

# Provable Non-convex Optimization for ML

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# Overview

$$\begin{aligned} \min_X f(X) \\ \text{s. t. } \text{rank}(X) \leq r \end{aligned}$$

- Projected gradient descent
- Alternating Minimization

# Our Results

- RIP/RSC based Linear Regression

$$\min_X \|A(X) - b\|_2^2 \quad s.t. \quad \text{rank}(X) \leq r$$

- $A(\cdot)$ : RIP operator
- $A(\cdot)$ : RSC operator (statistical setting)

- Matrix Completion

$$\min_X \|P_\Omega(X - M)\|_F^2 \quad s.t. \quad \text{rank}(X) \leq r$$

- $\Omega$ : randomly sampled,  $M$ : incoherent matrix

- Non-convex Robust PCA

$$\min_X \|M - X\|_0^2 \quad s.t. \quad \text{rank}(X) \leq r$$

- $M = L + S$ ,  $L$ : low-rank incoherent matrix,  $S$ : sparse matrix

# Foreground/Background Separation



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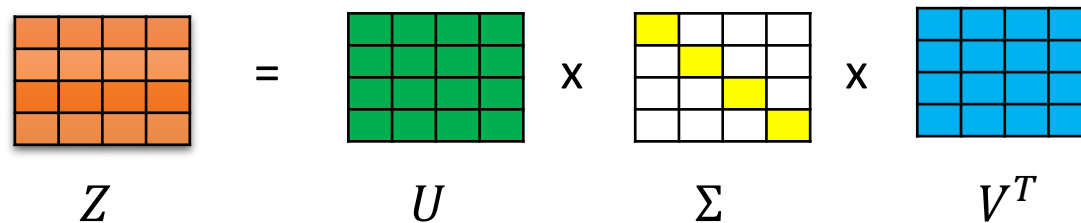
# Non-convexity of Low-rank manifold

$$0.5 \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} + 0.5 \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0.5 & 0 & 0 \\ 0 & 0.5 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

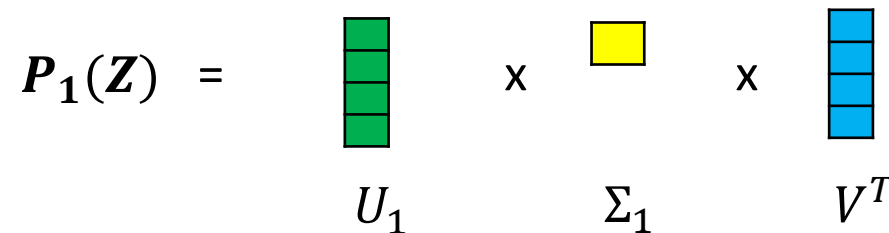
# Projection onto set of Low-rank Matrices

- Non-convex projections: NP-hard in general
- But  $P_r(Z)$  can be computed efficiently:

$$Z = U\Sigma V^T$$



- $P_r(Z) = U_r \Sigma_r V_r^T$









# Convex-projections vs Non-convex Projections

- For non-convex sets, we only have:

$$\forall Y \in C, \quad \|P_r(Z) - Z\| \leq \|Y - Z\|$$

- 0-th order condition

- But, for projection onto convex set  $C$ :

$$\forall Y \in C, \quad \|Z - P_C(Z)\|^2 \leq \langle Y - Z, P_C(Z) - Z \rangle$$

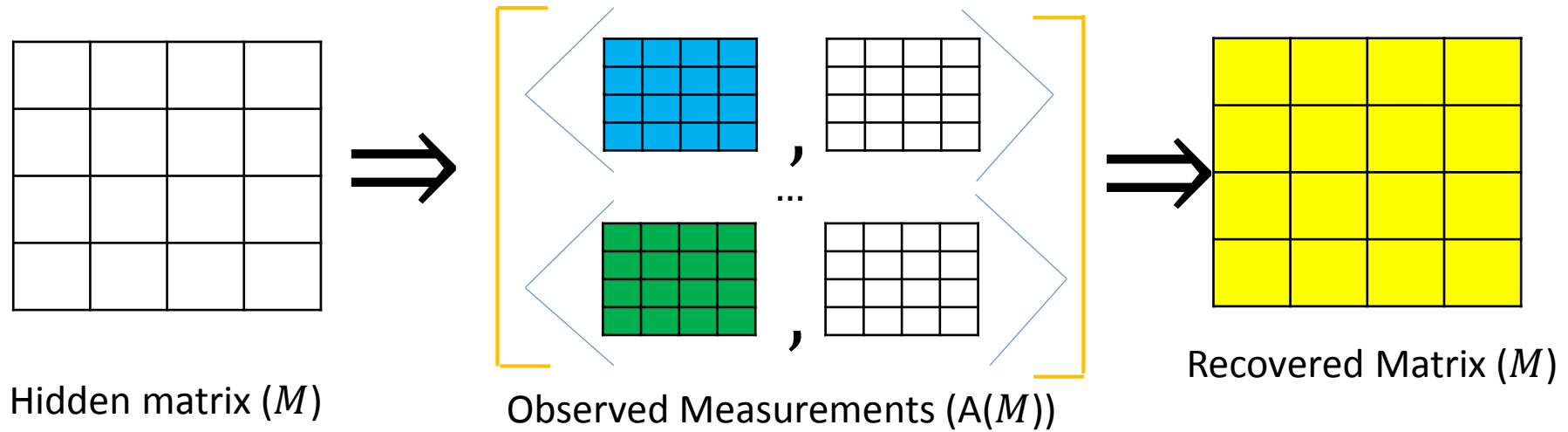
- 1-st order condition

- 0 order condition sufficient for convergence of Proj. Grad. Descent?

- In general, **NO** 😞

- But, for certain *specially structured* problems, **YES!!!**

# Low-rank Matrix Regression



# Matrix Linear Regression

$$\mathbb{A}(M) = b$$

- $\mathbb{A}: \mathbf{R}^{n \times n} \rightarrow \mathbf{R}^d$

- Linear operator

- $\mathbb{A} = \{\mathbf{A}_1, \mathbf{A}_2, \dots, \mathbf{A}_d\}$

$$\mathbb{A}(X) = \begin{bmatrix} \langle A_1, X \rangle \\ \langle A_2, X \rangle \\ \vdots \\ \langle A_d, X \rangle \end{bmatrix}$$

- Optimization Version:

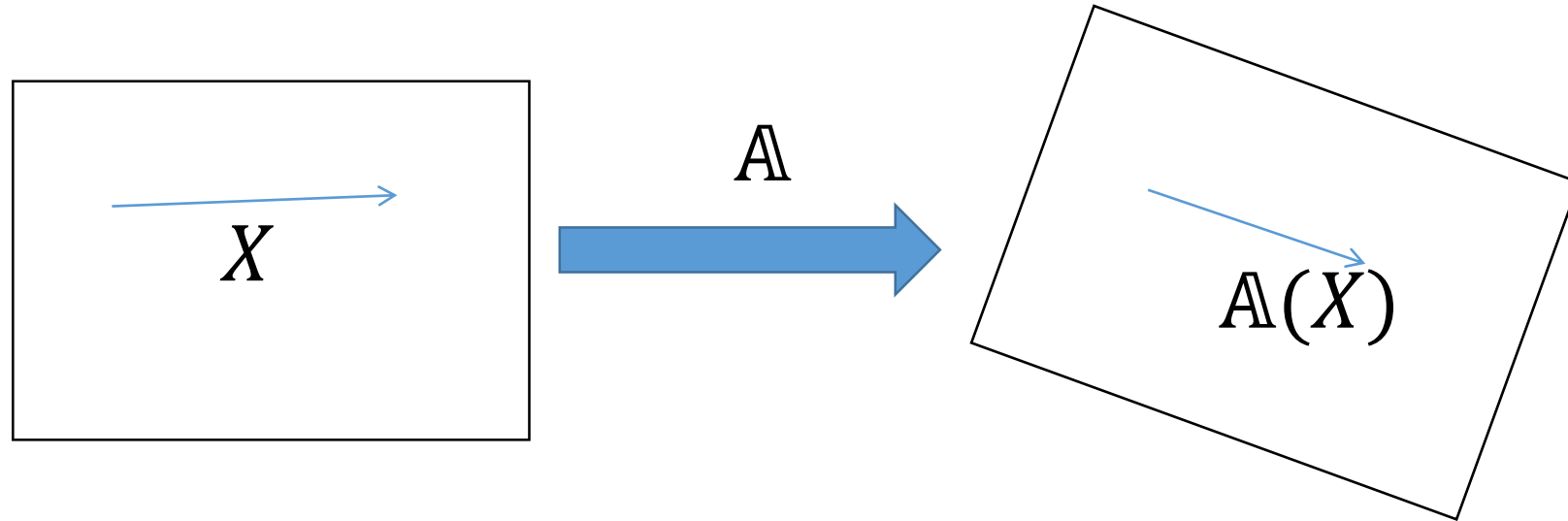
$$\begin{aligned} \min_X & \|\mathbb{A}(X) - b\|_2^2 \\ \text{s.t. } & \text{rank}(X) \leq r \end{aligned}$$

