

Lattice of atoms

① Subsumption is a quasi-order

② $\therefore$ Equivalence classes of atoms are related by partial order.

③ $A_1 \equiv A_2 \dots$ iff $\rightarrow$ same pred symbol
 $\rightarrow$ One-to-one correspondence between var names in A_1 & A_2 .
 $\underbrace{\{A_1\}}_{\text{eqn class}}$
 $\text{parent}(x, y) = \text{parent}(A, B)$

$\text{parent}(x, y) \neq \text{parent}(\text{sis}(\text{A}), B)$
 $\underbrace{A_1} \dots \underbrace{\{x/\text{sis}(A), y/B\}}_{A_2}$

$A_1 \not\equiv A_2 \leftarrow \text{X}$

Lattice structure over atoms contd.

$\mathcal{A}^+ = \mathcal{A} \cup \{\top, \perp\}$ is a quasi order
s.t

- $\top \succeq \underline{l}$ for all $\underline{l} \in \mathcal{A}^+$
- $\underline{l} \succeq \perp$ for all $\underline{l} \in \mathcal{A}^+$
conventional
- $\underline{l} \succeq m$ iff there is a substitution θ such that $\underline{l}\theta = m$, for $\underline{l}, m \in \mathcal{A}$

$\mathcal{A}_E^+$... partially ordered.

eg: $\underline{l} = \text{Mem}(x, [x, y])$ $m = \text{Mem}(1, [1, 2])$
Then -- $\underline{l} \succeq m$ with $\theta = \{x/1, y/2\}$ (3)?

$\underline{l} \equiv ?$ $\text{Mem}[x, [x, y_1]]$
 $\text{Mem}[x_2, [x_2, y_2]]$

Lattice structure over atoms contd.

Q: If l & m are "atoms" then
Is $l \preceq m$ iff $l \models m$?

for clauses $C_1 \in \mathcal{C}_2$

$C_1 \preceq C_2$ only if $C_1 \models C_2$

if did not hold because of
self-recursive clauses

$\Rightarrow$ Logical implication over atoms is
also a quasi-order over atoms

Lattice structure over atoms contd.

Embodies relations

- $[\perp] \sqcap [I] = [\perp]$, and $[\top] \sqcap [I] = [I]$

- If $I_1, I_2 \in \mathcal{A}$ have a most general unifier (see page 78) θ then $[I_1] \sqcap [I_2] = [I_1\theta] = [I_2\theta]$.

This can be proved as follows. Let $[u] \in \mathcal{A}_E^+$ such that $[I_1] \succeq [u]$ and $[I_2] \succeq [u]$, then we need to show that $[I_1\theta] \succeq [u]$. If $[u] = [\perp]$, this is obvious. If $[u]$ is conventional, then there are substitutions σ_1 and σ_2 such that $[I_1\sigma_1] = [u] = [I_2\sigma_2]$. Here we can assume σ_1 only acts on variables in I_1 , and σ_2 only acts on variables in I_2 . Let $\sigma = \sigma_1 \cup \sigma_2$. Notice that σ is a unifier for $\{[I_1], [I_2]\}$. Since θ is an mgu for $\{[I_1\sigma_1], [I_2\sigma_2]\}$, there is a γ such that $\theta\gamma = \sigma$. Now $[I_1\theta\gamma] = [I_1\sigma] = [I_1\sigma_1] = [u]$, so $[I_1\theta] \succeq [u]$.

- If $I_1, I_2 \in \mathcal{A}$ do not have a most general unifier θ then $[I_1] \sqcap [I_2] = [\perp]$.

Since I_1 and I_2 are not unifiable, there is no conventional atom u such that $[I_1] \succeq [u]$ and $[I_2] \succeq [u]$. Hence $[I_1] \sqcap [I_2] = [\perp]$.

- $[\perp] \sqcup [I] = [I]$, and $[\top] \sqcup [I] = [\top]$

- If I_1 and I_2 have an "anti-unifier" m then $[I_1] \sqcup [I_2] = [m]$; otherwise $[I_1] \sqcup [I_2] = [\top]$: Proof on [9]

Assume I_1 & I_2 standardized apart



Wont hold simultaneously

Anti-unification $\equiv$ Reverse of unification

Idea: Move from constants to variables.

