == k) add \$t1, \$t1, \$s6 // address of CS230[i] i+=1; lw \$t0, 0(\$t1) // t0 = CS230[i] bne \$t0, \$s5, Exit // go to Exit if CS230[i] not equals to k addi \$s3, \$s3,1 // i=i+1 j Loop // go to loop. j here is jump

Exit: // do nothing

# Textbook Chapter 2 P&H







# haben Sie einen guten Tag

