

# CS758: Advanced Tools from Modern Cryptography

## *Lattice-Based Cryptography*

Lecture 01 (06/Jan/26)

Instructor: Chethan Kamath

# Administrivia

- When and where: Slot 10 (14:00-15:25, Tuesdays and Fridays), CC101
- Contact hours: drop by my office (CC305) any time!
- Teaching assistant: Priyanshu Singh (24M2101)



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- Announcements/online discussion on Moodle:
  - [moodle.iitb.ac.in/course/view.php?id=9123](https://moodle.iitb.ac.in/course/view.php?id=9123)
- Any volunteers for class rep?
  - Please set up WhatsApp/Signal group for coordination

# Administrivia..

- Grading Scheme

<b>Weightage</b>	<b>Towards</b>
30%	End-sem
25%	Mid-sem
20%	Paper presentation (two students per paper)
15%	Two quizzes
5%	Class participation
5%	Scribing (1-2 lectures per student)

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15% +10	Two quizzes
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- Please submit **scribe** notes within two weeks
- Will **scrap paper presentation** and readjust grades if strength  $> 20$
- Attendance is not mandatory (but encouraged)

## Administrivia..

- Resources can be found on course website
  - [cse.iitb.ac.in/~ckamath/courses/2026s/CS758.html](http://cse.iitb.ac.in/~ckamath/courses/2026s/CS758.html)

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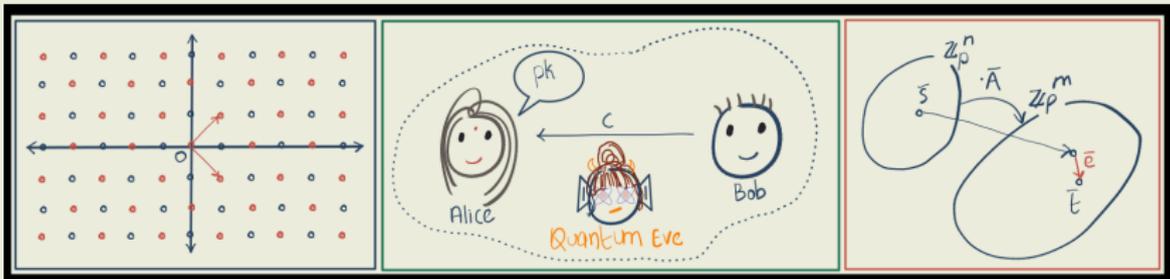
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## Pre-requisites

-  Will assume basic cryptography (CS409/CS409m/CS783)
-  Brush up your linear algebra
-  Useful to know basic coding theory and complexity theory
  - TA will conduct recitation sessions for these topics



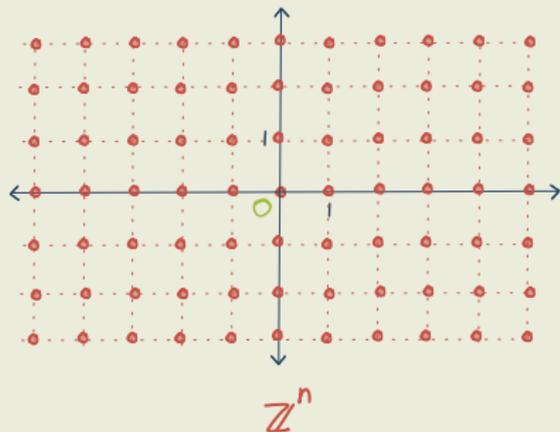
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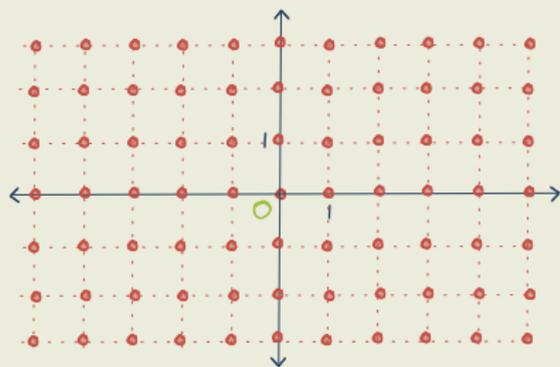
# What are Lattices?



