Computer Programming

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Session: Analyzing Selection Sort

Quick Recap of Relevant Topics

- Selection sort
  - Intuition
  - C++ implementation

Overview of This Lecture

- Analyzing performance of selection sort
  - Counting “basic” steps in sorting an array of size n

Selection Sort Animated

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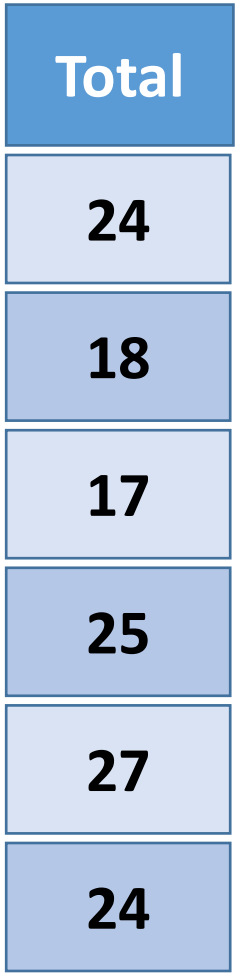

Selection Sort in C++

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```
int main() {
   … Declarations, input validation and reading elements of array A …
   // Selection sort
   int currTop, currMaxIndex;  // A[currTop] … A[n-1] is unsorted array
   for (currTop = 0; currTop < n; currTop ++) {
       currMaxIndex = findIndexOfMax(A, currTop, n);
       swap(A, currTop, currMaxIndex);
   }
    … Rest of code …
   return 0;
}
```

Selection Sort in C++

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```
// PRECONDITION: start < end
// start, end within array bounds of A
int  findIndexOfMax(int A[],  int start,  int end) {
    int i, currMaxIndex = start;
    for ( i = start ; i < end; i++ ) {
      if (A[i] >= A[currMaxIndex]) { currMaxIndex = i; }
    }
   return currMaxIndex;
 }
// POSTCONDITION: A[currMaxIndex] at least as large as
// all elements in A[start] through A[end-1], no change in A
```

