### A MAC for WiFi Mesh Networks

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# **IEEE 802.11 (WiFi)**

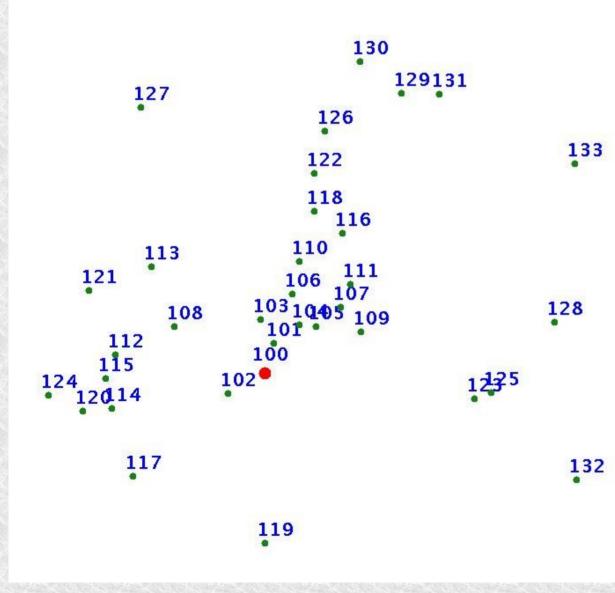
- Originally designed as a replacement (extension) for wired LANs
- Multi-hop mesh networks are very popular
  - Wireless community networks in cities
    - Share wired Internet broadband connections
    - Replacement for last-hop telephony?
  - Wireless Internet in rural areas
    - Developed as well as developing countries

# A WiFi Network in Djurslands, Denmark

- www.DjurslandS.net
- Remote area of Denmark
  - No broadband
  - Operators do not see returns to investment
  - 802.11 mesh network used for broadband connectivity

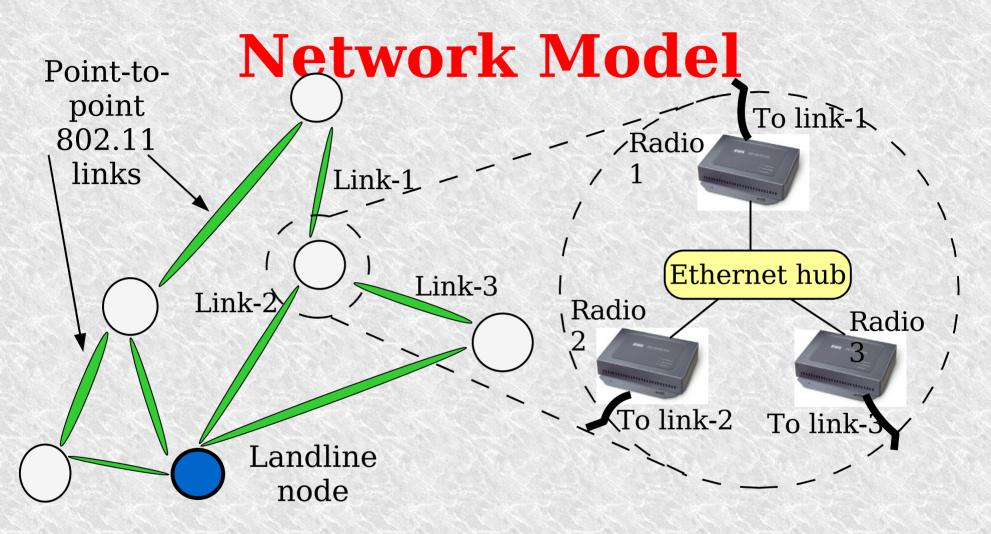


# A WiFi Network (planned) for Bhimavaram, A.P., India



#### Requirements:

- Need 384 Kbps video for telemedicine and educational apps
  - This is the requirement pervillage
- Cost should be as low as possible
  - So that the network model is replicable
- WiFi is the only serious option