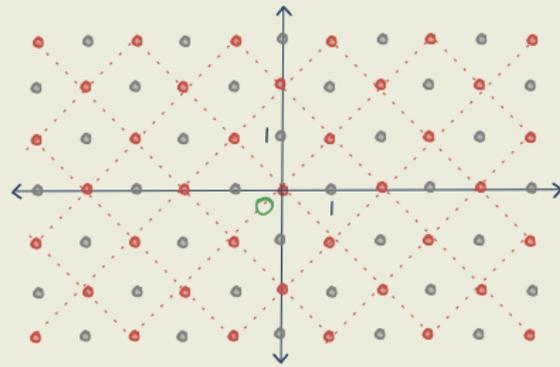
## Informal Definition 1

*Discrete* subspace of (vector space)  $\mathbb{R}^n$

# What are Lattices?



$\mathbb{Z}^n$

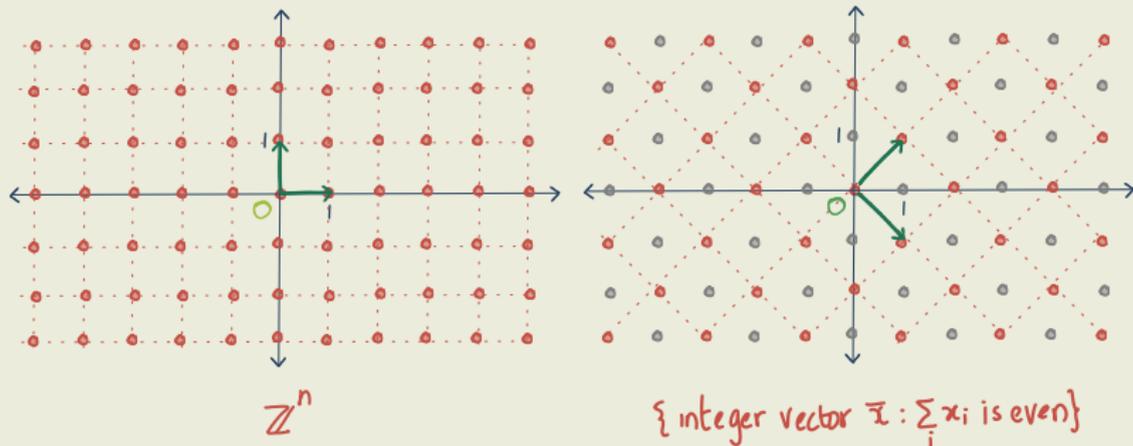


{ integer vector  $\vec{x} : \sum_i x_i$  is even }

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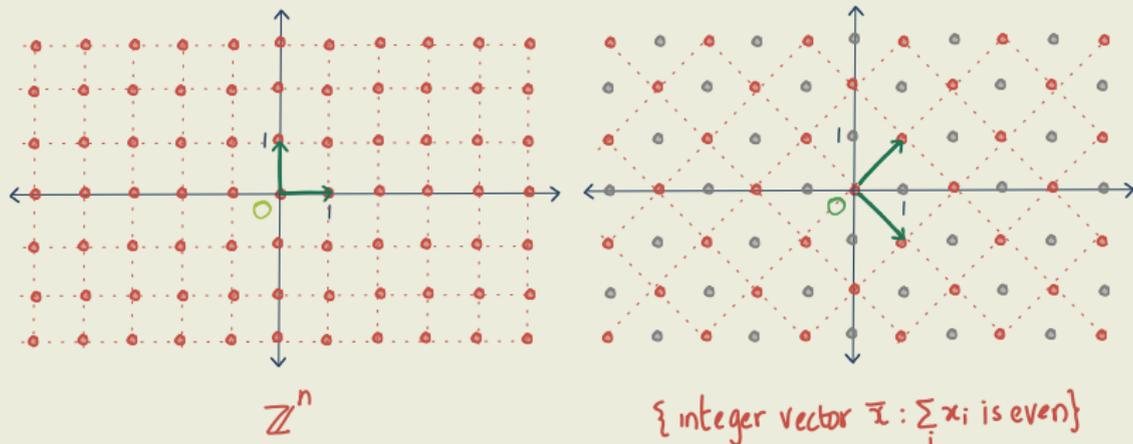
## Informal Definition 1

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## Informal Definition 2

Integer linear combinations of linearly-independent vectors  $\bar{b}_1, \dots, \bar{b}_m \in \mathbb{R}^n$

# What are Lattices?



## Informal Definition 1

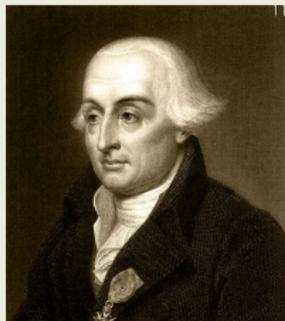
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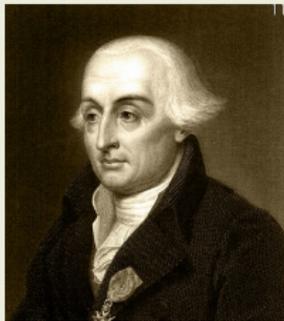
rank  $\rightarrow$  dimension

## How are Lattices Useful?...



- Historical applications
  1. Tool to study pure mathematics
    - E.g., as a bridge between number theory and geometry: can prove Fermat's theorem on sums of two squares using lattices

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Lagrange



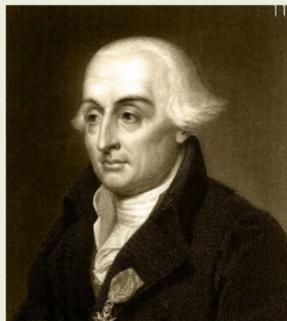
Gauß



Minkowski

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Minkowski



Bravais

- Historical applications
  1. Tool to study pure mathematics
    - E.g., as a bridge between number theory and geometry: can prove Fermat's theorem on sums of two squares using lattices
  2. Chemistry: to study crystal structure (crystallography)
- ...

