

Let us study the dual, slater's condition, strong duality, KKT conditions & application of all this through the example of LASSO

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2 + \lambda \|x\|_1$$

where  $\lambda$  is fixed (not a Lagrange variable)

### Recap

For Lasso, it can be shown that for every  $\theta$  there exists a  $\lambda > 0$  s.t. following two problems are equivalent:

$$\min_x \|Ax - b\|_2^2 + \lambda \|x\|_1 \quad \textcircled{1} \dots \text{say sum is } \hat{x} \text{ &} \\ \|x^*\|_1 = \beta$$

$$\begin{aligned} \min_x & \|Ax - b\|_2^2 \\ \text{s.t.} & \|x\|_1 \leq \theta \end{aligned} \quad \textcircled{2} \dots \text{say solution is } \hat{x}$$

Solution to  $\textcircled{2}$  with  $\theta = \beta = \hat{x}^*$  is also  $\hat{x}^*$ !

Solution to  $\textcircled{1}$  with  $\lambda$  as soln to  $A^T(b - Ax) = \lambda g_{\hat{x}}$  is also  $\hat{x}$ !  
 $g_{\hat{x}} \in \partial \|\hat{x}\|_1$

What about dual of Lasso?

Dual of  $\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - Ax\|_2^2 + \lambda \|x\|_1$  is a constant!

Redefine primal as:

$$\begin{array}{ll} \min & \frac{1}{2} \|y - z\|_2^2 + \lambda \|z\|_1 \\ z \in \mathbb{R}^m & \end{array}$$

$$\text{s.t. } z = Ax$$

$$\begin{aligned} L^*(u) &= \min_{z \in \mathbb{R}^m} \frac{1}{2} \|y - z\|_2^2 + \lambda \|z\|_1 + u^T(z - Ax) \\ &= \frac{1}{2} \|y\|_2^2 - \frac{1}{2} \|y - u\|_2^2 - I\left(\left\|\frac{A^T u}{\lambda}\right\|_\infty \leq 1\right) \end{aligned}$$

$\therefore$  Lasso dual problem is

$$\max_{\mu \in \mathbb{R}^m} \frac{1}{2} (\|b\|^2 - \|b - \mu\|^2)$$

$$\text{s.t. } \|A^\top \mu\|_\infty \leq \lambda$$

$$\stackrel{\text{Lc}}{=} \max_{\mu \in \mathbb{R}^m} \frac{1}{2} (\|\mu - b\|_2^2)$$

$$\text{s.t. } \|A^\top \mu\|_\infty \leq \lambda$$

Recall: Dual norm for  $p$  is  $q$  st

$$\frac{1}{p} + \frac{1}{q} = 1$$

Here  $p=1$  &  $q=\infty$

$\therefore$  KKT conditions are necessary & sufficient  
(since Slater's condition satisfied)

① Primal constraints:  $\hat{z} = \hat{A}\hat{x}$

② Dual constraints: None

③ Complementary slackness: None

④ Subdiff wrt primal variables = 0:

$$\partial_z L(\hat{x}, \hat{z}, \hat{\mu}) = \hat{z} - b - \hat{\mu} = 0$$

$$\partial_{\hat{x}} L(\hat{x}, \hat{z}, \hat{\mu}) = \lambda \begin{bmatrix} \text{sgn}(\hat{x}_i) & \text{if } \hat{x}_i \neq 0 \\ 0_i \in [-1, 1] & \text{if } \hat{x}_i = 0 \end{bmatrix} + A^T \hat{\mu} = 0$$

Q: How to use the dual formulation & KKT?

Ans.: We know by Slater's condition & KKT that if  $\hat{x}, \hat{\lambda}, \hat{z}$  satisfy KKT then

$$\frac{1}{2} \|A\hat{x} - b\|^2 + \lambda \|\hat{x}\|_1 = \min_{\substack{x \in \mathbb{R}^n \\ z \in \mathbb{R}^m \\ \text{s.t. } Ax = z}} \frac{1}{2} \|z - b\|_2^2 + \lambda \|x\|_1$$

$$= \max_{\substack{M \in \mathbb{R}^m \\ \text{s.t. } \|A^T M\|_\infty \leq \lambda}} \frac{1}{2} \|M - b\|_2^2$$

$$= \frac{1}{2} \|\hat{M} - b\|_2^2$$

Iterative algo for primal  
Claim: If  $\bar{a}_j = 2 \sum_{i=1}^n [A_{ij}]^2$  &  $\bar{b}_j = 2 \sum_{i=1}^n A_{ij}(y_i - \bar{x}_j^T A_j)$