[Term-place notation]

Mem[1, [1, 2]]

appears in 2 places

$\{ (1, \langle 1 \rangle), (1, \langle 2, 1 \rangle), (2, \langle 2, 2 \rangle) \}$

$t = \text{term}$

$P(a, f(g(b), p(g(c))))$

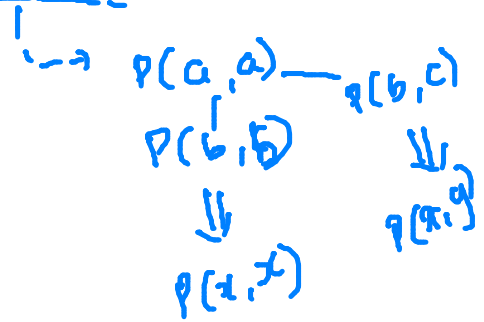
$\langle t, p \rangle = ?$

1. Let $\mathbf{l} = \mathbf{l}_1$ and $\mathbf{m} = \mathbf{l}_2$, $\theta = \emptyset$, $\sigma = \emptyset$
2. If $\mathbf{l} = \mathbf{m}$ return $\mathbf{l}$ and stop.
3. Try to find terms t_1 and t_2 that have the same (leftmost) place in $\mathbf{l}$ and $\mathbf{m}$ respectively, such that $t_1 \neq t_2$ and either t_1 and t_2 begin with different function symbols, or at least one of them is a variable.
4. If there is no such t_1, t_2 , return $\mathbf{l}$ and stop.
5. Choose a variable x that does not occur in either $\mathbf{l}$ or $\mathbf{m}$ and wherever t_1 and t_2 occur in the same place in $\mathbf{l}$ and $\mathbf{m}$, replace each of them by x
6. Set θ to $\theta \cup \{x/t_1\}$ and σ to $\sigma \cup \{x/t_2\}$
7. Go to Step 3

$\left. \begin{array}{l} \text{Parent}(\text{ann}, \text{mary}) \\ \text{Parent}(\text{ann}, \text{tom}) \end{array} \right\}$

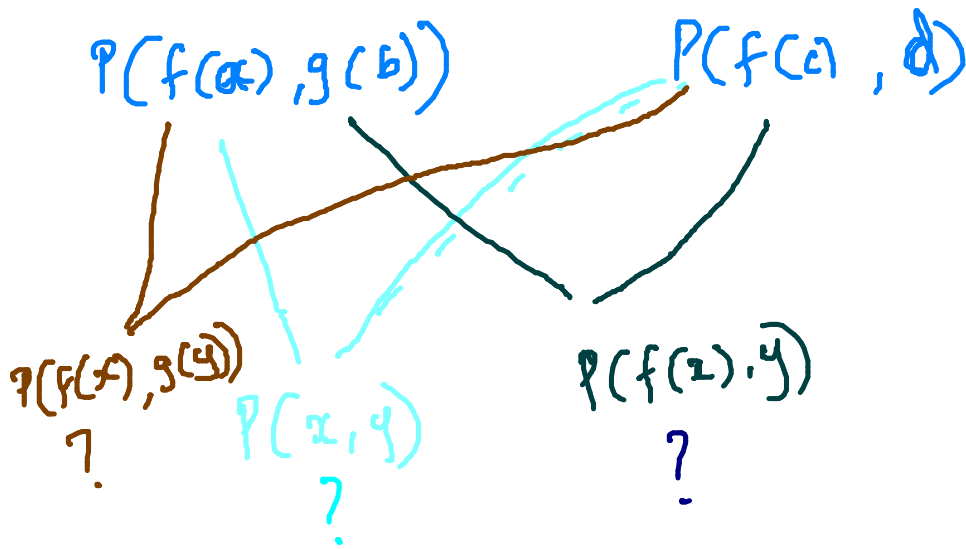


Parent(ann, x)



ANTI-UNIFICATION ALGO

Beware: You need to take minimal step



$$Q: \text{Mem}(1, [1, 2]) \sqcup \text{Mem}(2, [2, 4]) = ?$$

Theorem 21 After each iteration of the Anti-Unification Algorithm, there are terms $s_1, \dots, s_i$ and $t_1, \dots, t_i$ such that:

1. $\theta = \{z_1/s_1, \dots, z_i/s_i\}$ and $\sigma = \{z_1/t_1, \dots, z_i/t_i\}$.
2. $\mathbf{l}\theta = \mathbf{l}_1$ and $\mathbf{m}\sigma = \mathbf{l}_2$.
3. For every $1 \leq j \leq i$, s_j and t_j differ in their first symbol.
4. There are no $1 \leq j, k \leq i$ such that $j \neq k$, $s_j = s_k$ and $t_j = t_k$.