## How are Lattices Useful?..



Coppersmith



Lagarias



Odlyzko

- Modern applications
  1. Coding theory
  2. Cryptanalysis: breaking schemes using lattice algorithms

## How are Lattices Useful?..



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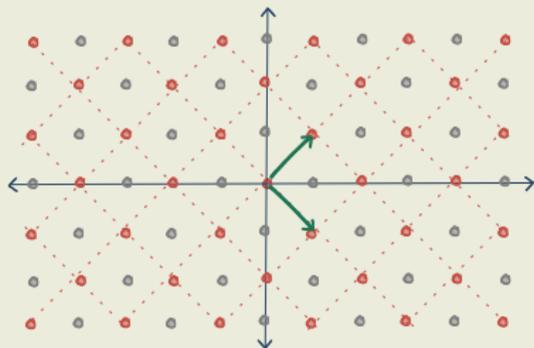
Ajtai



Regev

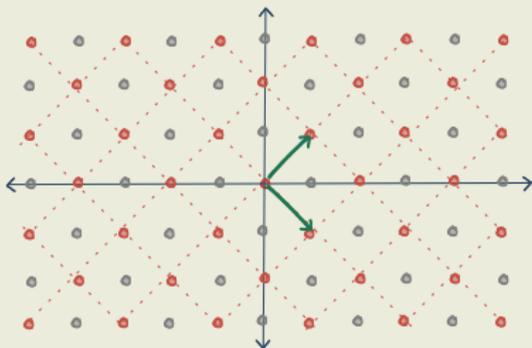
- Modern applications
  1. Coding theory
  2. Cryptanalysis: breaking schemes using lattice algorithms
  3. Lattice-based cryptography: constructing cryptographic primitives assuming hardness of lattice problems...

## How are Lattices Useful?..



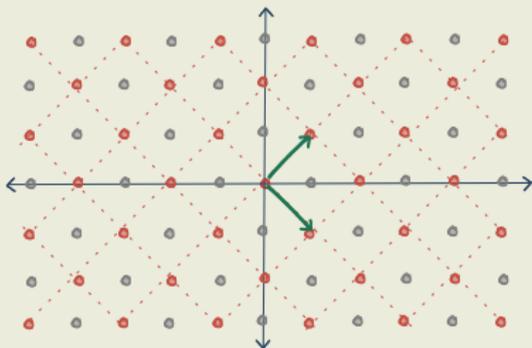
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  1. Fast (linear) operations ⚡

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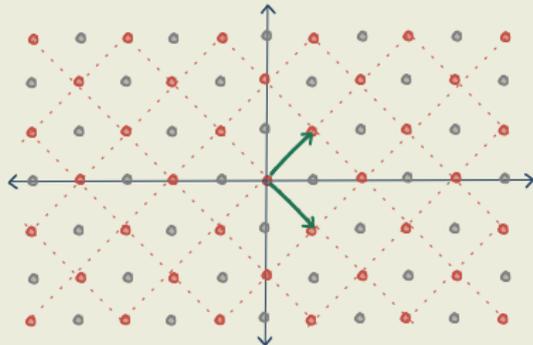
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    - ⇒ Cryptography based on worst-case assumptions! ★

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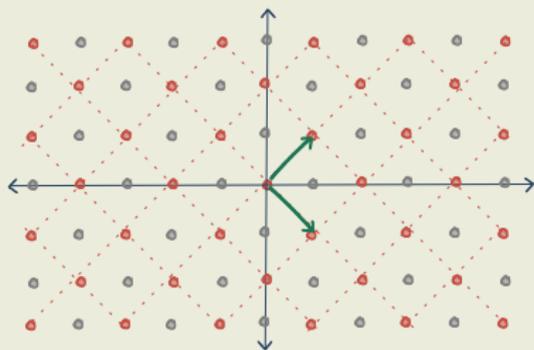
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    - ⇒ Cryptography based on worst-case assumptions! ★
  3. No efficient quantum algorithms known 
    - ⇒ Post-quantum cryptography
  4. Expressive! We know how to construct certain primitives only from lattices
    - E.g., Fully-Homomorphic Encryption (FHE) 

