# Question 1

# Framework for Strong Liveness Analysis

## Direction of the analysis

Backward

#### Lattice and Meet

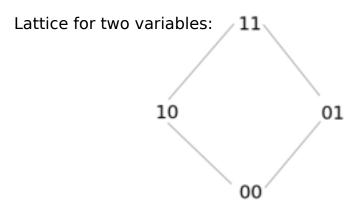


Figure 1: lattice

Meet operation for this frame work will be similar as live variables ( $\cup$ ) Note: In strong liveness analysis lattice is non-seperable.

### Gen and Kill Computation:

- if(e) or print(e)
  - Gen =  $\{v \mid v \in \text{operand}(e)\}\$ , e is an expression.
  - $\text{ Kill} = \phi$
- **•** y = e
  - $\begin{array}{l} \ Gen = if \ y \in Out(S) \\ \text{then } \{v \mid v \in operand(e)\} \\ \text{else } \phi \ , \ \text{where e is an expression} \end{array}$
  - Kill = if  $y \in \text{operand}(e)$ then  $\phi$ else y

# Question 2

No, the result of strong liveness analysis is not similar as iterative liveness analysis. Figure 2 represents a counter example. Live variable analysis will mark x as live at Out(Stmt) and hence no dead code elimination where as Strong liveness analysis will mark x as dead variable at Out(Stmt) and statement will be eliminated during dead code elimination.

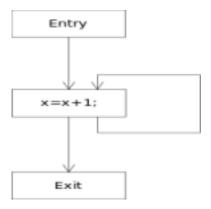


Figure 2: Counter Example