



CS305: Computer Architecture

World of Instructions-VI (The MIPS language)

<https://www.cse.iitb.ac.in/~biswa/courses/CS305/main.html>

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Quick recap

Register spilling, 32 MIPS registers, nested functions,

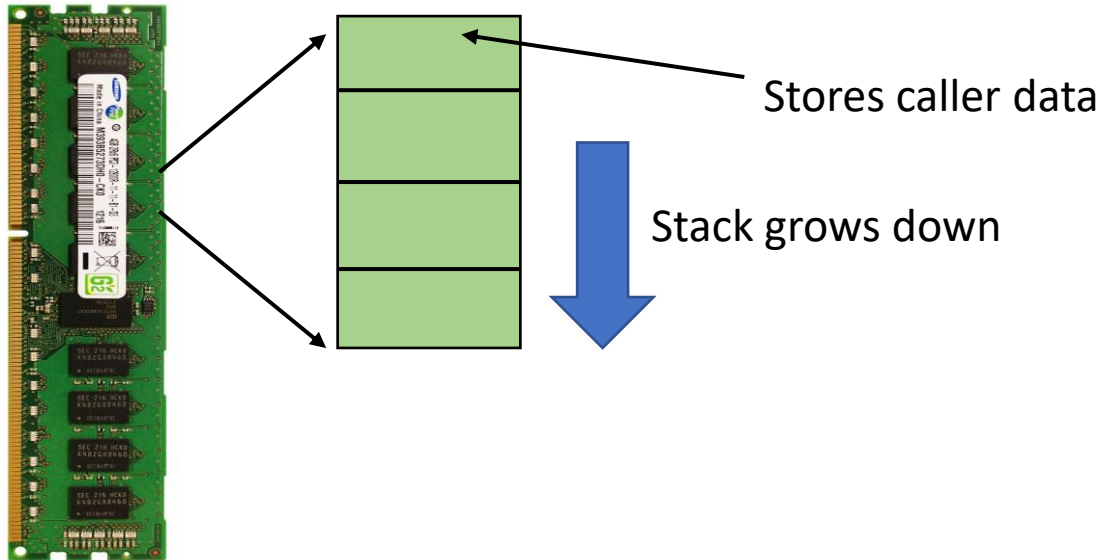
oh no!

Spilled registers: Where else can we store?

Where else can we store?

Remember previous lectures: registers or memory

MIPS way of handling it:
The Stack (part of DRAM, for each function call)



\$sp (stack pointer) points to the address where stack ends
One per function, private memory area, else the same
problem ☹️

Before that: Who does what?

In MIPS,

\$t0 to \$t9 (R8 to R15, R24, and R25) are temporary and *caller* saved registers. Register values not preserved across function calls (call-clobbered).

\$s0 to \$s7 (R16 to R23) are *callee* saved registers. Register values are preserved across function calls (call-preserved).

\$sp and **\$ra** are **caller or callee** saved registers ?

Before that: Who does what?

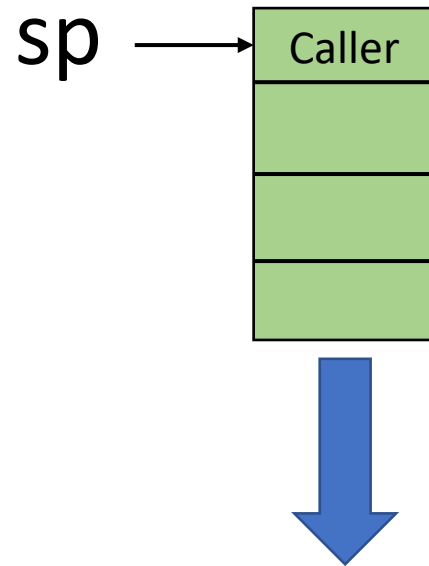
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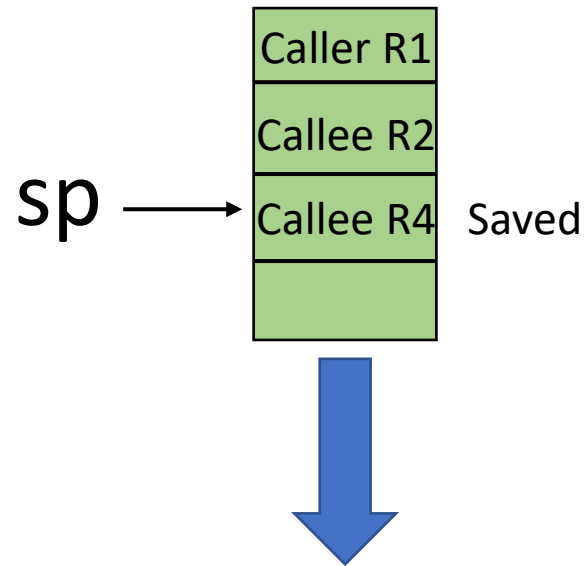
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\$sp and *\$ra* are *callee* saved registers.