## CS758: Course Overview



1. Module I: Introduction to Lattices
2. Module II: Basic Cryptography from Lattices
3. Module III: Advanced Cryptography from Lattices

# Module I: Introduction to Lattices

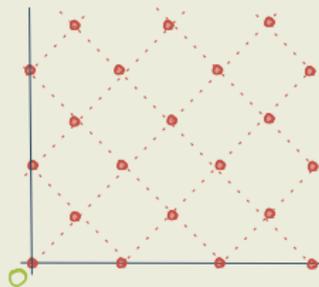
- Study computational problems on lattices

## Module I: Introduction to Lattices

- Study computational problems on lattices

### Shortest Vector Problem (SVP)

- $\mathbb{I}/p$ : lattice  $\Lambda$



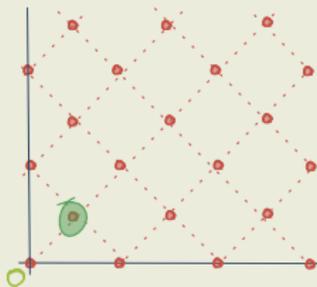
- Solution: *shortest* non-zero lattice vector  $\bar{x}$  in  $\Lambda$

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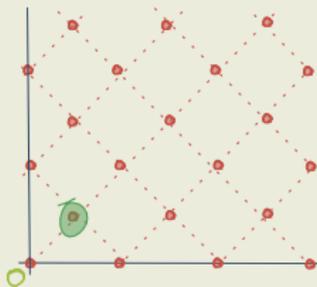
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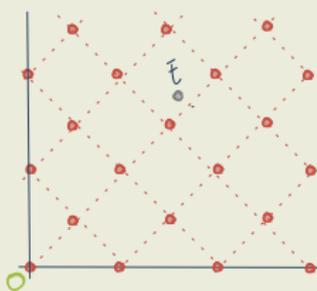
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## Closest Vector Problem (CVP)

- I/p: lattice  $\Lambda$ , target  $\bar{t} \in \mathbb{R}^n$



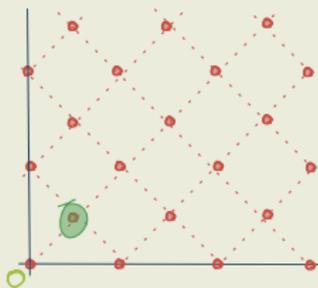
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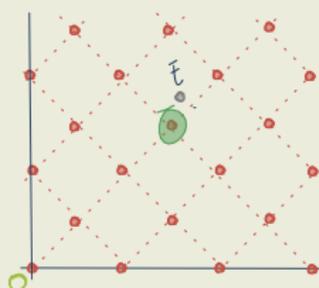
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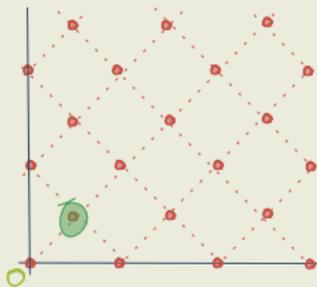
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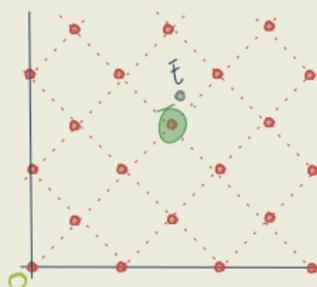
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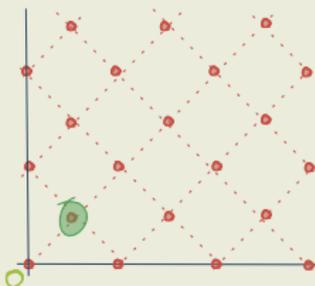
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- Solution: lattice vector  $\bar{x}$  *closest* to  $\bar{t}$

- Also look at *estimation* and *approximation* variants
- Most problems hard to solve in the *worst case* (**NP**-hard, in fact)

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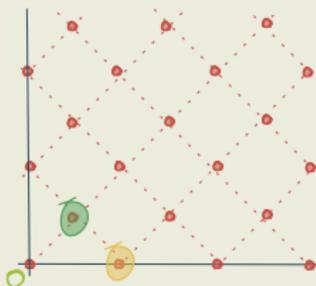


- Algorithms for lattice problems, e.g., LLL algorithm

### Informal Theorem 1 ([LLL82])

There exists a polynomial-time algorithm that solves SVP up to an *exponential* (in  $n$ ) approximation factor

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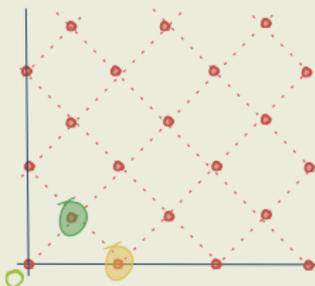