# Low-rank Matrix Estimation

$$\begin{aligned} \min_X & \quad \|A(X) - b\|_2^2 \\ \text{s. t.} & \quad \text{rank}(X) \leq r \end{aligned}$$

- NP-hard in general
  - Hard to even approximate within  $\log(n + d)$  [Meka, J., Caramanis, Dhillon'08]
- Tractable solutions under certain conditions
  - RIP conditions

# Restricted Isometry Property



- For all rank- $r$  matrix ( $X$ ):  
$$(1 - \delta_r) \|X\|_F^2 \leq \|A(X)\|_2^2 \leq (1 + \delta_r) \|X\|_F^2$$
- Examples:
  - $A$  : sampled from multivariate normal distribution
  - $m = O\left(\frac{r}{\delta_r^2} n \log n\right)$

# Approach 1: Trace-norm minimization

$$\begin{aligned} \min_X & \|\mathbb{A}(X) - b\|_2^2 \\ \text{s. t.} & \|X\|_* \leq \tau_r \end{aligned}$$

- $\|X\|_*$ : sum of singular values
- Provable recovery of  $M$ 
  - RIP based Matrix Sensing: [Recht, Fazel, Parrilo'07]
  - For Gaussian distributed samples:  $O(r n \log n)$
- However, convex optimization methods for this problem don't scale well
  - SVD computation per step
  - Intermediate iterates can have rank much larger than " $r$ "

# Approach 2: Alternating Minimization

$$\| \| \mathbf{b} - A \left( \begin{array}{c} \text{orange matrix} \\ \times \\ \text{blue matrix} \end{array} \right) \| \|_F^2$$

$$M \cong U \times V^T$$

$$V^{t+1} = \min_V \| \| b - A(U^t V^T) \| \|_2^2$$

$$U^{t+1} = \min_U \| \| b - A(U (V^{t+1})^T) \| \|_2^2$$

- Provable convergence to  $M$  [J., Netrapalli, Sanghavi'13]
  - RIP property satisfied
  - Gaussian distribution:  $O(nr^3 \log n)$ 
    - Suboptimal bounds

# Approach 3: Projected Gradient based Methods

- $X_0 = 0$
- For  $t=1:T$

$$X_t = P_r \left( X_{t-1} - \eta \mathbb{A}^T (\mathbb{A}(X_{t-1}) - b) \right)$$

- $P_r(Z)$ : projection onto set of rank-r projection
- Singular Value Projection
- Several other variants exist (ADMiRA [Lee, Bresler'09])



# Guarantees

- SVP converges to global optima
  - $\delta_{2r} \leq 1/3$
  - For Gaussians:  $O(r n \log n)$
  - Info. theoretically optimal
- Noisy case analysis also available
- Analysis: a simple extension of analysis of iterative hard thresholding [Garg, Khandekar'08]

# Extensions

- Optimize general  $f$

$$\begin{aligned} & \min_X f(X) \\ & \text{s.t. } \text{rank}(X) \leq r \end{aligned}$$

- Assume RSC-style condition:  $\forall X, \text{s.t. } \text{rank}(X) \leq r$   
 $(1 + \delta_r)I \succcurlyeq \nabla^2 f(X) \succcurlyeq (1 - \delta_r)I$

- SVP converges to the optima for such a case as well [J., Kar, Tewari'14]
- Extensions to the “statistical setting” as well

# Summary

$$\begin{aligned} \min_X f(X) \\ \text{s. t. } \text{rank}(X) \leq r \end{aligned}$$



- Projected gradient descent converges to the global optima
  - Assuming certain RSC/RIP style conditions
- Standard matrix sensing:
  - Information theoretic optimal bounds
- Analysis:
  - Only requires 0-th order property

$$\|Y - Z\| \geq \|P_r(Z) - Z\|, \quad \forall Y \in \mathcal{C}$$

# Low-rank Matrix Completion

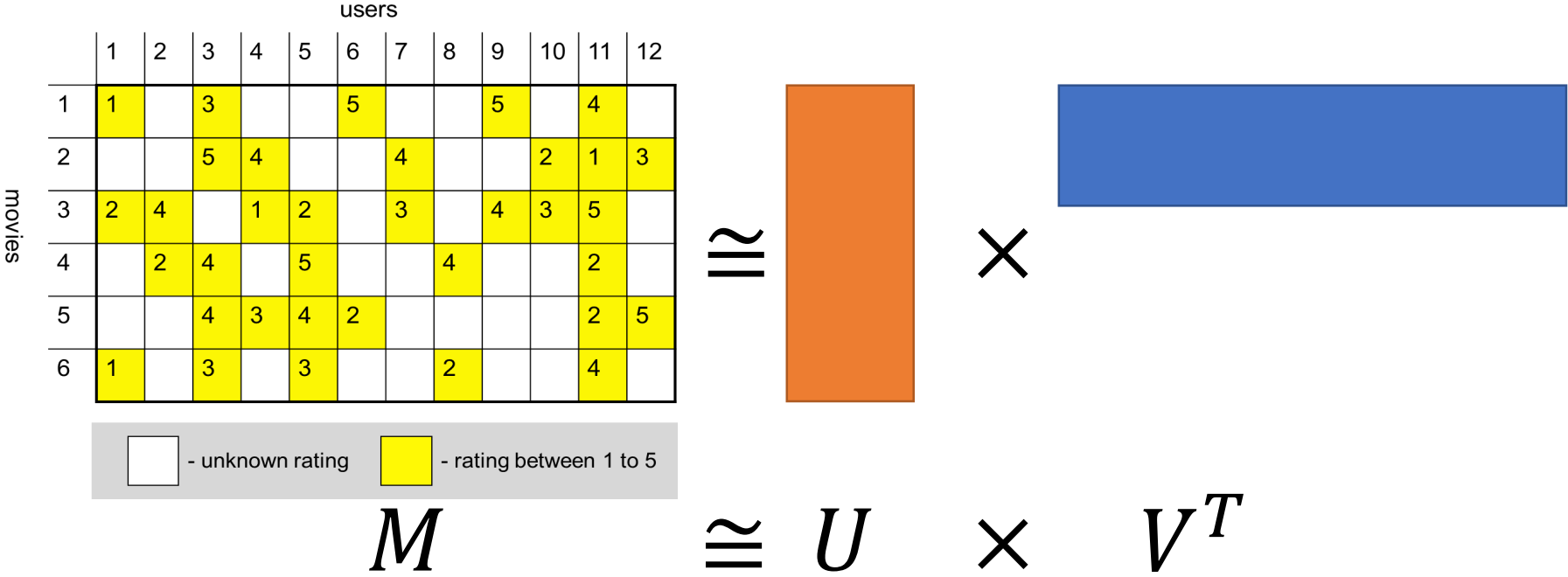
# Low-rank Matrix Completion

		users											
		1	2	3	4	5	6	7	8	9	10	11	12
movies	1	1		3			5			5		4	
	2			5	4			4			2	1	3
	3	2	4		1	2		3		4	3	5	
	4		2	4		5			4			2	
	5			4	3	4	2					2	5
	6	1		3		3			2			4	

 - unknown rating     - rating between 1 to 5

- **Task:** Complete ratings matrix
- Applications: recommendation systems, PCA with missing entries

# Low-rank



- M: characterized by U, V
- DoF:  $nr$
- No. of variables:
  - U:  $n \times r = nr$
  - V:  $n \times r = nr$

# Low-rank Matrix Completion

$$\min_X \text{Error}_\Omega(X) = \sum_{(i,j) \in \Omega} (X_{ij} - M_{ij})^2 = \|P_\Omega(X - M)\|_F^2$$

*s. t.*    **rank**( $X$ )  $\leq r$

- $\Omega$ : set of known entries
- $P_\Omega(X)_{ij} = X_{ij}, (i, j) \in \Omega$ 
  - 0 otherwise

1			
		2	
		1	
	4		

$M$



1	0	0	0
0	0	2	0
0	0	1	0
0	4	0	0

$P_\Omega(M)$

# Approach 1

- Convex relaxation: Replace  $\text{rank}(X)$  with  $\|X\|_*$
- Provably recovers  $M$  if:
  - $M$ : rank- $r$  incoherent matrix (non-spiky matrix)
    - $M = U\Sigma V^T$ ,  $\|U^i\|_2 \leq \frac{\mu\sqrt{r}}{\sqrt{n}}$
  - $\Omega$ : sampled uniformly at random and  $|\Omega| \geq O(r n \log^2 n)$
- Worst Computation time:  $O(n^3)$
- Refs: [Candes, Recht 2008], [Candes, Tao 2008], [Recht 2010]



Incoherence?