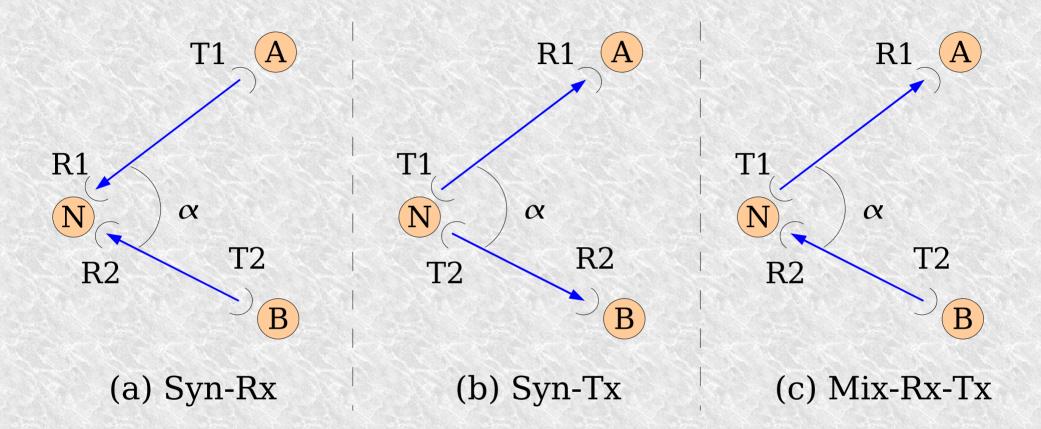


- Point-to-point links
- Multiple interfaces (radios) per node
- One directional antenna per link

# **Single Channel Operation**

- 802.11b has only three independent channels
- 802.11a has twelve independent channels
  - Four are meant for outdoor use
- Why only a single channel for the mesh?
  - Mitigation of "RF pollution"
  - The mesh may not be 3-edge-colourable
  - If the frequency is licensed, more channels could imply more cost