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Cryptanalysis using lattice algorithms:

- Factorise integers
- Solve low-exponent RSA
- Break Linear Congruential Generators (LCG)

## Module II: Basic Cryptography from Lattices.

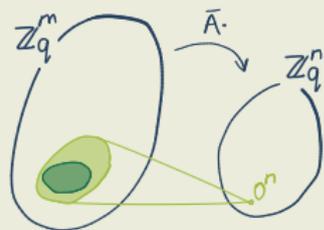
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  - E.g., integer factorisation, discrete-logarithm problem (DLP)

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- Lattice-based average-case hard problems (think  $m > n$  below):

### Short Integer Solution (SIS)

- I/p: matrix  $\bar{A} \leftarrow \mathbb{Z}_q^{n \times m}$



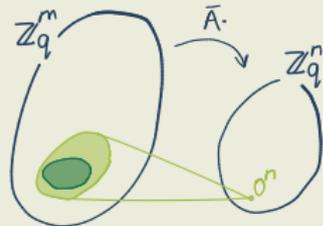
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such that  $\bar{A}\bar{x} = 0^n \pmod q$

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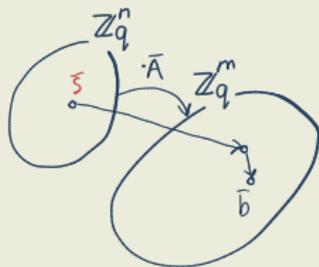
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### Learning With Errors (LWE)

- I/p: matrix  $\bar{A} \leftarrow \mathbb{Z}_q^{n \times m}$  and  $\bar{b}^t \approx \bar{s}^t \bar{A}$ , where  $\bar{s} \leftarrow \mathbb{Z}_q^n$



- Solution:  $\bar{s}$

## Module II: Basic Cryptography from Lattices.

- What is the connection to lattices?

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- What is the connection to lattices? Can be viewed as *average-case* variants of CVP

Informal Theorem 2 ([Ajt96] (resp., [Reg05]))

SIVP (resp., BDD) is worst-case hard  $\Rightarrow$  SIS (resp., LWE) is average-case hard

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### Informal Theorem 2 ([Ajt96] (resp., [Reg05]))

SIVP (resp., BDD) is worst-case hard  $\Rightarrow$  SIS (resp., LWE) is average-case hard

- What can we construct assuming SIS and LWE?

### Informal Theorem 3 ([Ajt96])

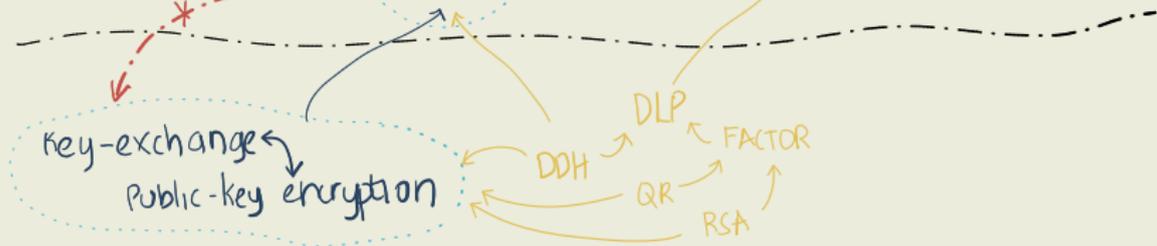
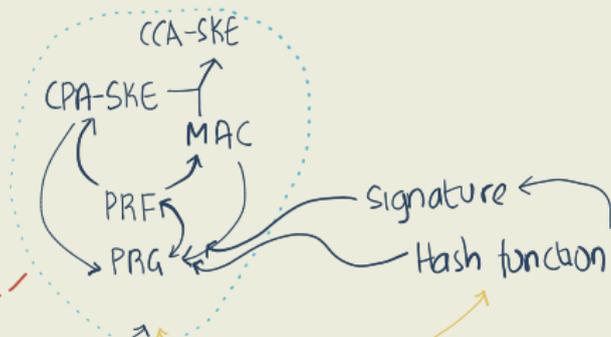
SIS  $\Rightarrow$  one-way function (OWF) and collision-resistant hash function

### Informal Theorem 4 ([Reg05])

LWE  $\Rightarrow$  public-key encryption (PKE)