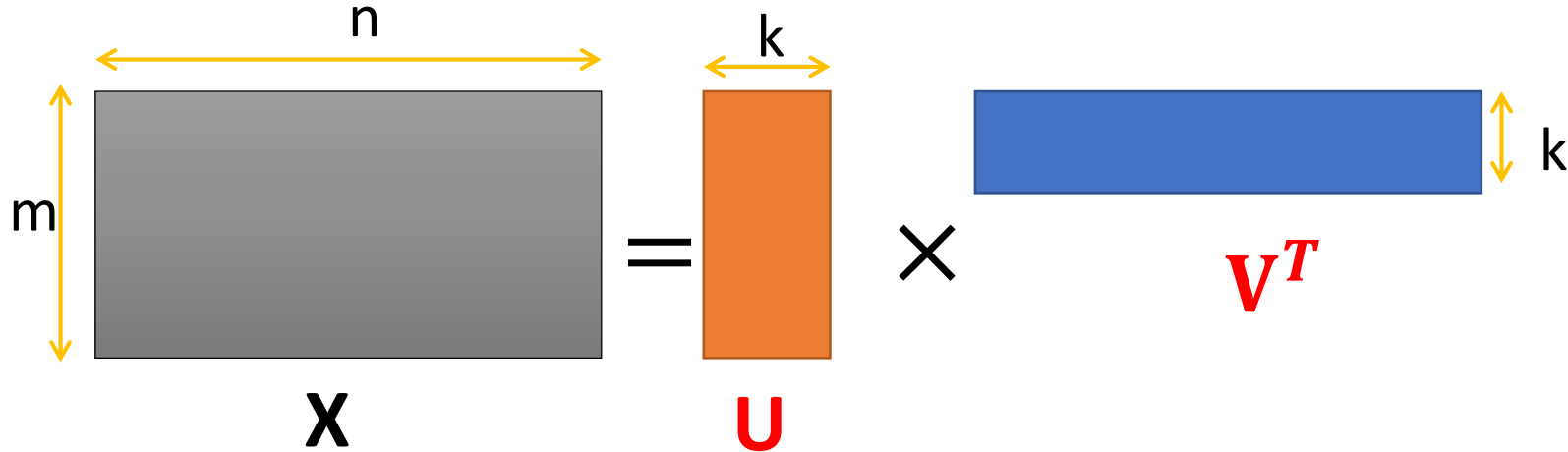


# Alternating Minimization

$$\min_X \text{Error}_\Omega(X) = \sum_{(i,j) \in \Omega} (X_{ij} - M_{ij})^2$$

s.t.  $\text{rank}(X) \leq r$

- If  $X$  has rank- $k$ :



$$V^{t+1} = \min_V \text{Error}_\Omega(U^t, V)$$

$$U^{t+1} = \min_U \text{Error}_\Omega(U, V^{t+1})$$

# Initialization [JNS'13]

- Initialization:
  - $SVD(P_{\Omega}(M), r)$

0	3	0
2	5	0
0	0	2

$P_{\Omega}(M)$

# Results [JNS'13]

- Assumptions:  $\Omega$ : set of known entries
  - $\Omega$  is sampled uniformly s.t.  $|\Omega| = O(k^7 n \log n \beta^6)$ 
    - $\beta = \sigma_1/\sigma_r$
  - $M$ : rank- $k$  “incoherent” matrix
    - Most of the entries are similar in magnitude
- Then,  $\|M - UV^T\|_F \leq \epsilon$  after only  $O(\log(\frac{1}{\epsilon}))$  steps
- Improved analysis by Hardt-Wooters'14

# Proof Sketch

- Assume Rank-1 case, i.e.,  $M = u^* v^{*T}$
- Fixing  $u$ , update for  $v$  is given by:

$$v = \arg \min_v \sum_{(i,j) \in \Omega} (u_i v_j - u_i^* v_j^*)^2$$

$$v_j = \frac{\sum_{(i,j) \in \Omega} u_i u_i^*}{\sum_{(i,j) \in \Omega} u_i^2} \cdot v_j^*$$

- If  $\Omega = [m] \times [n]$ ,

$$v_j = \langle u, u^* \rangle v_j^*$$

- Power method update!

# Proof Sketch

$$v = \underbrace{M^T u}_{\text{Power}} - \underbrace{B^{-1}(B \langle u, u^* \rangle - C)v^*}_{\text{Error Term}}$$

Method Term

Problems:

1. Show error term decreases with iterations
2. Also, need to show “incoherence” of each  $v$

Tools:

1. Spectral gap of random graphs
2. Bernstein-type concentration bounds

Bernstein?





Power Method?



# Approach 3: Singular Value Projection

$$\begin{array}{c} \text{Sample } \Omega \\ X_t = P_r(X_t - P_\Omega(X_t - M)) \end{array}$$

- Previous analysis applies only if  $P_\Omega(\cdot)$  satisfies RIP
  - RIP holds but *only* for incoherent matrices
  - $X_t - M$ : need not be incoherent

1	1	1
1	1	1
1	1	1

-

1	1	1
1	1	1
.5	.5	.5

=

0	0	0
0	0	0
.5	.5	.5

- Require:  $X_t \rightarrow M$  in  $L_\infty$  norm

# Guarantees

- Our approach:
  - Analyze  $\|X_t - M\|_\infty$  instead!
  - At first seems tricky:  $P_r(\cdot)$  optimal only w.r.t. spectral norm or Frobenius norm
- Three key tricks:
  - Use a Taylor series expansion technique by [Erdos et al' 2013]
  - Convert  $L_\infty$ -norm error bounds into  $\|\cdot\|_2$  error bounds
  - Analyze  $\|H^a u\|_\infty$

# Setting up the proof (Rank-one Case)

$$\begin{aligned} X_t &= P_1(X_{t-1} - P_\Omega(X_{t-1} - M)) \\ &= P_1(M + X_{t-1} - M - P_\Omega(X_{t-1} - M)) \\ &= P_1(M + E_t - P_\Omega(E_t)) \\ &= P_1(M + H_t) \end{aligned}$$

- $H_t = E_t - P_\Omega(E_t)$
- $E[H_t] = 0$  : assuming  $\Omega$  is independent of  $E_t$
- $E[H_t(i, j)^2] \leq \frac{\|M - X_{t-1}\|_\infty^2}{p}$
- $\|H_t\|_2 \leq \delta n \|M - X_{t-1}\|_\infty$  (assuming  $p \geq \log n / \delta^2$ )
- $\|M - X_t\|_2 \leq 2\|H_t\|_2$  (but only spectral norm bound)

Matrix Bernstein?





Matrix Perturbation?

Davis-Kahan?