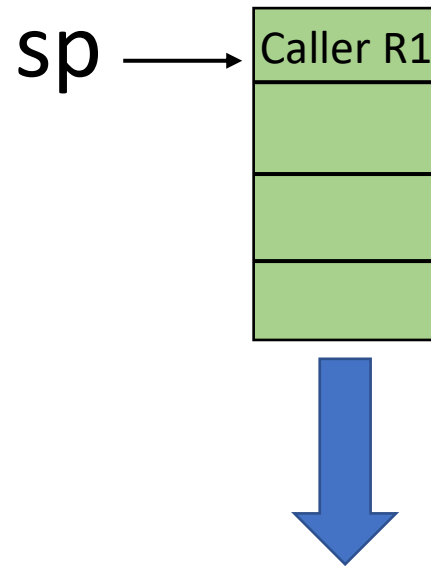
MIPS way of handling it: Before function call



MIPS way of handling it: Function call is ON



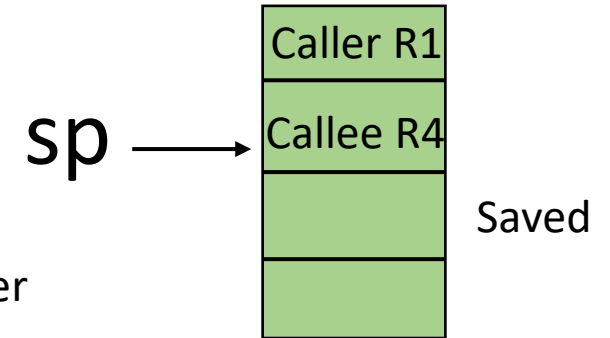
MIPS way of handling it: After the function call



How to save and restore?

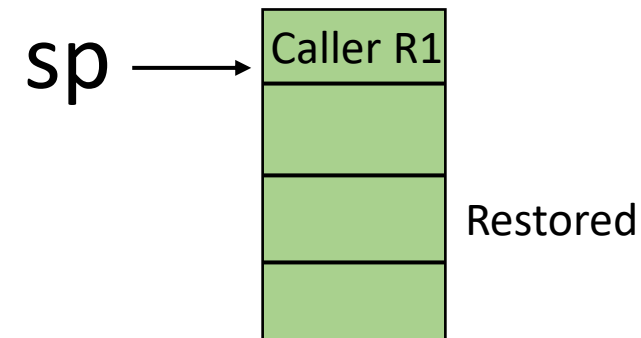
Save:

addi \$sp, \$sp, -4 → 32 bit registers, 4 bytes, one word, remember
sw R4, (\$sp)