Selection Sort in C++

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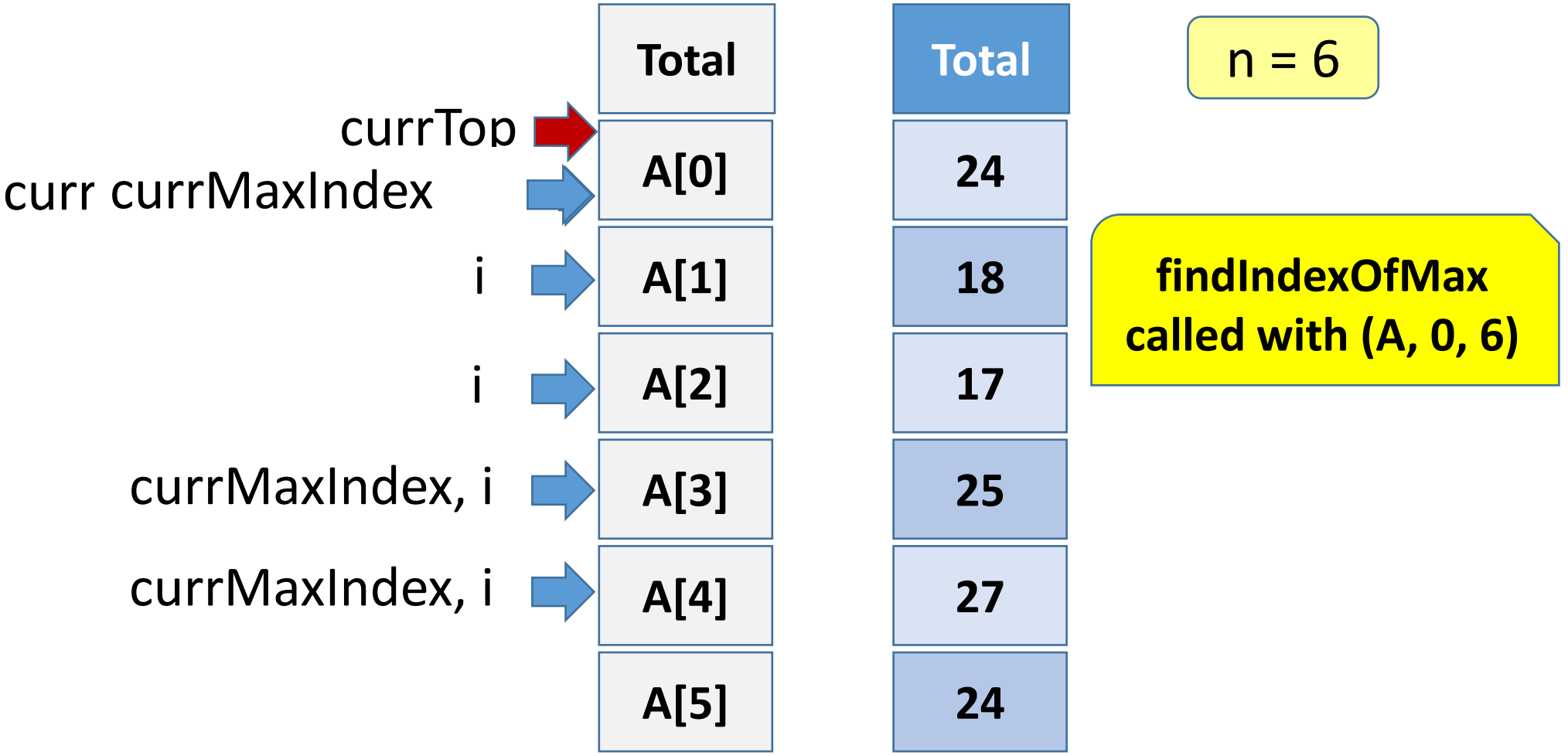
```
// PRECONDITION: index1, index2 within array
//                          bounds of A
void  swap(int A[],  int index1,  int index2) {
    int temp;
    temp = A[index1];
    A[index1] = A[index2];
    A[index2] = temp;
    return;
}
// POSTCONDITION: A[index1], A[index2] swapped
//                              Array A changed
```