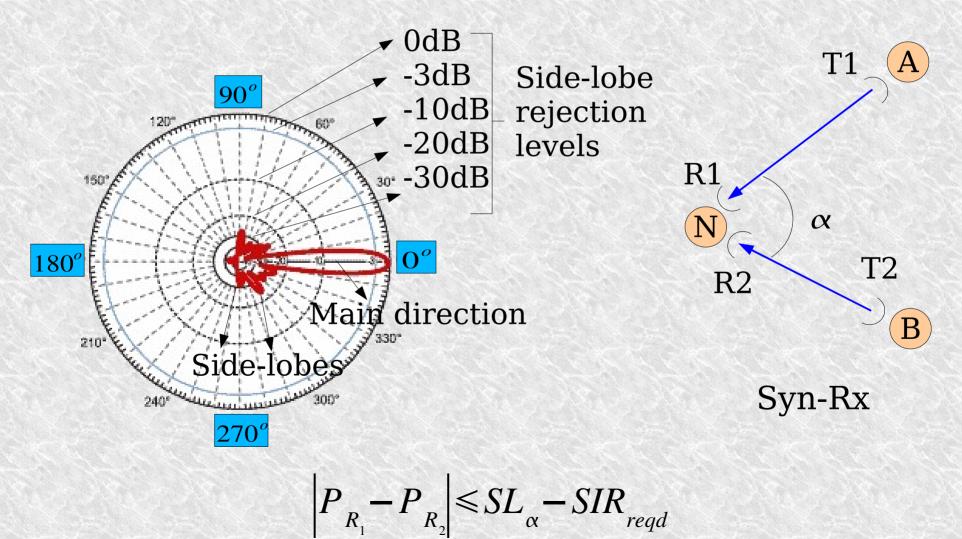
## SynRx, SynTx, and Mix-Rx-Tx

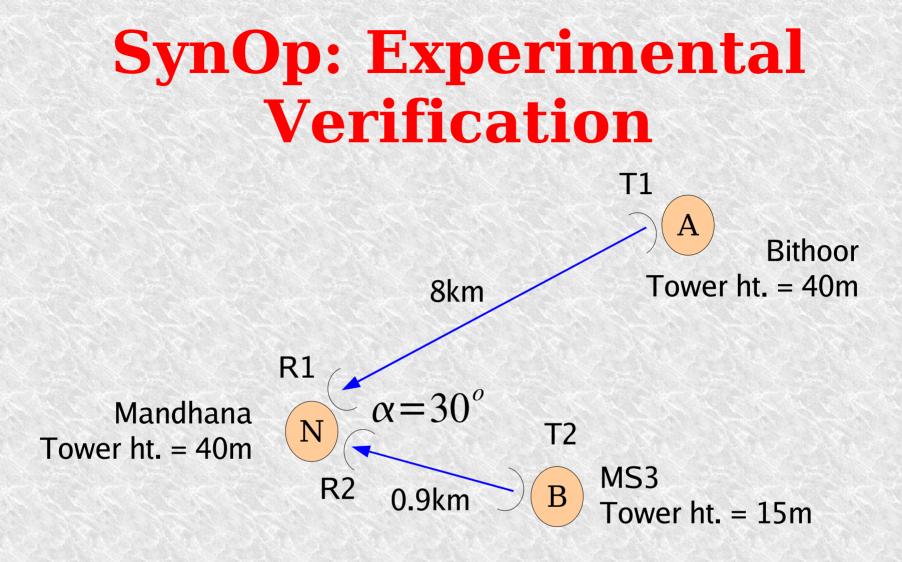


Exposed interface problem within a node: CSMA/CA (802.11 DCF) inherently allows only one link operation per node *Problems: (a) Immediate ACK, (2) CS back-off* 

# SynOp: SynRx + SynTx

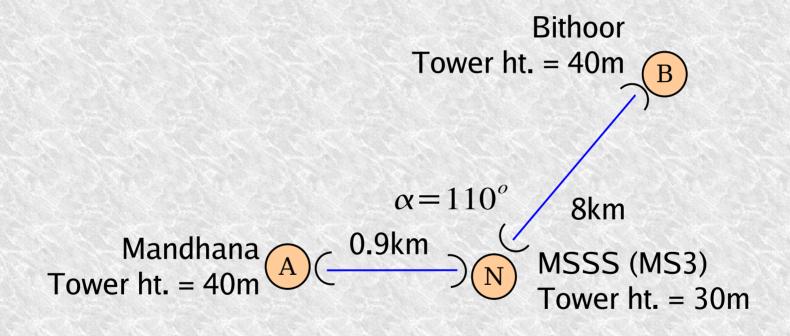
- Links at a node operating simultaneously, synchronously (on the same channel)
- Is this feasible? Yes, under certain conditions





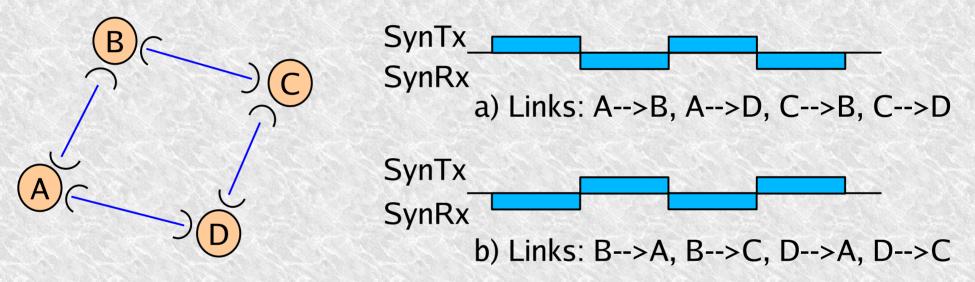
Used broadcast packets on both links (SynRx, SynTx) 6.5 Mbps with and without simultaneous operation SynTx also verified – using antenna diversity for the setup Experiments along with: A. R. Harish & Sreekanth Garigala

## SynOp: Another Experiment



## The 2P MAC Protocol

- Two phases: each node switches between SynRx and SynTx
- Topology has to be bipartite

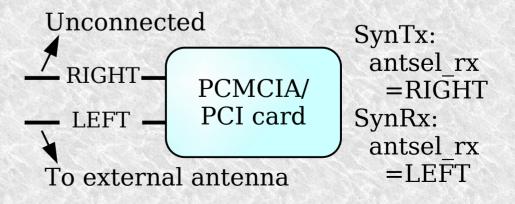


Note: diagram ignores system and propogation delays

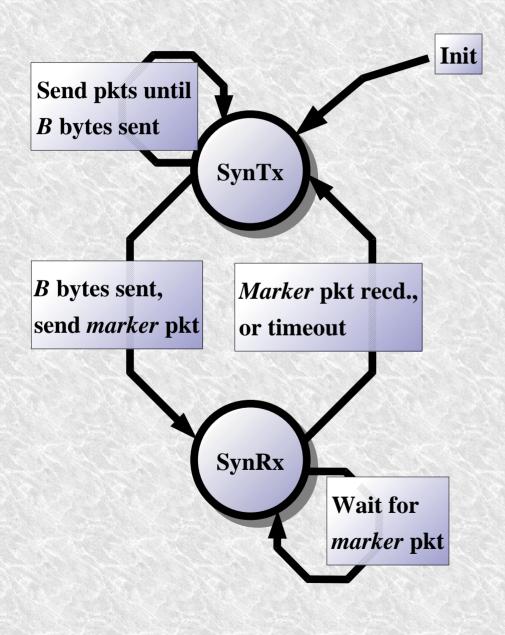
- How to achieve 2P on off-the-shelf hardware?
- Can 2P work without tight time synchronization?
- Relation between 2P and network topology
- 2P performance versus CSMA/CA