# The Cryptographic Landscape

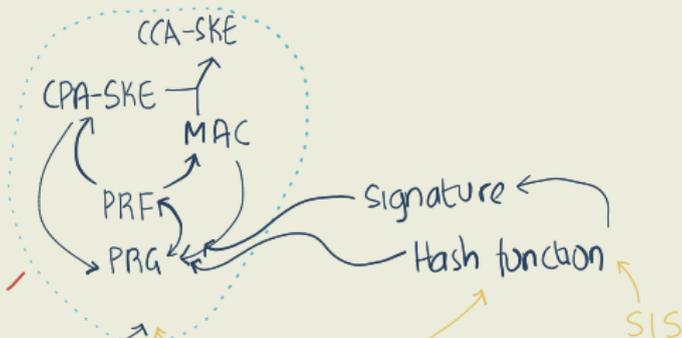
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"Cryptomania"

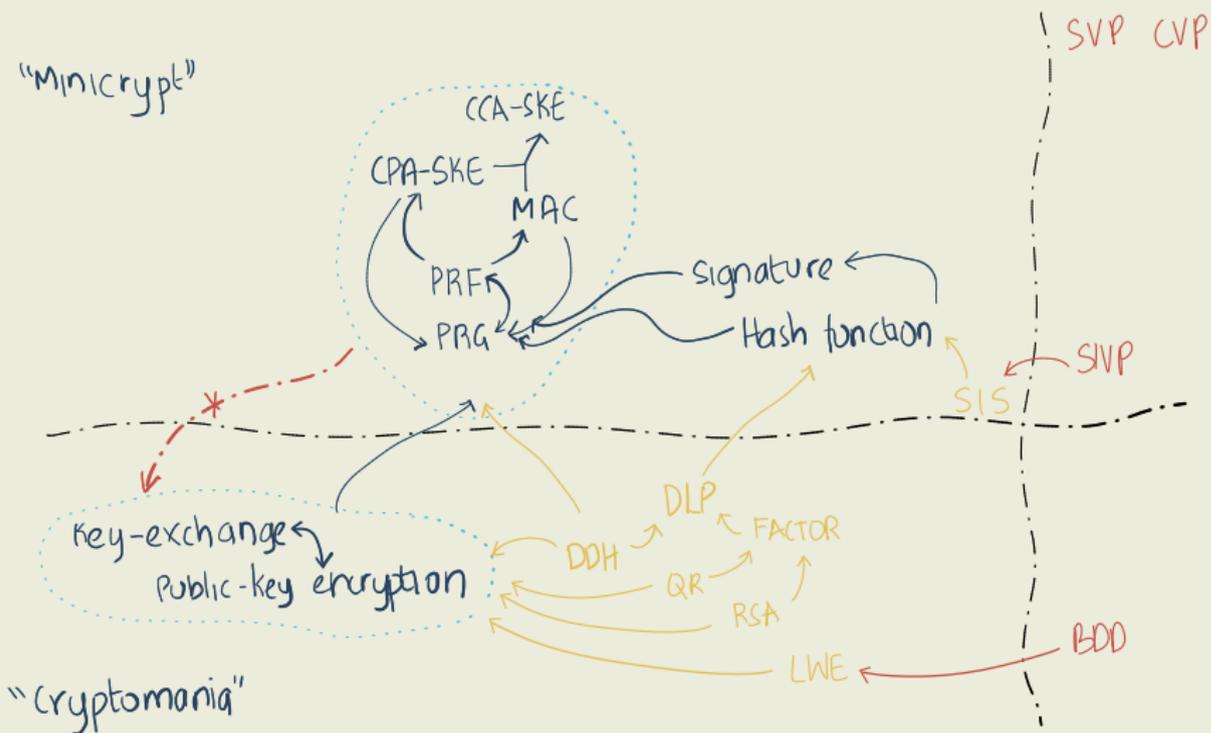
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# The Cryptographic Landscape



# Module III: Advanced Cryptography from Lattices..

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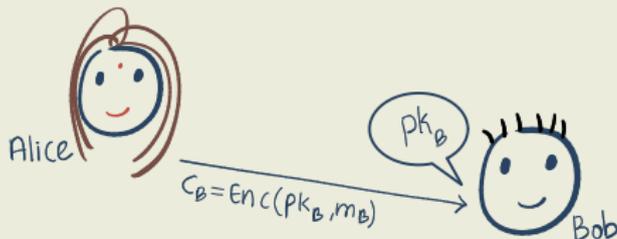
Advanced Primitive I: Identity-Based Encryption (IBE) [Sha84]

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## Advanced Primitive I: Identity-Based Encryption (IBE) [Sha84]

- Recall PKE (Public-Key Encryption)