Then:  $\bar{x}_j^* = \begin{cases} (\bar{b}_j + \lambda) / \bar{a}_j & \text{if } \bar{b}_j < -\lambda \\ 0 & \text{if } \bar{b}_j \in [-\lambda, \lambda] \\ (\bar{b}_j - \lambda) / \bar{a}_j & \text{if } \bar{b}_j > \lambda \end{cases}$

So to satisfy this, Lasso iterates on  $x^K$  as follows:  $x^{(0)} \rightarrow \bar{b}^{(0)} \rightarrow x^{(1)} \dots \dots$  until convergence --- We can understand through following simplification where  $A = I$

Iterative algo for primal ...

Stopping criterion :

$x^{(k)} \rightarrow$  solve for  $\mu^{(k)}$  &  $z^{(k)}$  using  
 KKT

Check gap  $\left\{ \frac{1}{2} \|Ax^{(k)} - b\|_2^2 + \lambda \|x^{(k)}\|_1 \right\} \leq C$   
 $- \frac{1}{2} \|\mu^{(k)} - b\|_2^2$

$$\min f(\mathbf{x})$$

$\mathbf{x} \in \mathcal{D}$

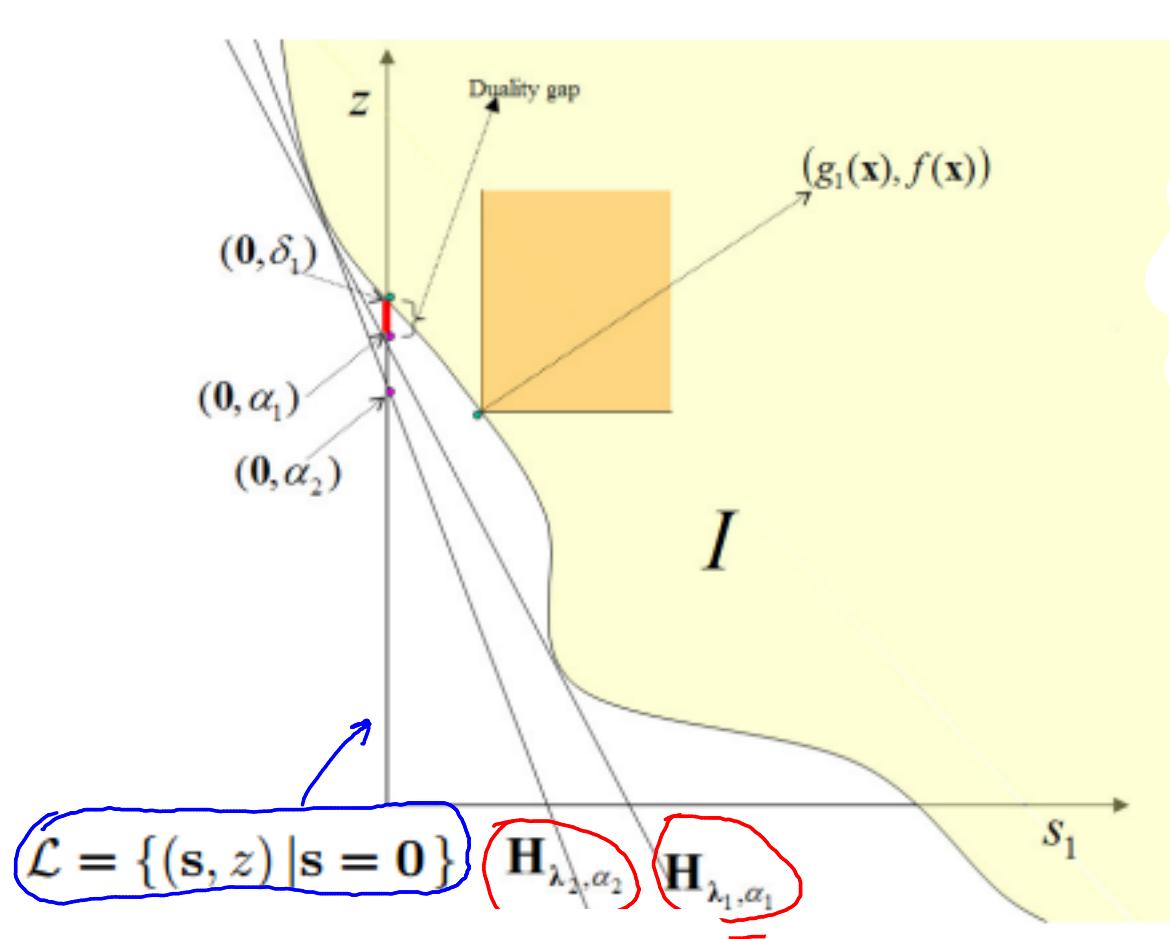
$$\text{s.t. } g_i(\mathbf{x}) \leq 0 \quad i=1 \dots m$$