# **Achieving SynOp**

- Goal: bypass DCF to achieve SynOp
- Two offending factors: immediate ACKs, CSMA backoff
- Avoiding immediate ACKs:
  - Use IBSS mode
  - IP unicast to/from MAC broadcast
- Avoiding CSMA backoff:
  - Make use of diversity antenna
  - Change antsel\_rx to the unconnected antenna before transmitting

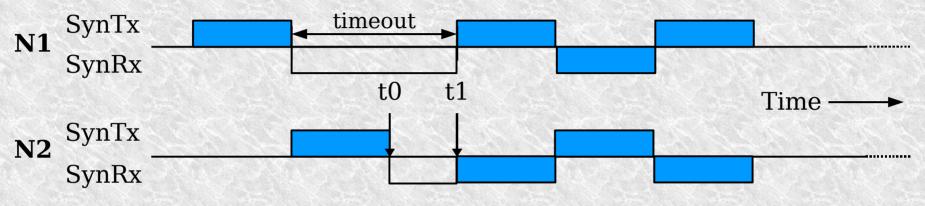


# **2P on a Single P2P Link**



- B bytes in each phase
- SynTx+SynRx = one round
- Marker packet acts as a "token"
- The two ends of the link are in loosesynchrony
- How do we handle:
  - Temporary loss of synchrony?
  - Link recovery or initialization?

## **The 2P Timeout Mechanism**

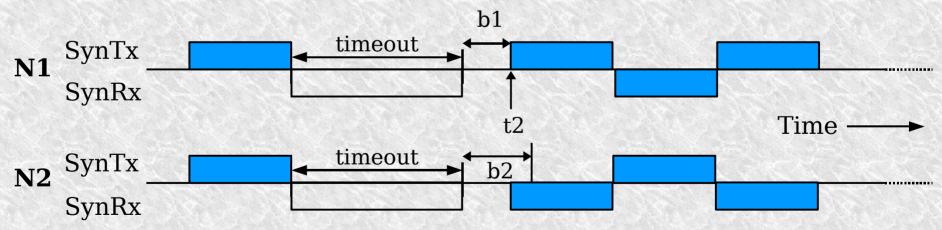


Note: diagram ignores system and propagation delays

- Timer started on entering SynRx
- Put on hold on starting to hear
- Link-resync takes only one round
- CRC errors of non-marker pkts immaterial

## **Bumping to Avoid Repeated Timeouts**

- If SynTx phases coincide, repeated timeouts occur
- Use random delay **bumping** to avoid this



Note: diagram ignores system and propagation delays

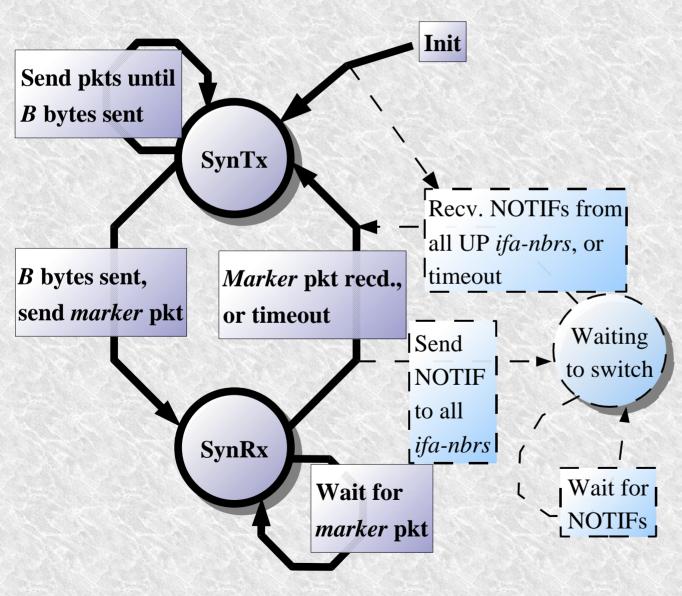
## **Parameters in 2P**

- Phase duration: B bytes
  - Large B implies lower % overhead, but higher latency
  - For B=10KB, 6% overhead, 13ms latency
  - For B=4.5KB, 11% overhead, 6ms latency