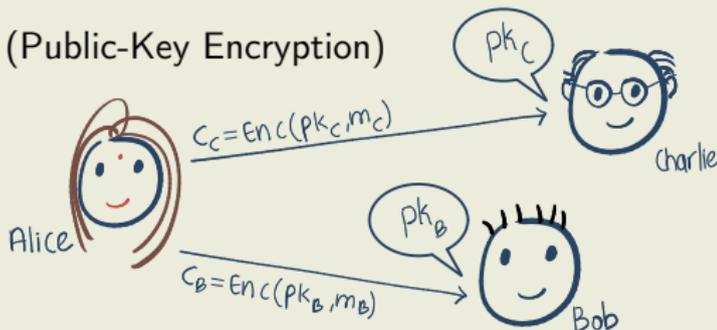


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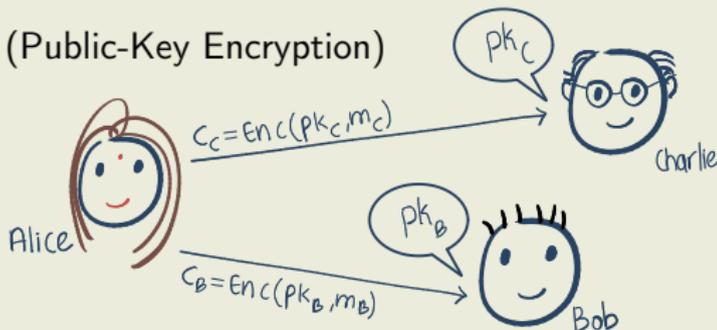
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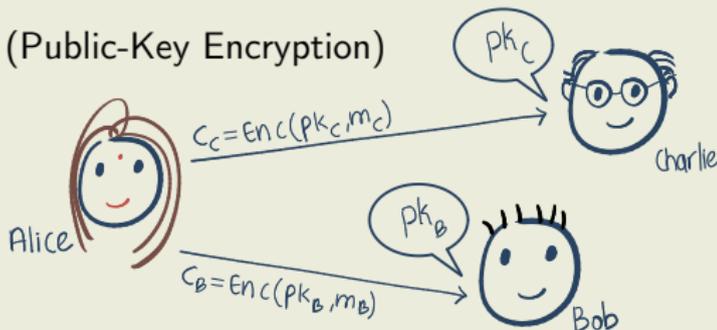


 **Drawback:** need to maintain (authenticated) directory with identity (e.g., e-mail address) and PK pairs

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## Advanced Primitive I: Identity-Based Encryption (IBE) [Sha84]

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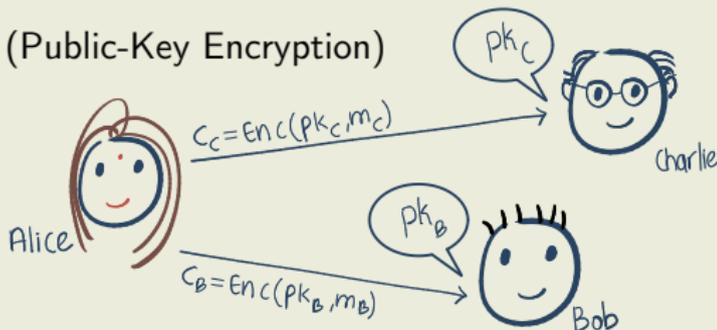
### Informal Theorem 5 ([GPV08])

LWE  $\Rightarrow$  IBE

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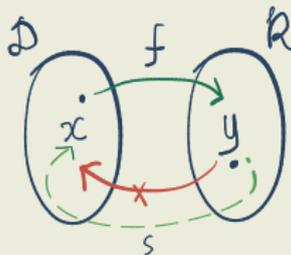
LWE  $\Rightarrow$  IBE

- Generalisation of IBE: attribute-based encryption (ABE) [SW04]
  - Application: secure access control (e.g., Cloudflare's GeoV2)

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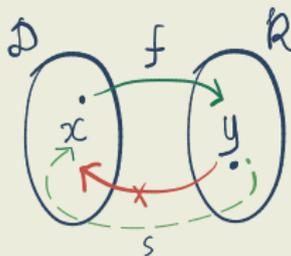
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# Module III: Advanced Cryptography from Lattices.

## Advanced Primitive I: Identity-Based Encryption (IBE) [Sha84]

- Recall trapdoor functions: OWF that is easy to invert given a “secret”  $s$



Key tool: lattice trapdoors, e.g., a “short” basis

### Informal Theorem 6 ([Ajt96; MP12])

There exists a PPT algorithm that samples a SIS matrix  $\bar{A} \leftarrow \mathbb{Z}_q^{n \times m}$  along with a “short” matrix  $\bar{S} \in \mathbb{Z}^{m \times m}$  such that  $\bar{A}\bar{S} = 0^{n \times m} \pmod{q}$

# Module III: Advanced Cryptography from Lattices..

Advanced Primitive II: FHE (Fully-Homomorphic Encryption) [RAD78]

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## Advanced Primitive II: FHE (Fully-Homomorphic Encryption) [RAD78]

- Recall secret-key encryption (SKE)