# Key Step 1

- Let  $v, \lambda$  be the largest eigenvector/value of  $M + H_t$

$$(M + H_t)v = \lambda v$$

$$\left(I - \frac{H_t}{\lambda}\right)v = \frac{Mv}{\lambda}$$

$$v = \left(I - \frac{H_t}{\lambda}\right)^{-1} \frac{Mv}{\lambda} = \frac{Mv}{\lambda} + \sum_{a=1}^{\infty} \left(\frac{H_t}{\lambda}\right)^a \frac{Mv}{\lambda}$$

- $X_t = \lambda v v^T$

$$M - X_t = M - \lambda v v^T$$

$$= M - M \frac{v v^T}{\lambda} M - \sum_{a \geq 0, b \geq 0, a+b \geq 1}^{\infty} \left(\frac{H_t}{\lambda}\right)^a \frac{M v v^T M^T}{\lambda} \left(\frac{H_t}{\lambda}\right)^b$$

# Key Step 2

$$\begin{aligned} & \|M - X_t\|_\infty \\ & \leq \|M - M \frac{vv^T}{\lambda} M\|_\infty + \sum_{a \geq 0, b \geq 0, a+b \geq 1}^\infty \left| \left(\frac{H_t}{\lambda}\right)^a \frac{Mvv^T M^T}{\lambda} \left(\frac{H_t}{\lambda}\right)^b \right|_\infty \end{aligned}$$

$M = u^* u^{*T}$

- $M = u^* u^{*T}$

$$\begin{aligned} \|M - M \frac{vv^T}{\lambda} M\|_\infty & \leq \max_{i,j} e_i^T u^* \left( 1 - u^{*T} \frac{vv^T}{\lambda} u^* \right) u^{*T} e_j \\ & \leq \max_{i,j} |e_i^T u^*| |e_j^T u^*| |1 - (u^{*T} v)^2 / \lambda| \\ & \leq \frac{\mu^2}{n} 4 \|H_t\|_2 \leq 8\mu^2 \delta \|M - X_{t-1}\|_\infty \end{aligned}$$

# Key Step 3

- Need to bound

$$\| (H_t)^a u^* \|_\infty$$

- $H_t = M - X_{t-1} - P_\Omega(M - X_{t-1})$
- $(H_t)^a$  has several correlated entries
  - Use technique of [Erdos et al'2013]
  - Intuitively, counts the total no. of paths between any pair of nodes
- Bound:  $\| (H_t)^a u^* \|_\infty \leq \frac{\mu}{\sqrt{n}} (\delta \|M - X_{t-1}\|_\infty c \log n)^a$
- Sum up terms to bound  $\|M - X_t\|_2$

# Guarantee for SVP

- At  $t$ -th step :

$$\|M - X_t\|_\infty \leq .5 \|M - X_{t-1}\|_\infty$$

- After  $\log\left(\frac{\mu}{\epsilon}\right)$  steps:  $\|M - X_t\|_\infty \leq \epsilon$

- Sample complexity:  $|\Omega| \geq nr^2 \mu^2 \left(\frac{\sigma_1}{\sigma_r}\right)^2 \log^2 n \log \frac{1}{\epsilon}$ 
  - Dependence on condition number!!!

# Stagewise-SVP

- $X_0 = 0$
- For  $k=1\dots r$ 
  - For  $t=1:T$ 
    - $X_t = P_r(X_{t-1} - P_\Omega(X_{t-1} - M))$
  - End For
  - $X_0 = X_T$
- End For

# Guarantees

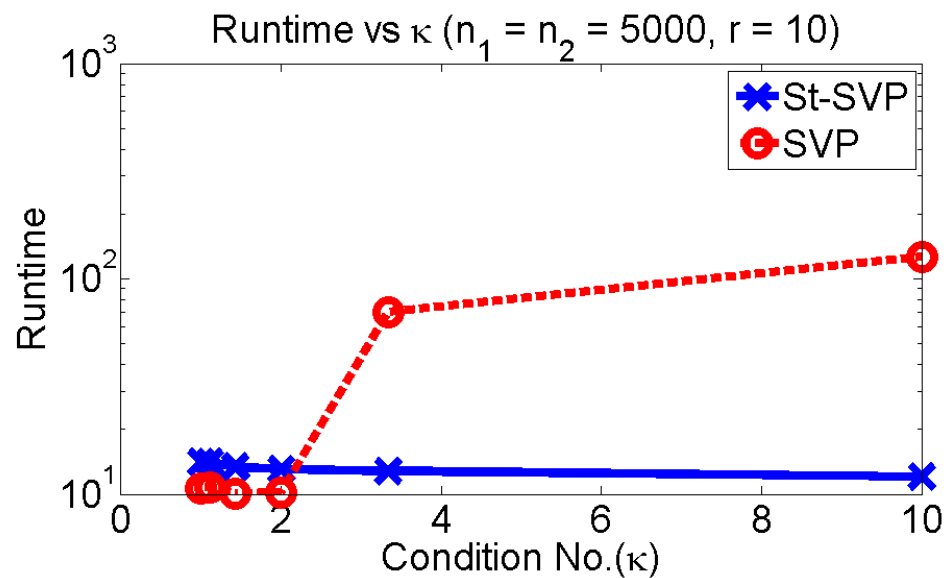
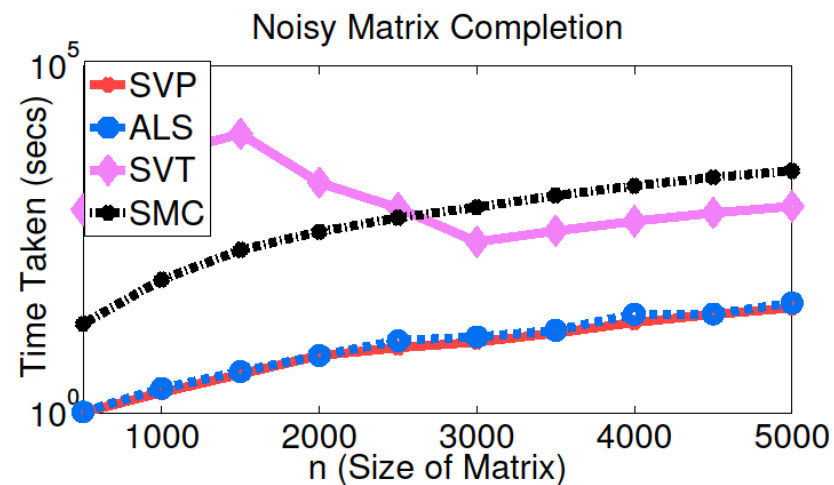
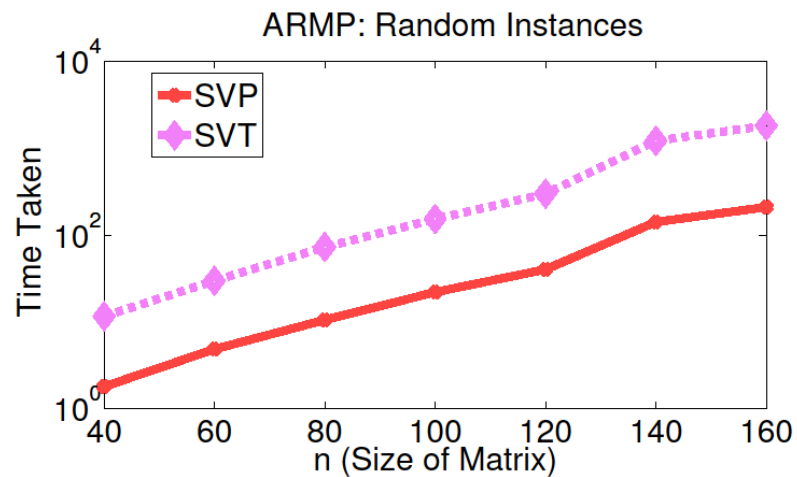
- After  $t$ -th step of  $k$ -th stage:

$$\|M - X_t\|_\infty \leq \frac{2\mu^2 r}{n} (\sigma_{k+1} + \left(\frac{1}{2}\right)^t \sigma_k)$$

- $M$ : rank- $r$  i.e.  $\sigma_{r+1} = 0$
- After  $T = \log\left(\frac{1}{\epsilon}\right)$  steps of  $r$ -th stage:  $\|M - X_T\|_\infty \leq \epsilon$
- Sample complexity:  $|\Omega| \geq nr^4 \mu^2 \log n \log 1/\epsilon$
- Computation complexity:  $O(nr^6 \mu^2 \log n \log \frac{1}{\epsilon})$ 
  - Linear in  $n$
  - No explicit dependence on  $\sigma_1/\sigma_r$