Terminally
by
construction

By virtue of
step 5.

Theorem 22 Let l_1 and l_2 be two atoms with the same predicate symbol. Then the Anti-Unification Algorithm with l_1 and l_2 as inputs returns $l_1 \sqcup l_2$.

① Algorithm will terminate after a finite # of steps (since finite terms)

② Let u be finally returned atom.

θ, σ be final substitutions

By Thm 21, $u\theta = l_1$ & $u\sigma = l_2$ } $u \geq l_1$
 $u \geq l_2$

By composition

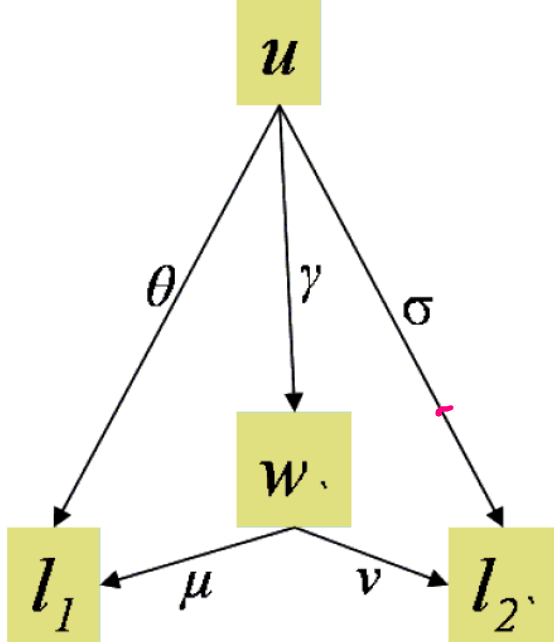
③ To show:- $u = l_1 \sqcup l_2$

Lets say $v \geq l_1$ $v \geq l_2$: Show $v \geq u$

Let $w = u \sqcap v$ (exists by prev lemma)

$u \geq w$ & $v \geq w \Rightarrow u\theta = w\theta$ & $v\sigma = w\sigma$

$\Downarrow$
 $w \geq l_1$ $w \geq l_2$ (by props of lattice)
(else l_1 would be g/b)



$$u\gamma = w$$

$$l_1: w\mu = \underline{u}\gamma\mu = \underline{u}\theta$$

$$l_2: w\nu = \underline{u}\gamma\nu = \underline{u}\sigma$$

$\Rightarrow$ If x is var in u
 then $x\theta = x\gamma\mu$
 $x\sigma = x\gamma\nu$

Claim: $w \in [u]$ ie γ is simply a renaming subst [1-1 map]

By contradiction:- say γ maps var "x" / γ unified
 to non-var "t" / "x" & "y"
 ① ②

$$\textcircled{1} \quad x\tau = t \leftarrow \text{---}$$

→ If x is not any z_j , then

$\bar{0}$ does: $x\tau u = x\theta = x$ --- contradiction to
'not act'
'on x '

$$\rightarrow x = z_j \Rightarrow x\theta = s_j = x\tau u = tu$$

$$x\sigma = t_j = x\tau u_j = tu$$

↓↓

s_j & t_j start with first symbol of
 $t \Rightarrow$ contradiction to thm 21 (5)

Suppose x and y are distinct variables in $\mathbf{u}$ such that γ unifies x and y .
Then,

1. If neither x nor y is one of the z_j 's, then $x\gamma\mu = x\theta = x \neq y = y\theta = y\gamma\mu$,
contradicting $x\gamma = y\gamma$

2

2. If x equals some z_j and y does not, then $x\gamma\mu = x\theta = s_j$ and $x\gamma\nu = x\sigma = t_j$, so $x\gamma\mu \neq x\gamma\nu$ by theorem 21, part 3. But $y\gamma\mu = y\theta = y = y\sigma = y\gamma\nu$, contradicting $x\gamma = y\gamma$.

3. Similarly for the case where y equals some z_j and x does not.

4. If $x = z_j$ and $y = z_k$, then $j \neq k$, since $x \neq y$. Furthermore, $s_j = x\theta = x\gamma\mu = y\gamma\mu = y\theta = s_k$ and $t_j = x\sigma = x\gamma\nu = y\gamma\nu = y\sigma = t_k$. But this contradicts theorem 21, part 4 $\rightarrow s_j \neq t_j$ cannot start with same symbol

-- $\rightarrow i \neq j$ then $s_i \neq s_j$ & $t_i \neq t_j$ not possible