(The general dual problem & its geometric interpretation)

pg 292, sec 4.4.3 of <http://www.cse.iitb.ac.in/~cs709/notes/BasicsofConvexOptimization.pdf>

Consider the set:

$$\mathcal{I} = \{(\mathbf{s}, z) \mid \mathbf{s} \in \mathbb{R}^m, z \in \mathbb{R}, \exists \mathbf{x} \in \mathcal{D} \text{ with } g_i(\mathbf{x}) \leq s_i \forall 1 \leq i \leq m, f(\mathbf{x}) \leq z\}$$



$$\mathcal{H}_{\lambda, \alpha} = \{ (\mathbf{s}, z) \mid \lambda^T \cdot \mathbf{s} + z = \alpha \}$$

$$\min f(\mathbf{x})$$

$\mathbf{x} \in \mathcal{D}$

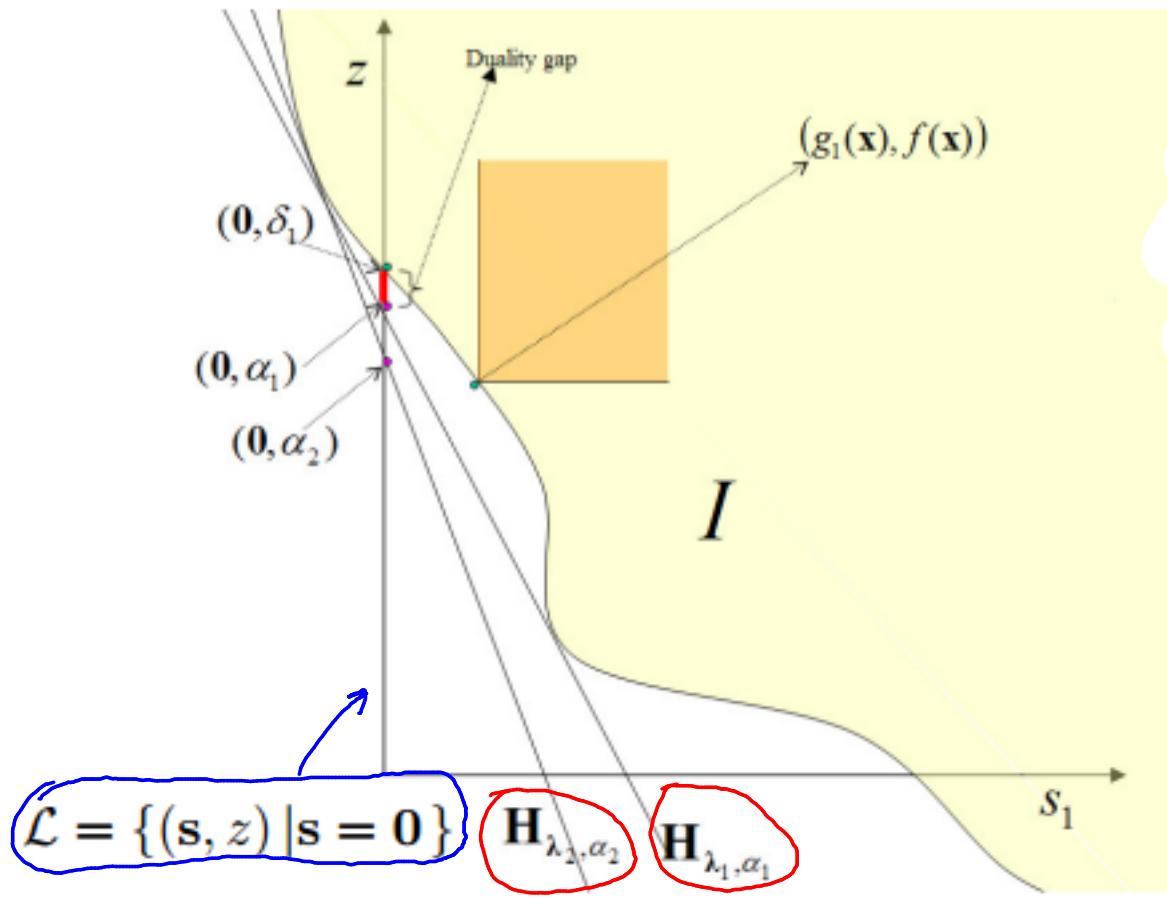
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smallest  $z$  value  
in  $\mathcal{I}$  for  
 $s_i \leq 0$  will be