#### • Timeout:

- Lower bound: one phase duration
- Simulation: 1.25 times the phase duration
- Implementation: 25ms (kernel jitter ~ 10ms)

# **Communication Across Interface-Neighbours**



- NOTIF msgs to indicate end of SynRx
- Wait for NOTIF
  msgs from all ifa nbrs before
  SynTx
- UP/DOWN state w.r.t. each ifa-nbr
- Communication through sharedmemory, or ethernet

## **Some Remarks on 2P**

- Dummy bytes sent when no IP data
  - Power consumption not a major concern
  - Embedded platform ~ 4-6W at least 802.11 radio ~ 0.1-0.2W only
- Unequal phase durations possible
  - But not really useful for more than a single hop network
- RF leakages: not too many interfaces can be placed close to each other

# **Topology Constraints**

- 2P has two main constraints:
  - Topology should be bipartite
  - Power constraints

### **Power Constraints**

- Denote by  $P_i$ , the txpower at antenna  $A_i$
- Each transmission acts as interference to all other transmissions
- Write a set of linear equations with variables Pi
  - SIR >= SIR  $_{reqd}$
  - Probably should have some head-room too
  - Feasibility of a solution to this implies that the topology is 2P-compatible

a B Overall gain from  $a_i$  to  $a_i =$ (Gain of  $a_i$ 's Tx in  $a_i$ 's dirn)×  $(Gain of a_{i}' s Rx in a_{i}' s dirn) =$ Gain at angle  $\alpha \times$  Gain at angle  $\beta$ 

# **Topology Formation**

- Tree topology:
  - Trivially bipartite
  - Only one landline ==> tree is natural
    - Only a tree is active at any time
- Heuristics:
  - H1: use short links
  - H2: avoid short angles between links
  - H3: minimize the number of hops
- Mimic a natural deployment pattern
  - Nodes close to landline connected first, then the next level

# **Evaluation of 2P**

- Topology formation
- Simulation studies
- Implementation

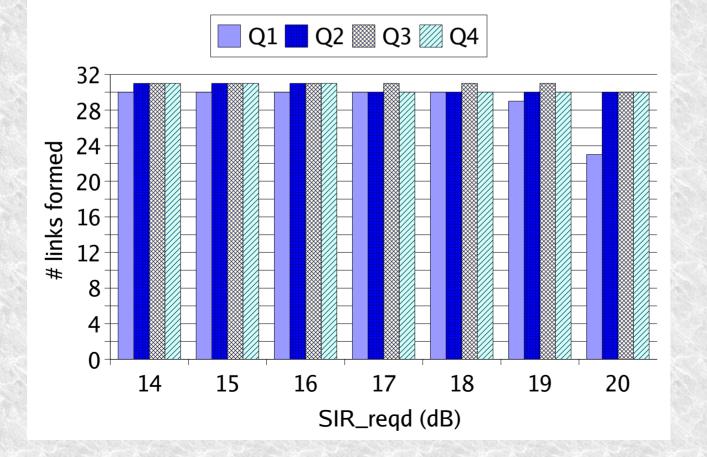
## **Evaluation of Topology Creation**

- Aspects of interest:
  - How well does the algorithm scale?
  - How much head-room in SIR<sub>reqd</sub> is possible?
- Evaluation:
  - Using parts of the map of Durg district, Chattisgarh, U.P.
  - Using random topologies