“Basic” Steps in Selection Sort

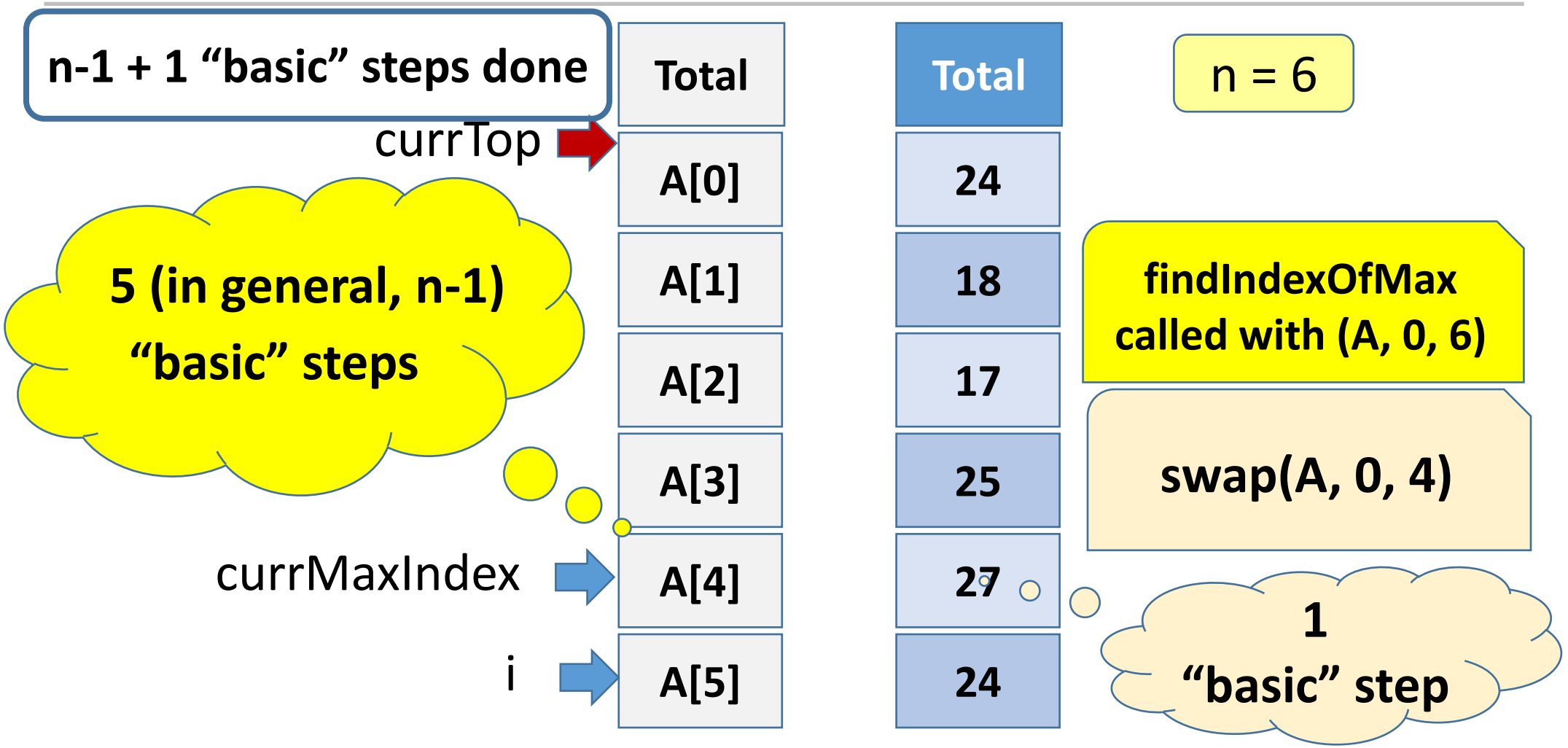
- Reading two elements of array A, comparing them and updating currMaxIndex, if necessary
- Swapping two specified elements of array A

Given an array of n integers, how many “basic” steps (as a function of n) are needed to sort by selection sort?