for  $s_i = 0$

since  $(s_i, z_i) \in \mathcal{I}$

$\Rightarrow (s_2, z_2) \in \mathcal{I}$

$\& s_2 \geq s_1 \& z_2 \geq z_1$

thus  $\Rightarrow$    
is not possible!

$$\mathcal{H}_{\lambda, \alpha} = \{ (\mathbf{s}, z) \mid \lambda^T \cdot \mathbf{s} + z = \alpha \}$$

$$\begin{array}{ll} \max & \alpha \\ \text{subject to} & \mathcal{H}_{\lambda, \alpha}^+ \supseteq \mathcal{I} \end{array}$$

$$\begin{array}{ll} \max & \alpha \\ \text{subject to} & \lambda^T \cdot \mathbf{s} + z \geq \alpha \quad \forall (\mathbf{s}, z) \in \mathcal{I} \end{array}$$

$$\begin{array}{ll} \max & \alpha \\ \text{subject to} & \lambda^T \cdot \mathbf{s} + z \geq \alpha \quad \forall (\mathbf{s}, z) \in \mathcal{I} \\ & \lambda \geq \mathbf{0} \end{array}$$

(a)

If  $\nexists \alpha_b$  soln to (b),  
 st  $\lambda^T g(\alpha_b) + f(\alpha_b) = \alpha_b$   
 then  $s = g(\alpha_b)$   $z = f(\alpha_b)$   
 is soln to (a)  
 complete proof of equivalence

(b)

$$\begin{array}{ll} \max & \alpha \\ \text{subject to} & \lambda^T \cdot \mathbf{g}(\mathbf{x}) + f(\mathbf{x}) \geq \alpha \quad \forall \mathbf{x} \in \mathcal{D} \\ & \lambda \geq \mathbf{0} \end{array}$$