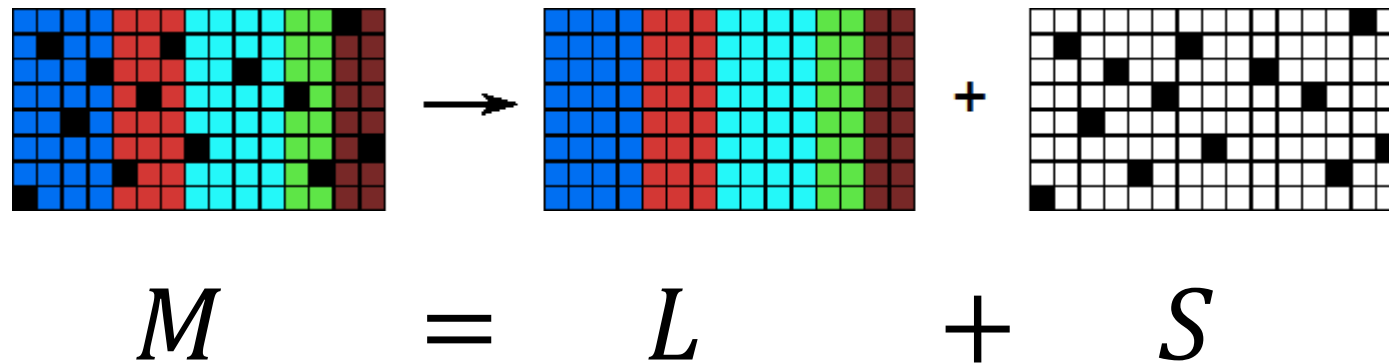
# Simulations



# Summary

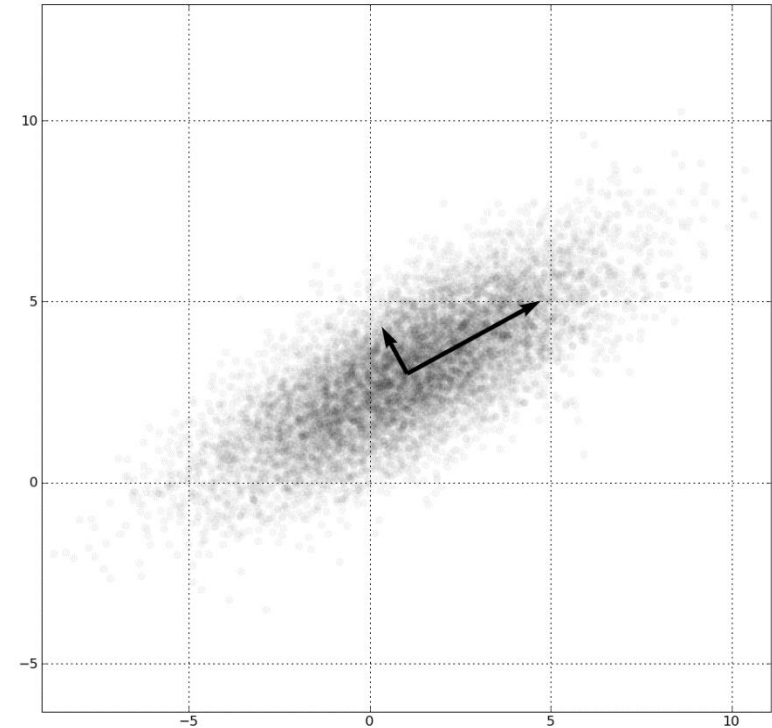
- Study matrix completion problem
- Projected gradient descent works!
- With some tweaks, obtain a nearly linear time algorithm for matrix completion
  - No explicit dependence on condition number
- Future work:
  - Remove dependence on  $\epsilon$  for sample complexity
  - AltMin: remove condition no. dependence using similar techniques?

# Robust Principal Component Analysis



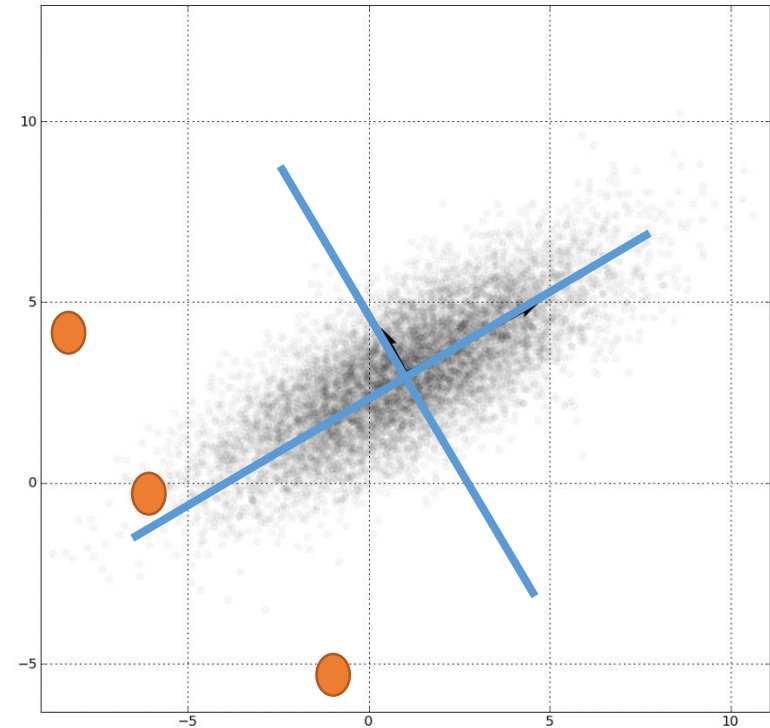
# Principal Component Analysis

- $X = [x_1 \ x_2 \ \dots \ x_n]$
- PCA: find best rank- $r$  approx. of  $X$ 
  - Top  $r$  –singular components of  $X$
  - $X_r = P_r(X)$
- $\|X - X_r\|_2 = \sigma_{r+1}$ 
  - Frobenius norm guarantees



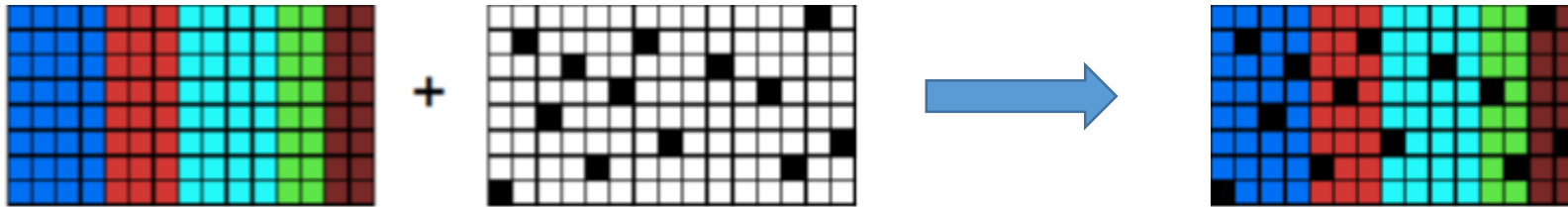
# PCA with Corruption?

- $X = [x_1 \ x_2 \ \dots \ x_n] + E$
- $\|X - P_r(X + E)\|_2 \leq \sigma_{r+1} + 2\|E\|_2$



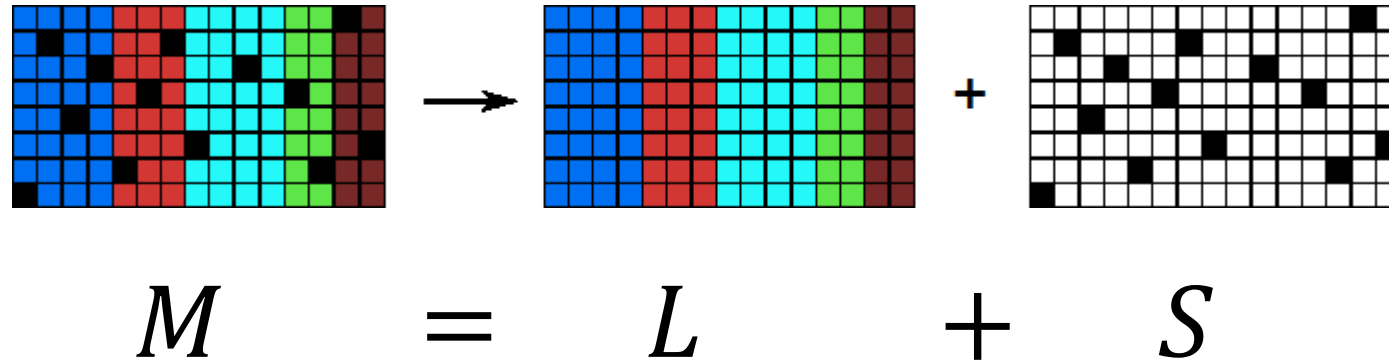
# Sparse Corruptions?