Counting “Basic” Steps In Selection Sort



Counting “Basic” Steps In Selection Sort



Recall: Selection Sort in C++

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```
int main() {
    … Declarations, input validation and reading elements of array A …
   // Selection sort
   int currTop, currMaxIndex;  // A[currTop] … A[n-1] is unsorted array
   for (currTop = 0; currTop < n; currTop ++) {
       currMaxIndex = findIndexOfMax(A, currTop, n);
       swap(A, currTop, currMaxIndex);
   }
     … Rest of code …
   return 0;
}
```

Counting “Basic” Steps In Selection Sort

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n-1 + 1 “basic” steps done

currTop

curr currMaxIndex

i

currMaxIndex, i

	Total
A[0]	27
A[1]	18
A[2]	17
A[3]	25
A[4]	24
A[5]	24

Total

n = 6

findIndexOfMax called with (A, 1, 6)

Counting “Basic” Steps In Selection Sort

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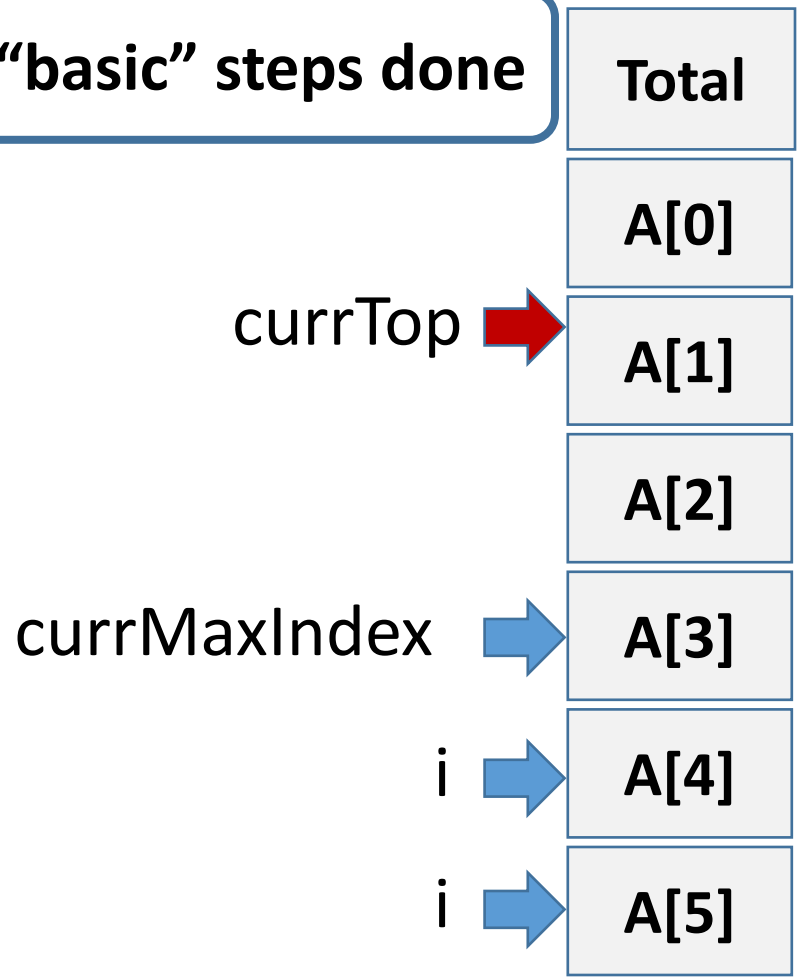
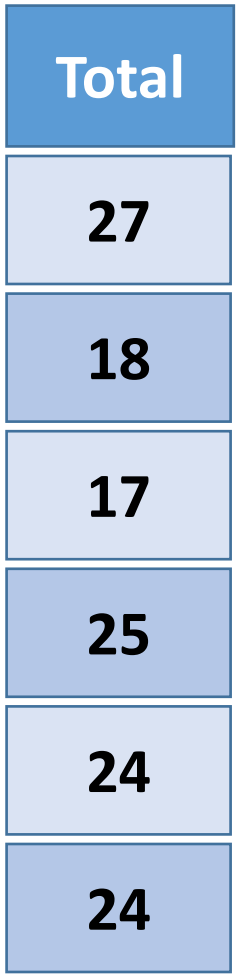
n-1 + 1 “basic” steps done





n = 6