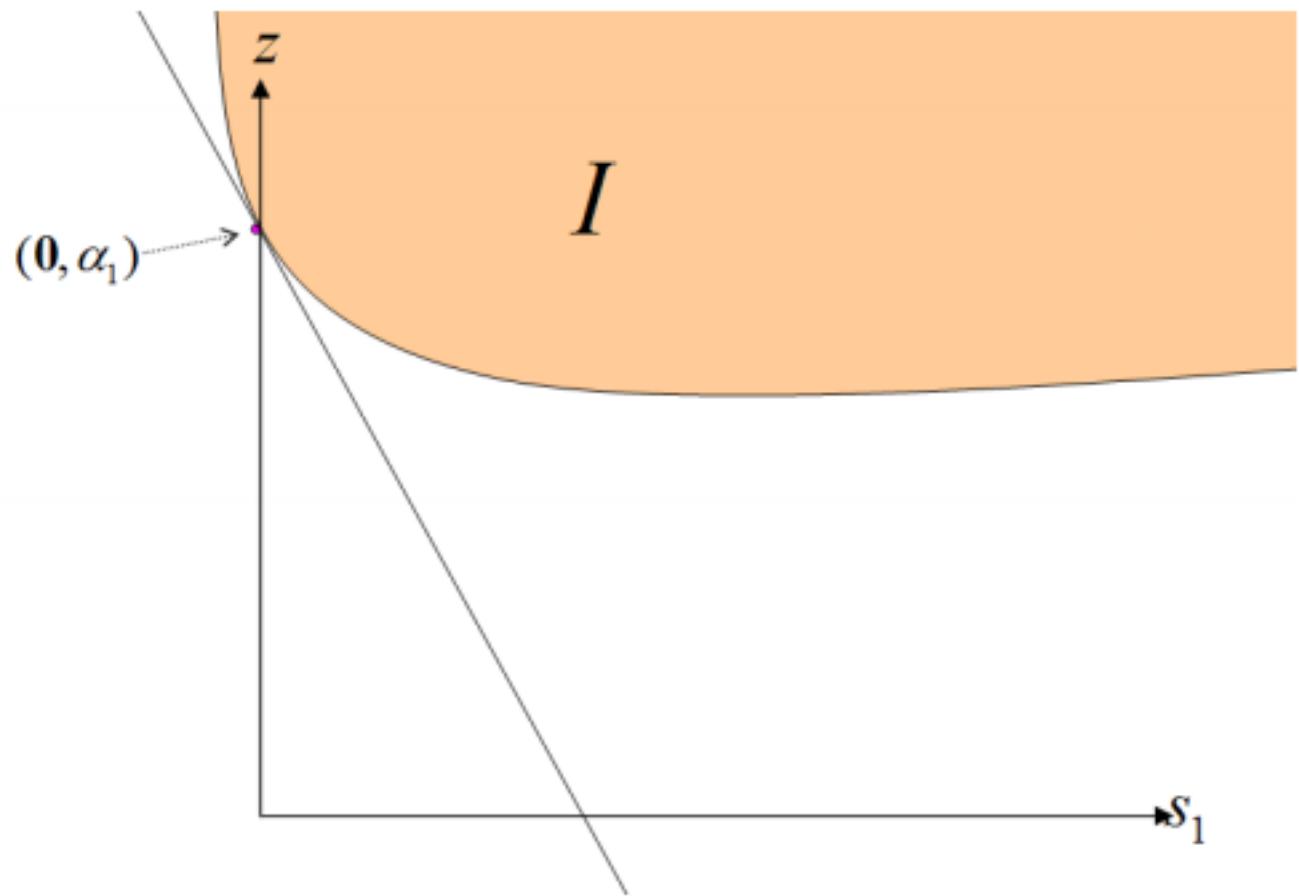
$$\begin{aligned} & \max && \alpha \\ \text{subject to} & & L(\mathbf{x}, \lambda) \geq \alpha \quad \forall \mathbf{x} \in \mathcal{D} \\ & & \lambda \geq \mathbf{0} \end{aligned}$$

Since,  $L^*(\lambda) = \min_{\mathbf{x} \in \mathcal{D}} L(\mathbf{x}, \lambda)$ , we can deal with the equivalent

$$\begin{aligned} & \max && \alpha \\ \text{subject to} & & L^*(\lambda) \geq \alpha \\ & & \lambda \geq \mathbf{0} \end{aligned}$$