- Can we do better?
  - If  $E$  is sparse?



- E.g.
  - Each point can be corrupted in a few random co-ordinates

# Robust PCA



- $M$ : given matrix
  - $L$ : low-rank matrix
  - $S$ : sparse matrix

# Foreground + Background Separation



Original Video

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Background

+



Foreground



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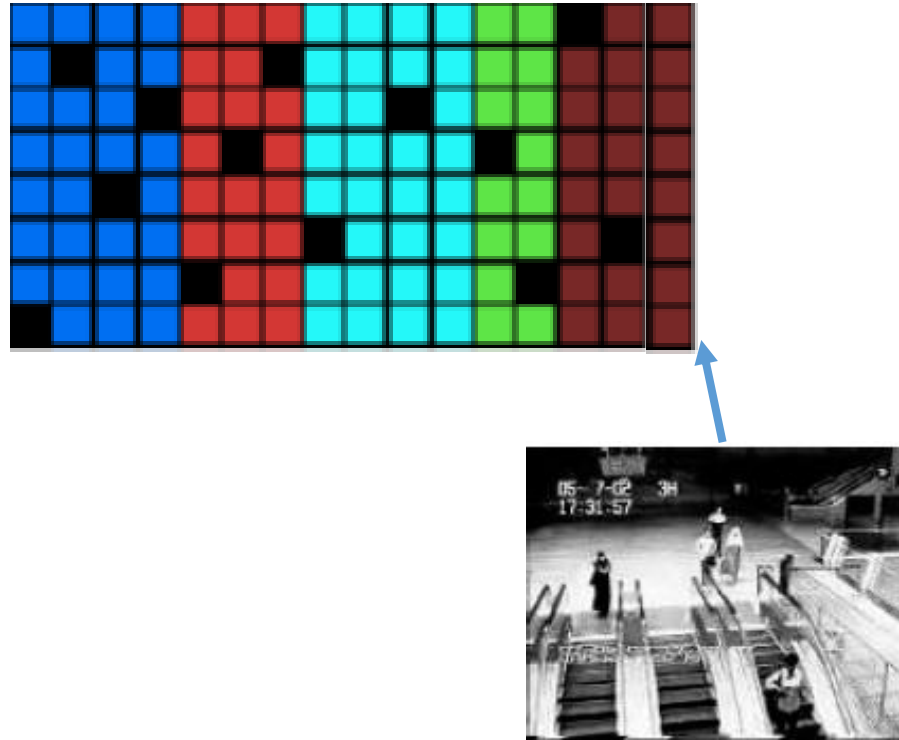
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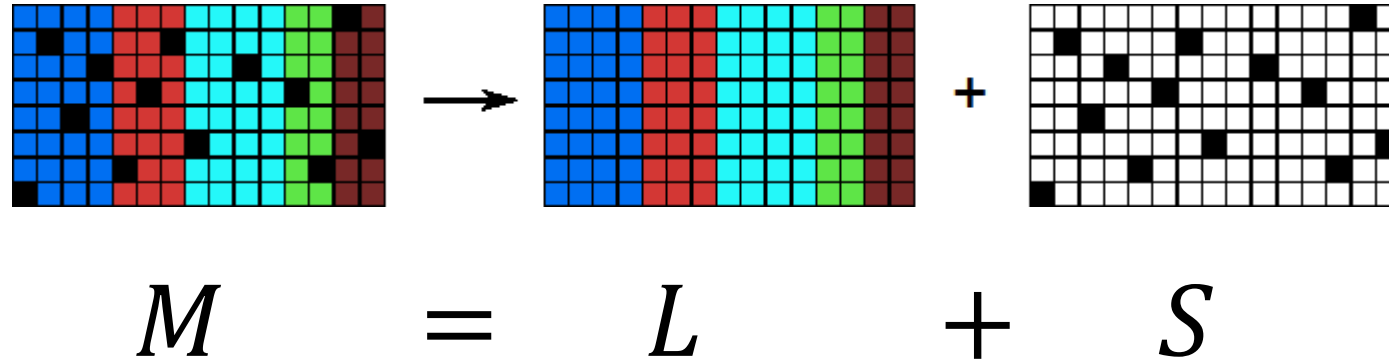


# Foreground + Background Separation

- Each  $64 \times 64$  frame: 4096-dimensional vector



# Robust PCA



- $M \in R^{n \times n}$ : given matrix
  - $L$ : low-rank matrix
  - $S$ : sparse matrix
- NP-hard problem in general

# Identifiability?

1	0	0
0	0	1
0	0	0

=

1	0	0
0	0	0
0	0	0

+

0	0	0
0	0	1
0	0	0

$M = L + S$

- Assumptions:

- $L$  is incoherent---  $L_{ij} \leq \mu \|L\|_F / n$
- $S$  is row and column sparse

# Existing Method

$$\min_{\hat{L}, \hat{S}} \text{rank}(\hat{L}) + \lambda \|\hat{S}\|_0$$
$$\text{s. t. } M = \hat{L} + \hat{S}$$

- $\|\hat{L}\|_* = \sum_i \sigma_i(\hat{L})$

- Convex program,

- Assumption

- $M$

- $s \leq \frac{n}{\mu^2 r}$

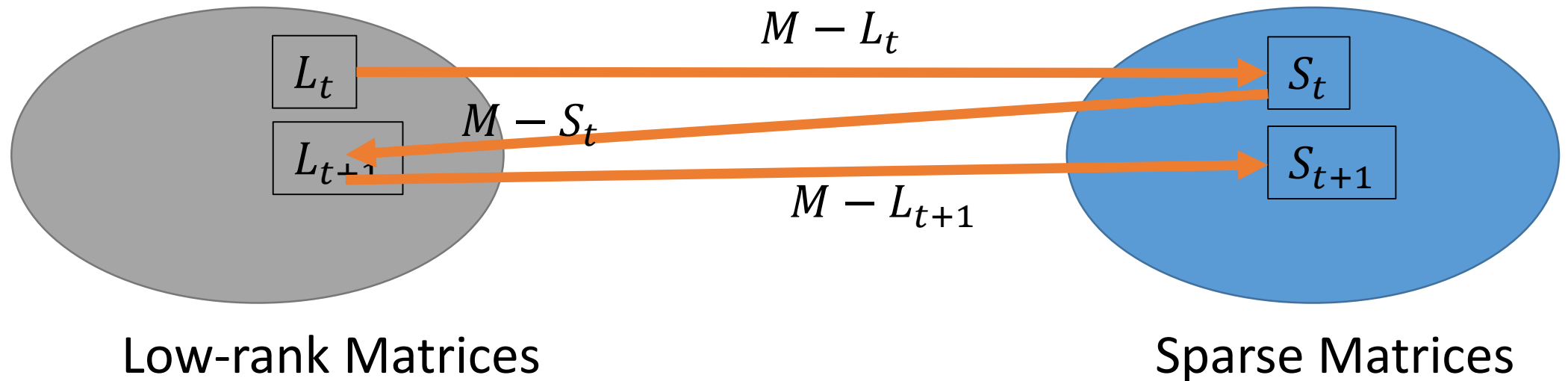
- Recover  $L, S$  [Chandrasekharan et al'2009, Candes et al'2009]

Question: PCA time complexity  
for Robust PCA?

That is,  $O(n^2 r)$  algorithm?