## **Topology Creation on Durg District**

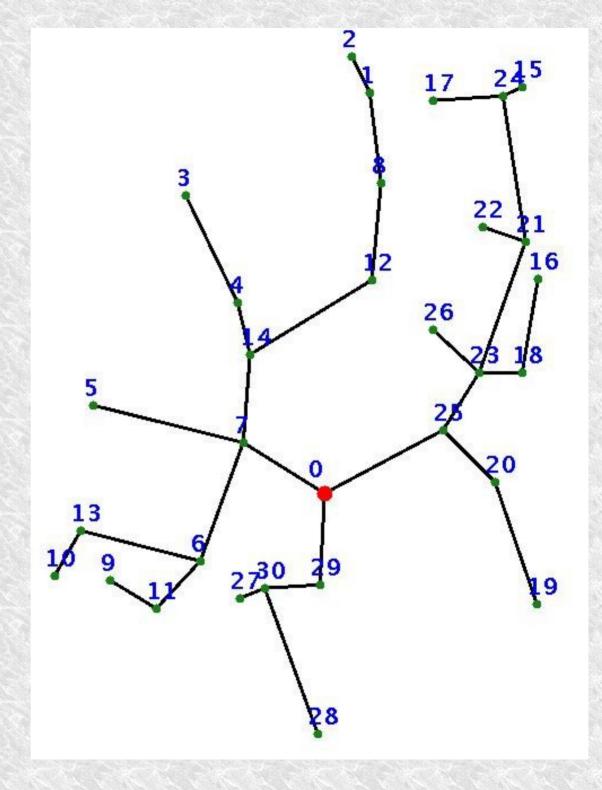
Four clusters of villages

→ Q<sub>i</sub> (i=1..4) 31, 32, 32, and 32 villages each

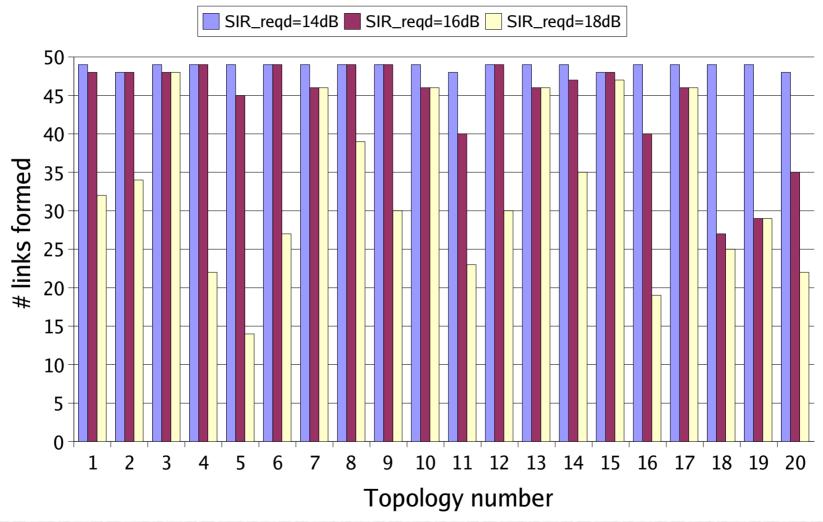


SIR<sub>reqd</sub> of 18-20dB easily possible

# The Topology on Q<sub>1</sub>



## **Topology Creation on Random Scenarios**



SIR<sub>reqd</sub> of 16-18dB mostly possible for up to 30-50 node topologies

## Simulation-based Evaluation

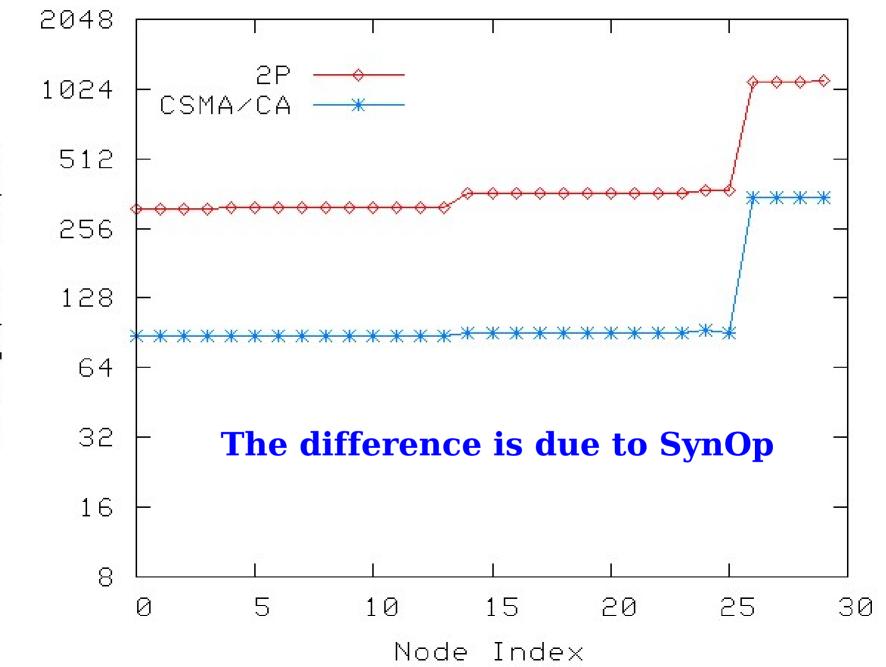
#### TeNs:

- http://www.cse.iitk.ac.in/~bhaskar/tens/
- Channel interference, grey regions, multiple interface support, directional antennas
- Further extensions:
  - Populating the ARP table appropriately
  - 24dBi directional antenna support
  - MAC modifications: air propagation delay, ACK timeout
  - LLC: sliding-window protocol

## **Simulation Parameters**

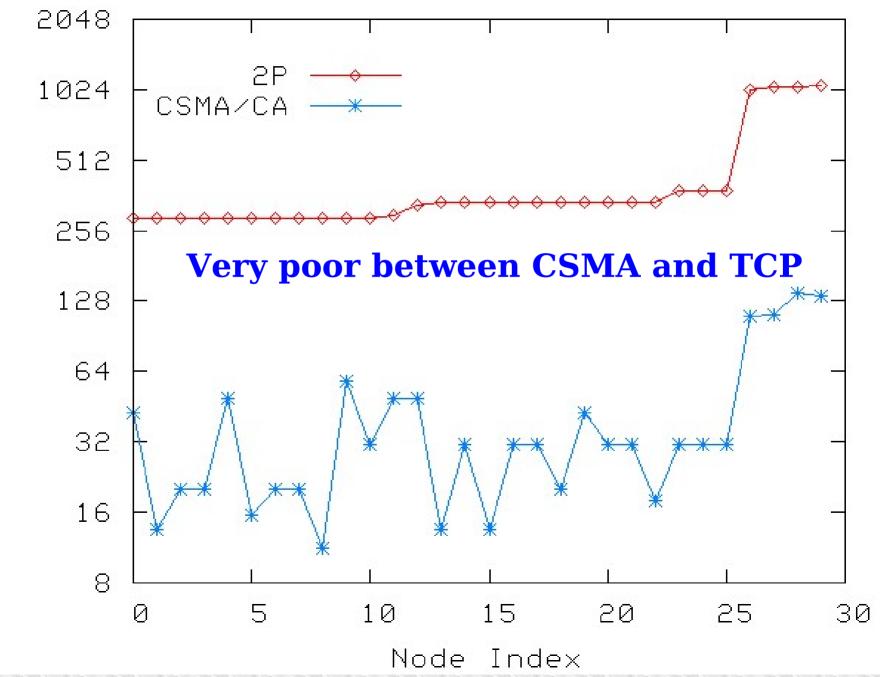
- $Q_1$ 's 31-node topology used
- UDP or TCP traffic
  - Packet size: 1400 bytes
  - UDP: saturating CBR traffic (every 2ms)
  - TCP: NewReno used
- Simulated time duration: 10sec

## **Saturation Throughput (UDP)**



Throughput (kbps)

#### **TCP Performance**



Throughput (kbps)

## Implementation-based Evaluation

- Implementation using HostAP v0.2.4, Linux 2.4 (also works on Linux 2.6)
- 2P on a single link: 6.1Mbps
  - Less than the max. possible 6.5Mbps
  - Overhead in antsel\_rx, marker pkt, CW<sub>min</sub> being 32
- 2P performance on a pair of links:

A <--> N1, N2 <--> B, UDP traffic

100 0 0 0 00 A				Avg (SD) thrpt at N2 (Mbps)	
10000	2P	2.70 (0.31)	2.06 (0.24)	2.81 (0.15)	2.81 (0.10)
No. all	CSMA	2.07 (0.13)	1.13 (0.22)	1.90 (0.15)	3.11 (0.14)

# **Concluding Remarks**

- 2P good for 802.11 mesh networks
  - Reuse of spectrum for max. throughput
  - Applicable in a wide-range of deployments
- Can be extended to P2MP scenarios as well
  - Provided the antenna is suitable
- Topology creation is an interesting aspect of study