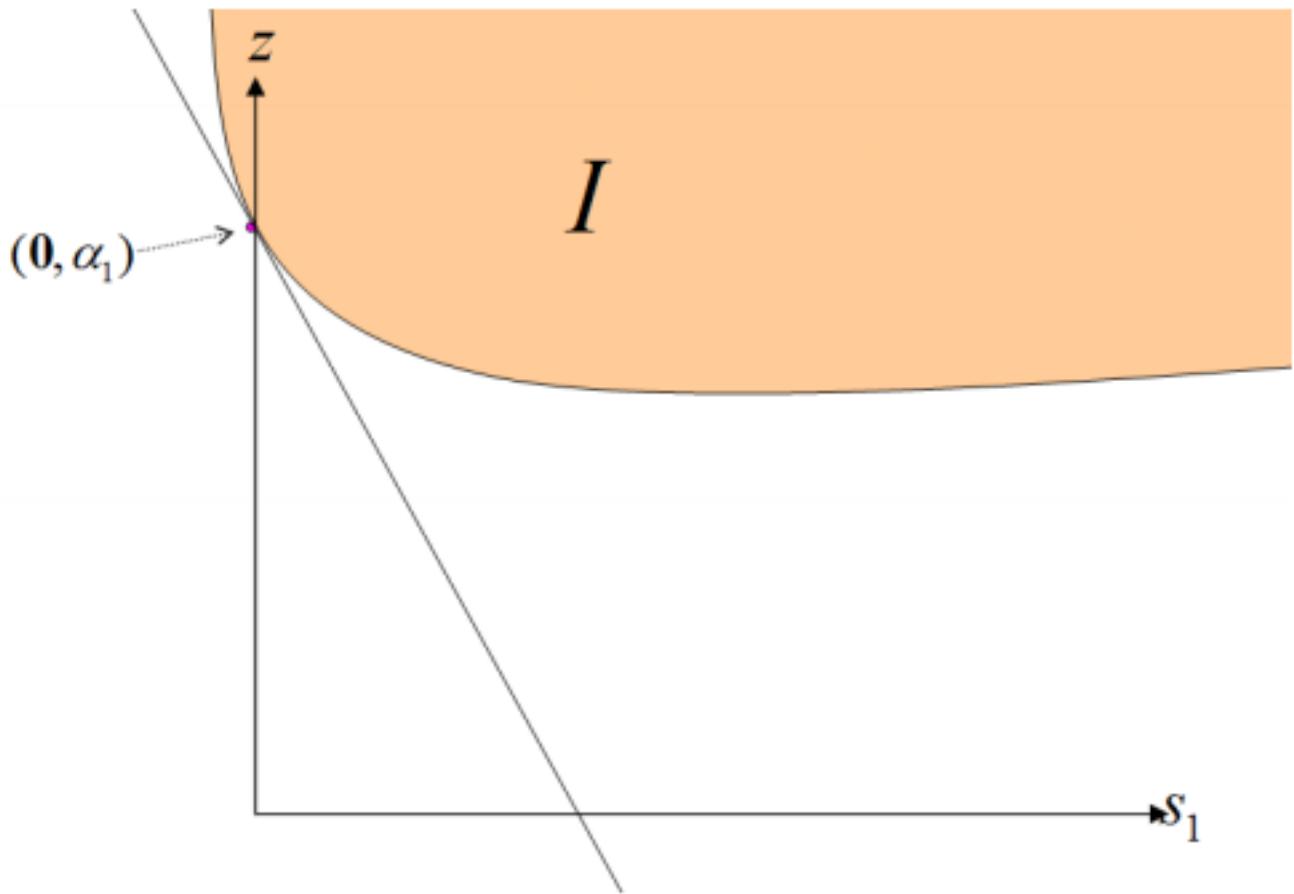
This problem can be restated as

$$\begin{aligned} & \max && L^*(\lambda) \\ \text{subject to} & & \lambda \geq \mathbf{0} \end{aligned}$$



Q: What is desirable of the set  $I$  for zero duality gap?

- ①  $I$  should be convex  $\Leftrightarrow f$  &  $g_i$ 's are convex
- ②  $(0, \delta_1)$  should exist  $\Leftrightarrow I$ 's intersection with  $s=0$  should be closed



Q: What is desirable of the set  $I$  for zero duality gap?

Ans:  $\exists (0, \alpha) \in I$  and  $\lambda$  s.t  $\lambda^T s + z \geq \alpha \forall (s, z) \in I$   
& intersection of  $I$  with  $z$  axis is closed below

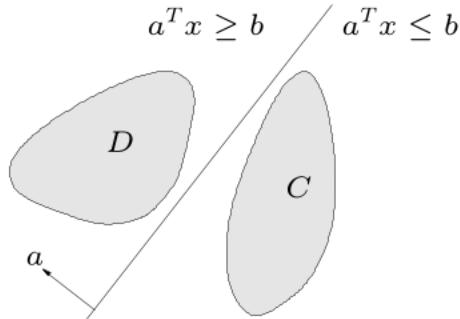
$\Leftrightarrow \exists$  a supporting hyperplane to  $I$  at  $(0, \alpha)$   
& intersection of  $I$  with  $z$  axis is closed below (with  $(0, \alpha)$  being boundary pt)

$\Leftarrow I$  is closed &  $\exists$  a supporting hyperplane to  $I$  at every boundary point

## Separating hyperplane theorem

if  $C$  and  $D$  are disjoint convex sets, then there exists  $a \neq 0, b$  such that

$$a^T x \leq b \text{ for } x \in C, \quad a^T x \geq b \text{ for } x \in D$$



the hyperplane  $\{x \mid a^T x = b\}$  separates  $C$  and  $D$

strict separation requires additional assumptions (e.g.,  $C$  is closed,  $D$  is a singleton)

Convex sets

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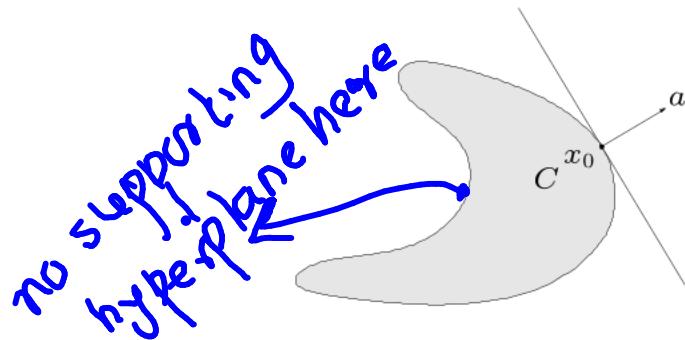
Consequence