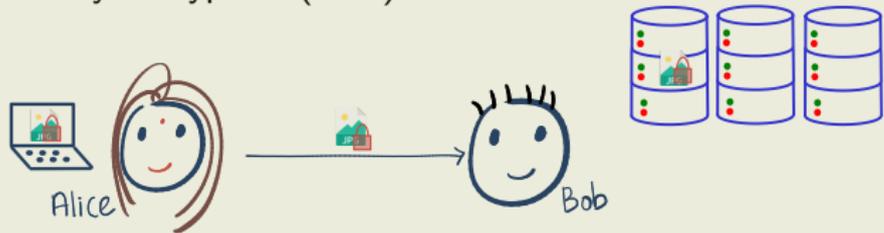


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## Advanced Primitive II: FHE (Fully-Homomorphic Encryption) [RAD78]

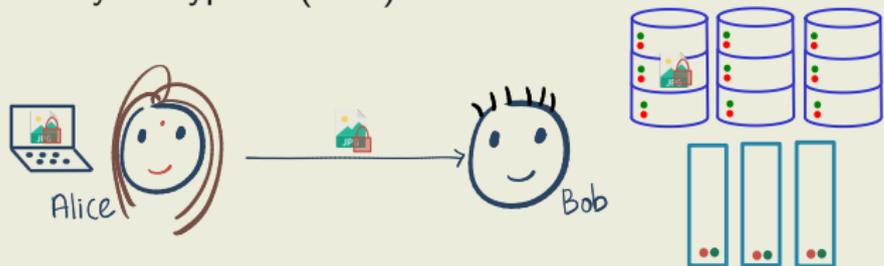
- Recall secret-key encryption (SKE)



# Module III: Advanced Cryptography from Lattices..

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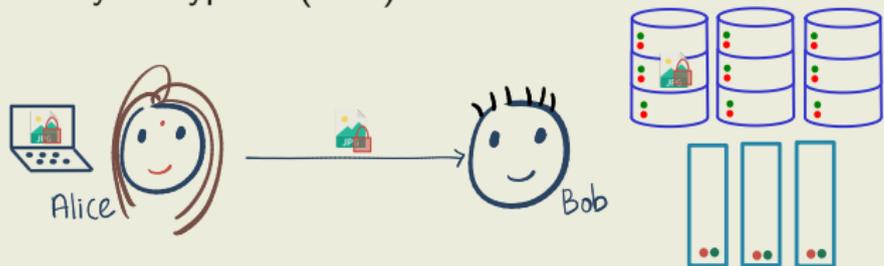
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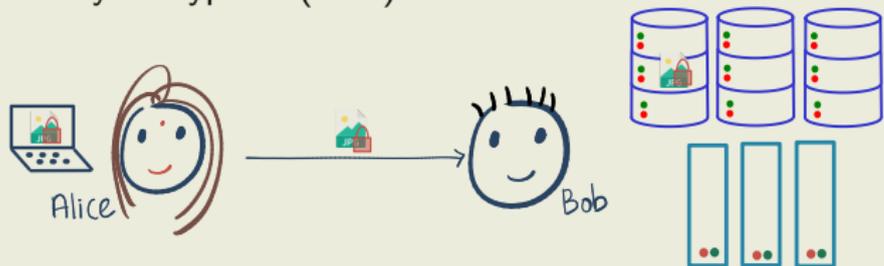


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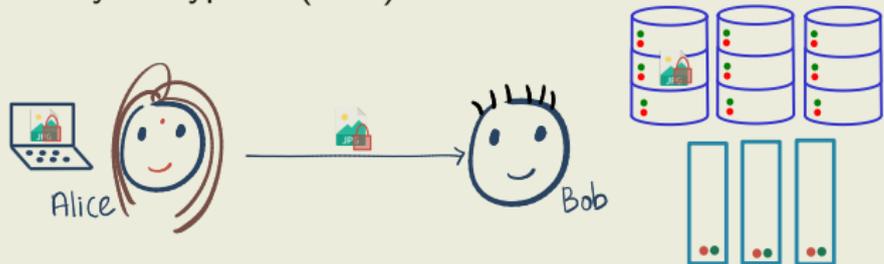
Informal Theorem 7 ([GSW13])

$LWE \Rightarrow FH(PK)E$

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### Informal Theorem 7 ([GSW13])

$LWE \Rightarrow FH(PK)E$

- Applications:
  - Private (zero-trust) outsourcing of computation
  - Encrypted search (e.g., see Zama)

## Lecture Resources

1. *Lecture 1: Introduction of CSE206A (Spring 2007)*, by Micciancio (PDF)
2. *Historical Talk on Lattice-Based Cryptography*, Micciancio (YouTube)
3. *Mathematics of Lattices*, Micciancio (YouTube)

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