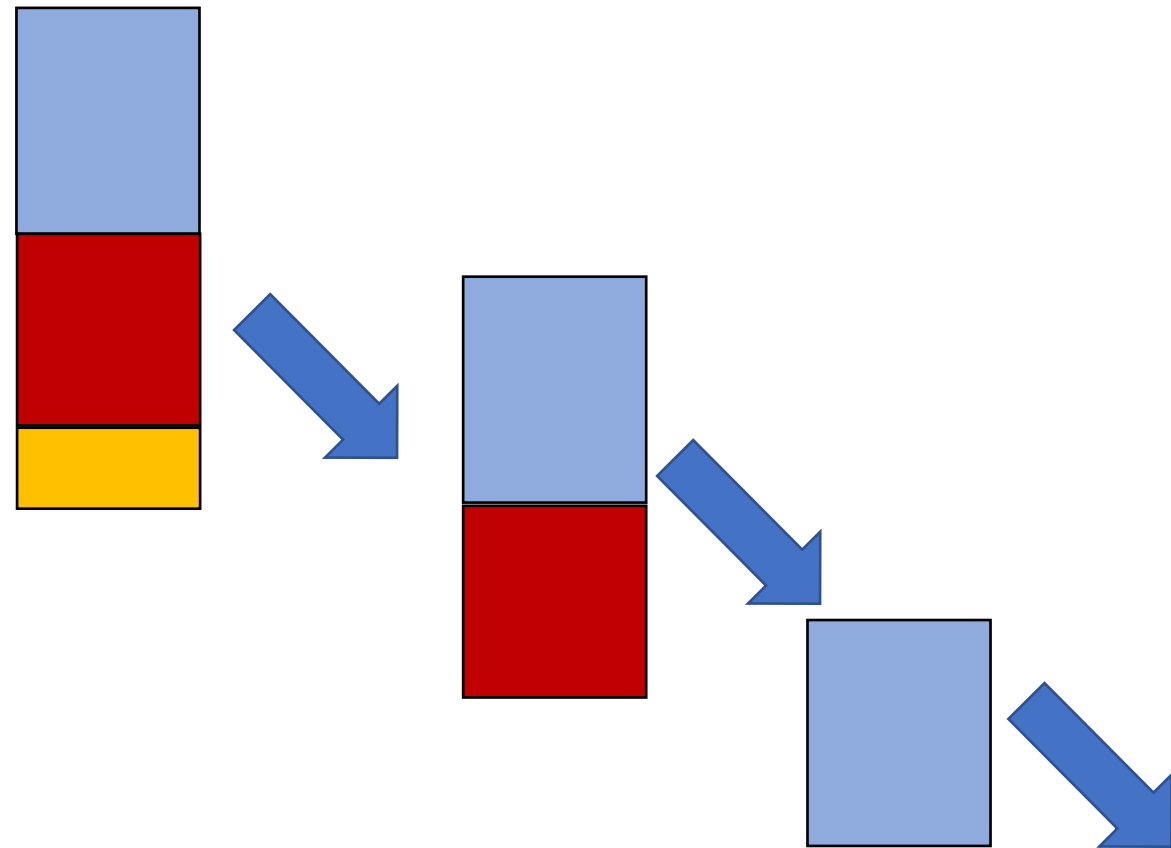
Restore:

lw R4, (\$sp)
addi \$sp, \$sp, 4



Nested Functions (Remember main() is a function too 😊)

```
CS305 // jal cs305
{
  CS405 // jal cs405
  {
    CS505 // jal cs505
    {
      } //jr
    } //jr
  } //jr
} //jr
```

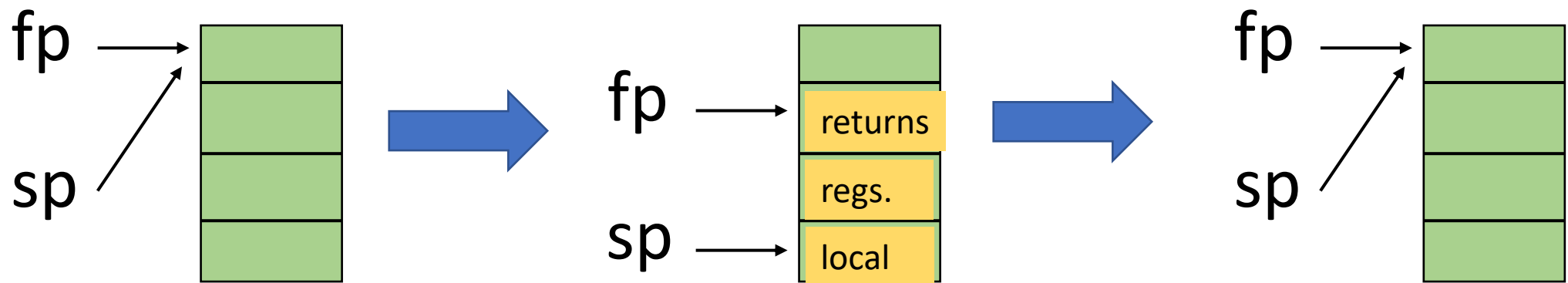


The final one: Frame pointer

Stack also stores local variables and data structures (local arrays and structures) for a function along with the **return address(es)**.

The final one: Frame pointer

Frame pointer: Points to local variables and saved registers. Points to the **highest address** in the **procedure frame**. **Stays there** throughout the procedure. Stack pointer, **moves** around.



Awesomeness: You can access any using fp/sp and an offset

Try This Out! Discuss on Piazza

Page no A-27 to A-29 P&H

Recursive function fact(n)

Look for sp, fp, ra, jal, and jr

More from the TAs during the labs.

For the Curious Ones (Beyond CS305)

Stack buffer overflow - 101:

https://en.wikipedia.org/wiki/Stack_buffer_overflow

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