## Supporting hyperplane theorem

**supporting hyperplane** to set  $C$  at boundary point  $x_0$ :

$$\{x \mid a^T x = a^T x_0\}$$

where  $a \neq 0$  and  $a^T x \leq a^T x_0$  for all  $x \in C$



**supporting hyperplane theorem:** if  $C$  is convex, then there exists a supporting hyperplane at every boundary point of  $C$

Proof (from separating hyperplane theorem):

a)  $\text{Interior}(C) \neq \emptyset$  (so that  $\text{int}(C) \cap \{x_0\} = \emptyset$ )

Apply separating hyperplane theorem

to sets  $C' \subseteq \{x_0\}$  and  $D' = \text{Interior}(C)$

$$\exists a, b \text{ s.t. } a^T x \geq b \quad \forall x \in D'$$

$$a^T x_0 \leq b$$

Strict separation not applicable since  $\text{int}(C)$  is open

This inequality extends to all boundary (limit pts) leading to  $a^T x_0 \geq b$

$$\Rightarrow \text{combined: } a^T x_0 = b$$

b)  $\text{int}(C) = \emptyset$

$\Rightarrow$  Any hyperplane containing that affine set contains  $C \cup x_0$  (H/w: Prove)

$\Rightarrow$  This hyperplane is a trivial supporting hyperplane

**Some topological concepts:** [Topological set is set with concept of nbrhood]

① A set  $U$  is called an open set if it does not contain any of its boundary pts. If  $S$  is a metric space (eg an inner product space) with distance metric  $d(x,y)$ , then a subset  $U$  of  $S$  is called open if, given any  $x \in U$ ,  $\exists \epsilon > 0$  such that given any  $y \in S$  with  $d(x,y) < \epsilon$ ,  $y \in U$

② A set  $V \subseteq S$  is called closed if its complement  $S \setminus V$  is an open set

③  $x \in S$  is called an **interior point** of  $S$  if there exists a neighborhood of  $x$  contained in  $S$ . If  $S$  is a metric space, then  $x \in S$  is an **interior pt** if  $\exists \epsilon > 0$  s.t.  $\forall y$  s.t  $d(x,y) < \epsilon, y \in S$

The set of all interior pts of  $S$  form the **interior of  $S$** . Thus, if  $S$  is a metric space:

$$\text{int}(S) = \underbrace{\{x \mid \exists \epsilon > 0 \text{ s.t. } \forall y \text{ s.t. } d(x,y) < \epsilon, y \in S\}}_{\text{interior of } S}$$

What can I say if  $\text{interior}(C) = \emptyset$

- e.g.: sufficient conditions**
- ①  $C \subseteq$  hyperplane. In particular, if  $S$  is affine this is necessary & sufficient
  - ② e.g. a shell

- ③ e.g.  $C = \partial K$  in topological space  $S$   
(see next page for defn of boundary of set  $K$  denoted by  $\partial K$ )

④ The set of pts of a set  $S$  s.t every neighborhood of a point from the set consists of atleast one point in  $S$  and one point not in  $S$  is called the **boundary**  $\partial S$  of  $S$ . If  $S$  is a metric space

$$\partial S = \{x \in S \mid \forall \epsilon > 0, \exists y \text{ s.t } d(x,y) < \epsilon \text{ and } \exists y' \text{ s.t } d(x,y') < \epsilon \text{ and } y' \notin S\}$$

⑤ Let  $S$  be a subset of a topological space  $X$ . A point  $x \in X$  is a limit point of  $S$  if every neighborhood of  $x$  contains at least one point of  $S$  different from  $x$  itself.