findIndexOfMax called with (A, 1, 6)

Counting “Basic” Steps In Selection Sort

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n-1 + 1 “basic” steps done

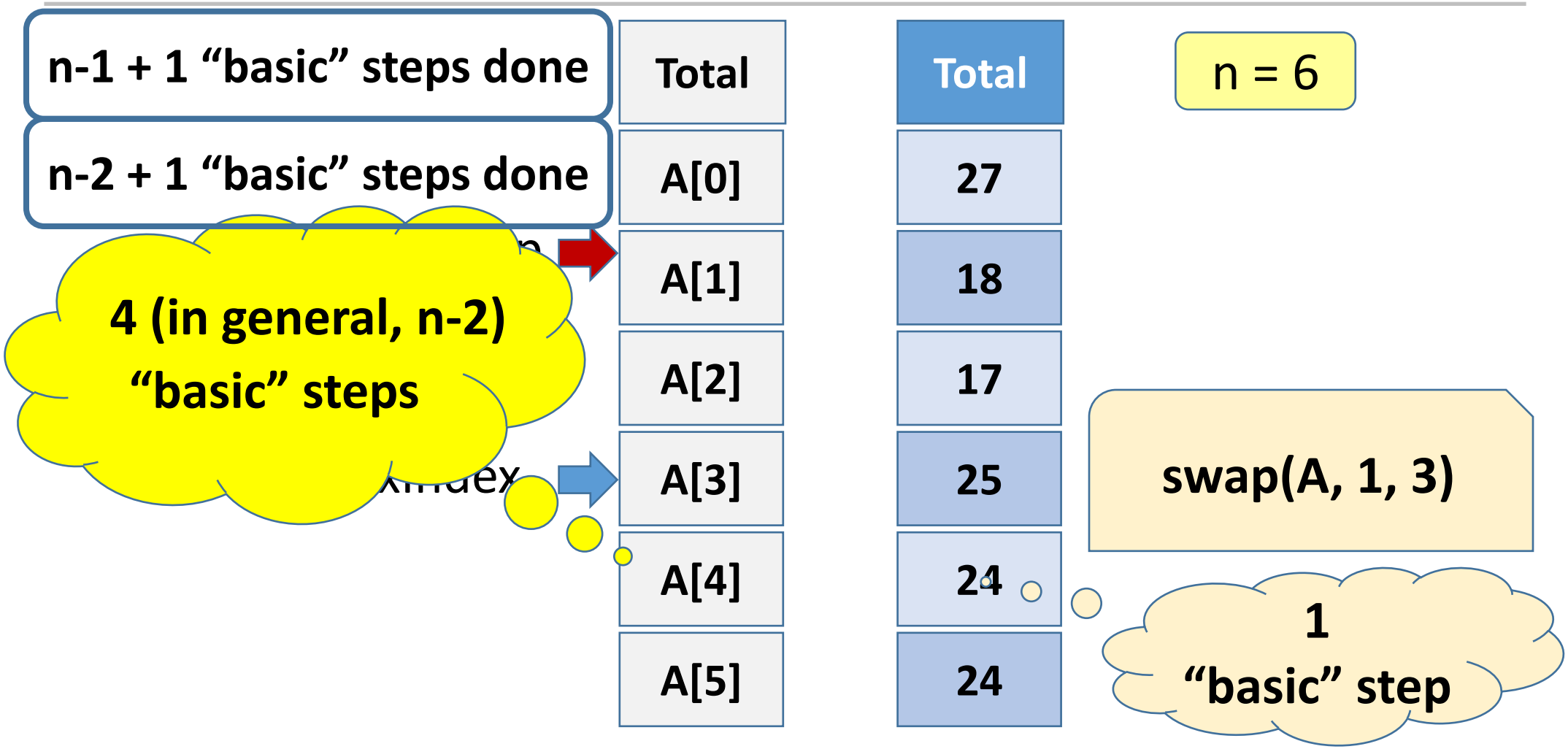
4 (in general, n-2) “basic” steps

	Total	Total
	A[0]	27
n	A[1]	18
	A[2]	17
xIndex	A[3]	25
	A[4]	24
i	A[5]	24

n = 6

findIndexOfMax called with (A, 1, 6)

Counting “Basic” Steps In Selection Sort



Counting “Basic” Steps In Selection Sort

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n-1 + 1 “basic” steps done

n-2 + 1 “basic” steps done

n-3 + 1 “basic” steps done

currTop

Total
A[0]
A[1]
A[2]
A[3]
A[4]
A[5]

Total
27
18
17
25
24
24

n = 6

Counting “Basic” Steps In Selection Sort

n-1 + 1 “basic” steps done

n-2 + 1 “basic” steps done

n-3 + 1 “basic” steps done

⋮

n-(n-1) + 1 “basic” steps done

currTop

Total	Total
A[0]	27
A[1]	18
A[2]	17
A[3]	25
A[4]	24
A[5]	24

n = 6

Counting "Basic" Steps in Selection Sort

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- Count of "basic" steps to sort an array of n elements:

(n-1 + 1) +

(n-2 + 1) +

⋮

(n-(n-1) + 1)

Increases quadratically with n

= (1 + 2 + … n-1) + n-1 = (n - 1) x (n+2)/2

Quadratic Growth With n

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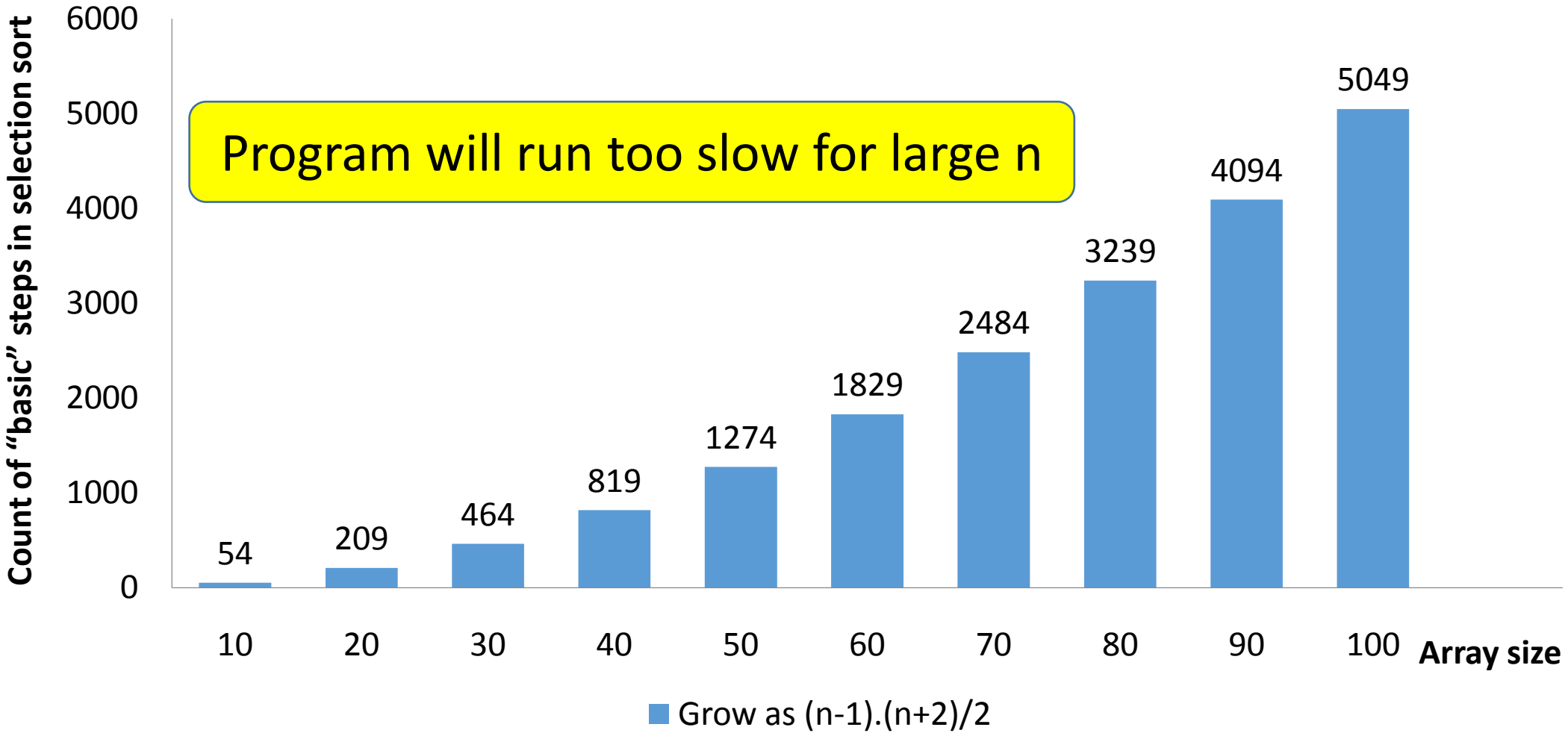

Is Selection Sort Fast Enough?

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- Real-world sorting requirements
  - Query generating 1 million data items, each with a score
  - Selection sort too slow for such applications
    - With $n = 10^6$, $(n-1).(n+2)/2 \approx 5 \times 10^{11}$
    - If each "basic" step takes 20 ns (memory reads and writes, comparison, etc.), we need 10^4 seconds (approx. 2.78 hours)!!!
- Can we do better?
  - Yes, much better !!!
  - Approximately $(n. \log_2 n)$ "basic" steps to sort an array of size n
    - **10^6 elements can be sorted in no more than a few seconds!**
  - Topic of next few lectures …

Summary

- Analysis of performance of selection sort
  - Count of “basic” steps grows quadratically with size of array
- Need for faster sorting techniques