



# CS305: Computer Architecture

World of Instructions-III (The MIPS language)

<https://www.cse.iitb.ac.in/~biswa/courses/CS305/main.html>

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# Logistics

Get your hands dirty with lab assignments.

TAs (Rwitaban and Shashank) have done all the hard work to make it happen. You will learn a lot by doing things. Assignment-1 will be MIPS instructions in action.

A digression/clarification on the next slide: As per assignment zero.

# Instructor Biswa and THE Biswa



The only thing that is in common:  
we, both are from the state of Odisha 😊

Do not expect a stand-up from the instructor 😊

# A quick recap

## Von Neumann (stored program) concept

As registers are limited, data can be there in the registers or in the memory

Register accesses are through register names/numbers

Memory accesses are through addresses stored in registers

# Let's move on: Decision Making Instructions

- Decisions: if, else ....

Two instructions: beq (branch equals to) and  
bne (branch not equals to)

beq \$t0, \$t1, L1

bne \$t0, \$t1, L1

# Branch Instructions: Conditional branches

**beq** \$t0, \$t1, **L1**

*goto* L1 (statements labeled as L1) if \$t0 equals \$t1

**bne** \$t0, \$t1, **L1**

*goto* L1 (statements labeled as L1) if \$t0 does not equal to \$t1

# The slt instruction (Set on less than)

if (a < b) // beq and bne won't work here

    c=1

else

    c=0

**slt** \$t3, \$t1, \$t2 // t1 and t2 contain a and b

We can slti too; one of the operand will be a constant

# Loops: How to deal with it?

```
while(CS305[i] == k)
```

```
    i+=1;
```

say i and k are in \$s3 and \$s5, and the  
base of CS305 in \$s6



## Loops continued

```
while(CS305[i] == k)
```

```
    i+=1;
```

1. LOAD CS305[i], base address of CS305 is in \$s6
2. We need to go to CS305[i]
3. Assuming CS305 is an integer array, each index is of 4 bytes. We need to go to CS305 [i\*4 bytes]

Loops contd. (\$s3=i, \$s5=k, \$s6=base address)

```
sll $t1, $s3, 2      // i*4
add $t1, $t1, $s6    // address of CS305[i]
lw  $t0, 0($t1)      // t0 = CS305[i]
bne $t0, $s5, Exit   // go to Exit if CS305[i] not equals to k
addi $s3, $s3, 1     // i=i+1
```

while(CS305[i] == k)  
i+=1;

Exit: // do nothing

Where is the Loop?

# Loops continued

```
Loop: sll $t1, $s3, 2    // i*4
      add $t1, $t1, $s6  // address of CS305[i]
      lw  $t0, 0($t1)    // t0 = CS305[i]
      bne $t0, $s5, Exit // go to Exit if CS305[i] not equals to k
      addi $s3, $s3, 1   // i=i+1
```

while(CS305[i] == k)  
    i+=1;

Exit:                   // do nothing

How to jump to the Loop?

# Loops continued

```
Loop: sll $t1, $s3, 2      // i*4                                while(CS305[i] == k)
    add $t1, $t1, $s6      // address of CS305[i]                i+=1;
    lw $t0, 0($t1)         // t0 = CS305[i]
    bne $t0, $s5, Exit     // go to Exit if CS305[i] not equals to k
    addi $s3, $s3, 1       // i=i+1
    j     Loop             // go to loop. j here is jump
```

```
Exit:                // do nothing
```

# Sequential execution and jumps

PC, PC+4, PC+8, .....

PC, PC+4, {if condition here, TRUE} PC+32, .....

j instruction loads an immediate into the PC. It can be either specified as an offset or the label (assembler will convert this label into an offset). Next lecture: jr, jal, ..

# Dhanyavad