# Our Approach: Alternating Projections

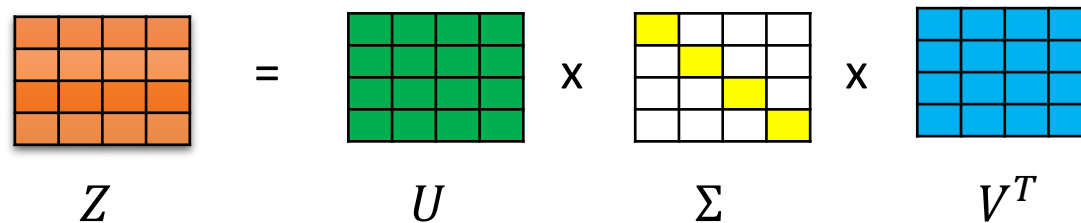
- Goal:  $M = L + S$ 
  - $L$ : low-rank matrix
  - $S$ : sparse matrix
- $M = L_t + S_t$



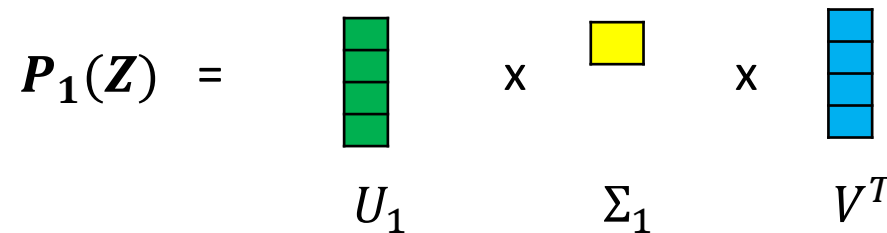
# Projection onto Low-rank Matrices

- Non-convex projections: NP-hard in general
- But  $P_r(Z)$  can be computed efficiently:

$$Z = U\Sigma V^T$$



- $P_r(Z) = U_r \Sigma_r V_r^T$




- Time complexity:  $O(n^2 r)$

# Projection onto Sparse Matrices

- Non-convex projection
- $HT_{\zeta}(Z)$ : removes all elements with magnitude smaller than  $\zeta$

1	0.1	0.22
0.1	0.01	.9
0.11	0.02	0.12

$HT_{0.5}$  

1	0	0
0	0	.9
0	0	0

# Non-convex RPCA

- $L_0 \rightarrow 0$
- $\zeta = \mu^2 r / n$
- For  $t=1, 2, \dots, T$ 
  - $\zeta = \frac{1}{4} \cdot \zeta$
  - $S_t = HT_{\zeta}(M - L_t)$
  - $L_{t+1} = P_r(M - S_t)$
- Output,  $L_T, S_T$



# Computation Time

- Each round: 1 SVD + 1 Hard Thresholding
- Time complexity per round:  $O(n^2r)$
- No. of rounds?

# Results

- After t-th step:

$$\|L - L_{t+1}\|_\infty \leq \frac{1}{2} \|L - L_t\|_\infty$$

- $T = \log\left(\frac{\|L\|_\infty}{\epsilon}\right)$ ,  $\|L_T - L\|_\infty \leq \epsilon$

- Computation complexity:  $O(n^2 r \log \frac{1}{\epsilon})$

- $O(\log \frac{1}{\epsilon})$  more expensive than PCA

- Assumption:  $M = L + S$

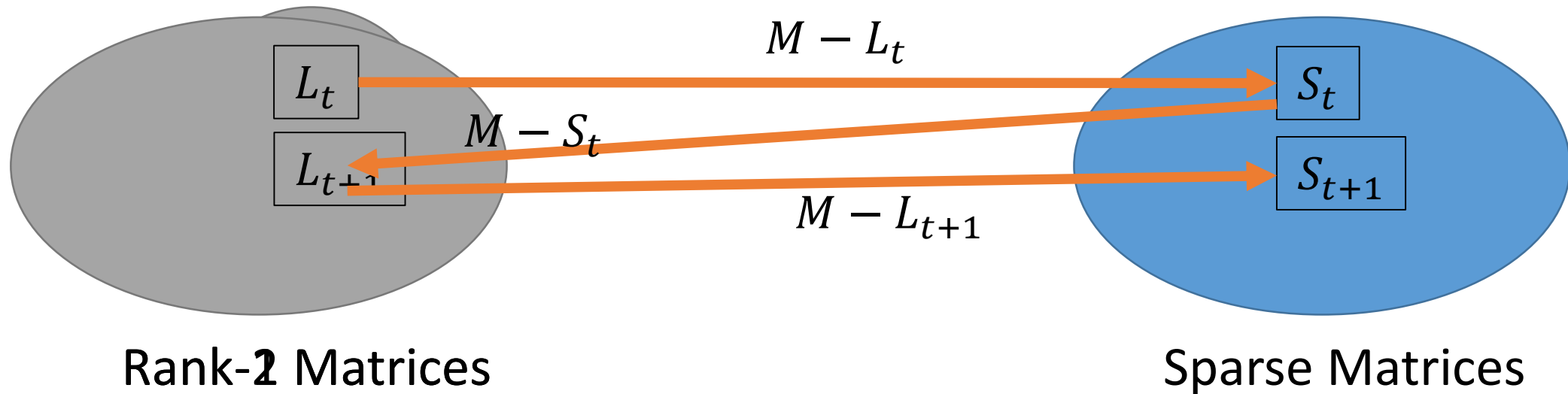
- $s \leq \frac{n}{\mu^2 r} \cdot \frac{\sigma_r^2}{\sigma_1^2}$

- Worse requirement than Hsu et al'2011

# Remove Condition No. Dependence?

- Stagewise procedure
  - k-th stage projects onto rank- $k$  matrices
  - $1 \leq k \leq r$

## 2<sup>nd</sup> Stage



# Result

- $T = \log\left(\frac{1}{\epsilon}\right)$

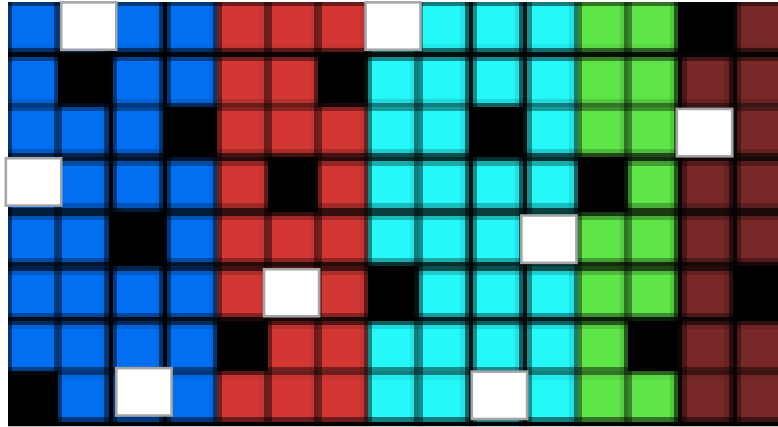
$$\|L_T - L\|_2 \leq \epsilon$$

- Assumption:  $s \leq \frac{n}{\mu^2 r}$

- $s$ : number of corrupted entries in any row or column
- Same as convex relaxation approach (Hsu et al'2011)

- Running time:  $O(n^2 r^2 \log \frac{1}{\epsilon})$

# Missing Entries?



- Assuming missing entries are corrupted entries
- Allows for  $O\left(\frac{n^2}{r}\right)$  missing entries

# Proof Technique

- $L_{t+1} = P_r(M - S_t) = P_r(L + S - S_t) = P_r(L + E_t)$
- Standard SVD guarantees:
  - $\|L_{t+1} - L\|_2 \leq \|E_t\|_2 \sim O(1)$
  - $\text{supp}(S_{t+1}) \neq \text{supp}(S)$
  - Hence,  $E_{t+1} = S - S_{t+1}$  can be dense
- Goal: ensure
  - $\text{supp}(S_{t+1}) \subseteq \text{supp}(S)$
  - $\|S - S_{t+1}\|_\infty \leq .5 \|S - S_t\|_\infty$
- But for this, we need  $\|L_{t+1} - L\|_\infty \leq .5 \|E_t\|_\infty$

# A Novel Perturbation Lemma

$$\|P_r(L + E_t) - L\|_\infty \leq .5 \|E_t\|_\infty$$

- If:
  - $E_t$ : sparse
  - $L$ : incoherent
- Much tighter than the standard matrix perturbation results
  - $\|P_r(L + E_t) - L\|_2 \leq 2\|E_t\|_2$

# Proof Sketch (Rank-1 case)

- $L = uu^T$
- $L_{t+1} = P_1(L + E_t), \quad L_{t+1} = vv^T$

$$(L + E_t)v = v$$

$$(I - E_t)v = Lv$$

$$v = (I - E_t)^{-1} Lv = Lv + \sum_{a=1}^{\infty} (E_t)^a Lv$$

$$\begin{aligned} L - L_{t+1} &= L - vv^T \\ &= L - Lv v^T L - \sum_{a \geq 0, b \geq 0, a+b \geq 1} (E_t)^a Lv v^T L^T (E_t)^b \end{aligned}$$