*Can be relaxed to open neighborhoods without loss*

If  $S$  happens to have an associated metric  $d$ , and  $A \subseteq S$ , then  $a \in S$  is a limit point of  $A$  iff:

$$\forall \epsilon > 0 : \{x \in A \text{ s.t } 0 < d(x,a) < \epsilon\} \neq \emptyset$$

[Note:  $a$  need not belong to  $A$ ]

⑥ Closure of  $S$   $[cl(S)] = S \cup \{\text{limit points of } S\}$

Informally speaking,  $a$  is a limit point of  $A$  if there are points in  $A$  that are different from  $a$  but arbitrarily close to it

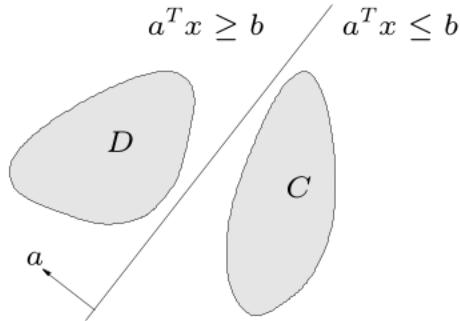
Some standard results that we will regularly invoke for topological spaces

- ① Intersection of (even uncountable)  
closed sets is closed
- ② Union of (even uncountable)  
open sets is open
- ③ Intersection of finite number of  
open sets is open
- ④ Union of finite number of  
closed sets is closed
- ⑤  $S$  is closed iff  $S^c$  is open

## Separating hyperplane theorem

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**Proof:** Let  $S = \{x - y \mid x \in C, y \in D\}$ .

Now we can prove (see <http://www.cse.iitb.ac.in/~cs709/notes/eNotes/ExtraProblems-1.pdf>, Q1) that  $S$ , being a sum of two convex sets, is convex.

Since  $C \cap D = \emptyset$ ,  $0 \notin S$ .

(a) Suppose  $0 \notin \text{cl}(S)$ : Consider the sets  $\{0\}$  and  $\text{cl}(S)$ . We will prove that  $\exists a \neq 0$  s.t  $a^T z > 0 \forall z \in \text{cl}(S)$  &  $a^T w = 0$  for  $w \in \{0\}$

$\underbrace{a^T z > 0 \forall z \in \text{cl}(S)}$

Q: How to choose  $a$ ?

$\underbrace{a^T w = 0 \text{ for } w \in \{0\}}$

obvious

complete proof given in class:  $k/w$

i.e  $\exists \alpha$  s.t  $\alpha^T(x-y) > 0 \quad \forall x-y \in S$

i.e  $\alpha^T x > \alpha^T y \quad \forall x \in C \text{ & } y \in D$

Let  $b = \inf_{x \in C} \alpha^T x$ . Then we proved existence

of  $a \in \mathbb{R}$  s.t

$\alpha^T x \geq b \quad \forall x \in C \text{ & } \alpha^T y \leq b \quad \forall y \in D$

(b) suppose  $0 \in \text{cl}(S)$ . Since  $0 \notin S$ ,  $0 \in \text{bdry}(S)$   
If  $\text{interior}(S) = \emptyset$ ,  $S$  must be  $\subseteq \{z | \alpha^T z = b\}$   
& the hyperplane must include  $0$  on  $\text{bdry}(S)$  +  
 $\Rightarrow b = 0$ . i.e  $\alpha^T x = \alpha^T y \quad \forall x \in C \text{ & } y \in D$  A hyperplane  
 $\Rightarrow$  we have a trivial separating hyperplane

limit by  $\bar{a}$ , we have

$$\bar{a}(\epsilon_k)^T z > 0 \quad \forall z \in S_{-\epsilon_k}$$

for all  $k$  & therefore

$$\bar{a}^T z > 0 \quad \forall z \in \text{interior}(S)$$

and

$$\bar{a}^T z \geq 0 \quad \forall z \in S \quad \begin{matrix} \text{proof by} \\ \text{contradiction} \end{matrix}$$

that is

$$\bar{a}^T x > \bar{a}^T y$$

$$\forall x \in C \text{ } \forall y \in D$$

(use the property  
that a convex  
set is connected?)

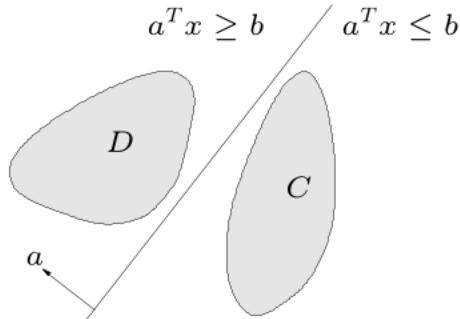
Hence proved!

## Separating hyperplane theorem

*Thus*

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*Consequence*

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**supporting hyperplane theorem:** if  $C$  is convex, then there exists a supporting hyperplane at every boundary point of  $C$