# Proof Sketch

- Using  $L = uu^T$

$$L - L_t = (1 - \langle u, v \rangle^2)L + \langle u, v \rangle^2 \sum_{a \geq 0, b \geq 0, a+b \geq 1}^{\infty} (E_t)^a uu^T (E_t)^b$$

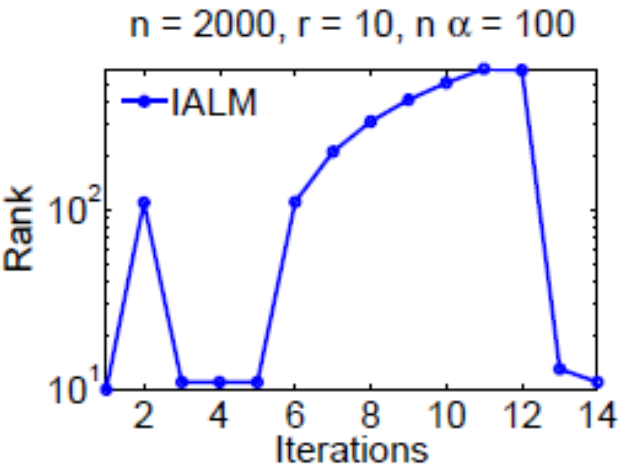
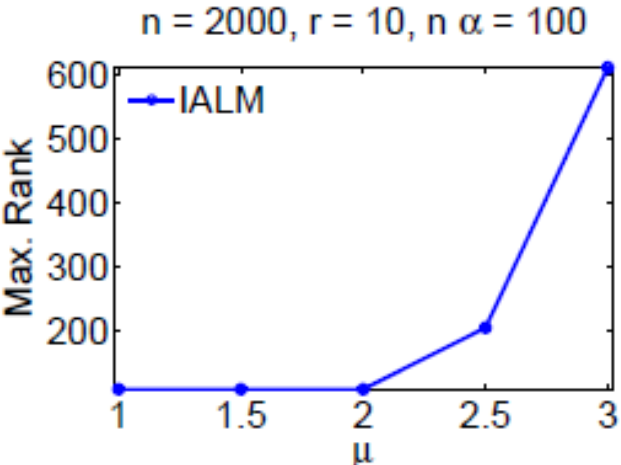
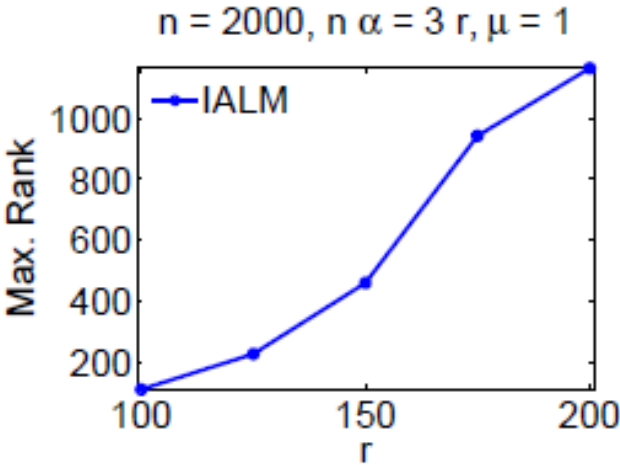
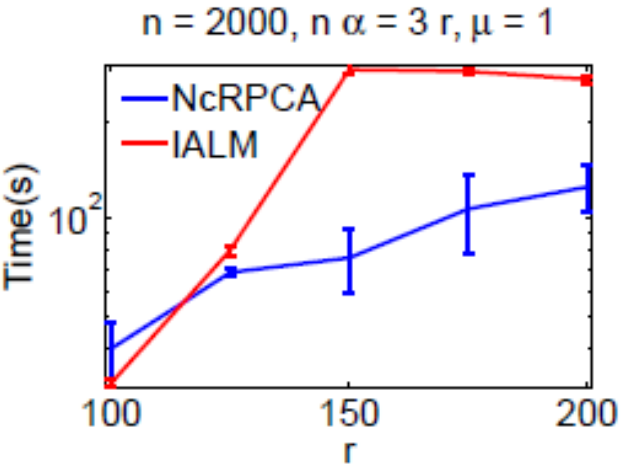
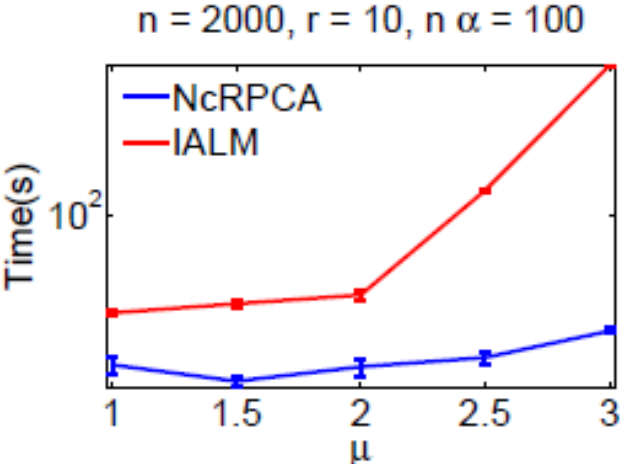
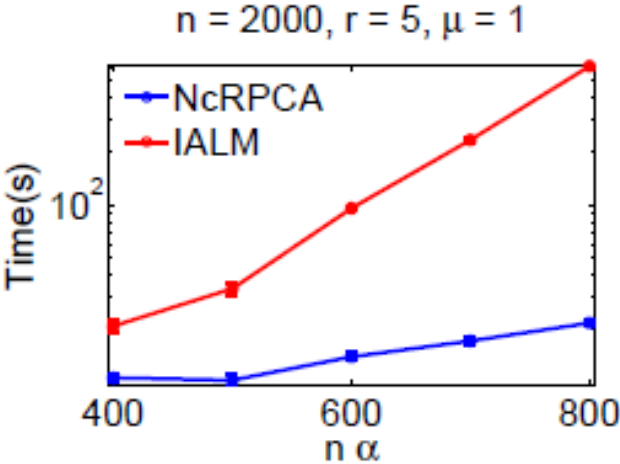
- $\|(E_t)^a u\|_{\infty}$ : small
  - $E_t$ : sparse
  - $u$ : incoherent ( $\|u\|_{\infty} \leq \mu/\sqrt{n}$ )
- Bound  $\|\cdot\|_{\infty}$  of each term







# Empirical Results (Synthetic Datasets)



# Empirical Results

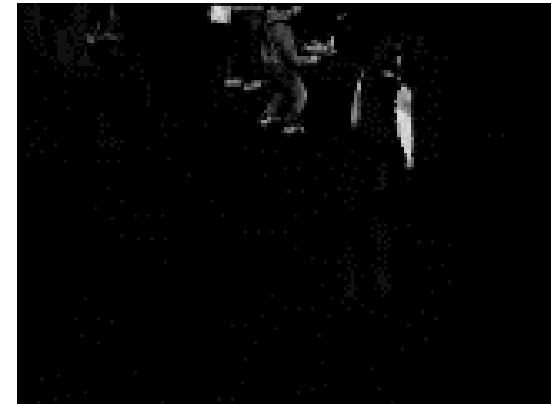
Convex Method. Runtime: 1700 sec



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+



Non-Convex Method. Runtime: 70 sec



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# Summary

- Robust PCA
  - Low-rank+Sparse Decomposition
- Alternating Projection Method
- Under standard assumptions
  - Linear rate of convergence
  - Computation time: Recovery in  $O(\text{PCA})$ , for constant rank matrices
- Key analysis tool: a strong perturbation bound for SVD

# Future Work

- RIP/RSC based Matrix sensing:
  - Necessity of the required RIP/RSC conditions?
- Matrix completion:
  - Remove dependence of  $|\Omega|$  on error  $\epsilon$
  - Optimal dependence of  $|\Omega|$  on  $r$
- Robust PCA:
  - Extension to [Candes et al'09] style conditions
  - Can handle  $O(\frac{n}{\mu^2})$  corruptions per row (currently,  $O(\frac{n}{\mu^2 r})$ )
- Develop a more generic framework to jointly analyze these problems
  - Similar to unified M-estimator technique of [Negahban et al'09]



Thanks!