Public-Key Cryptography

Lecture 9

Public-Key Encryption

Diffie-Hellman Key-Exchange

PKE scheme

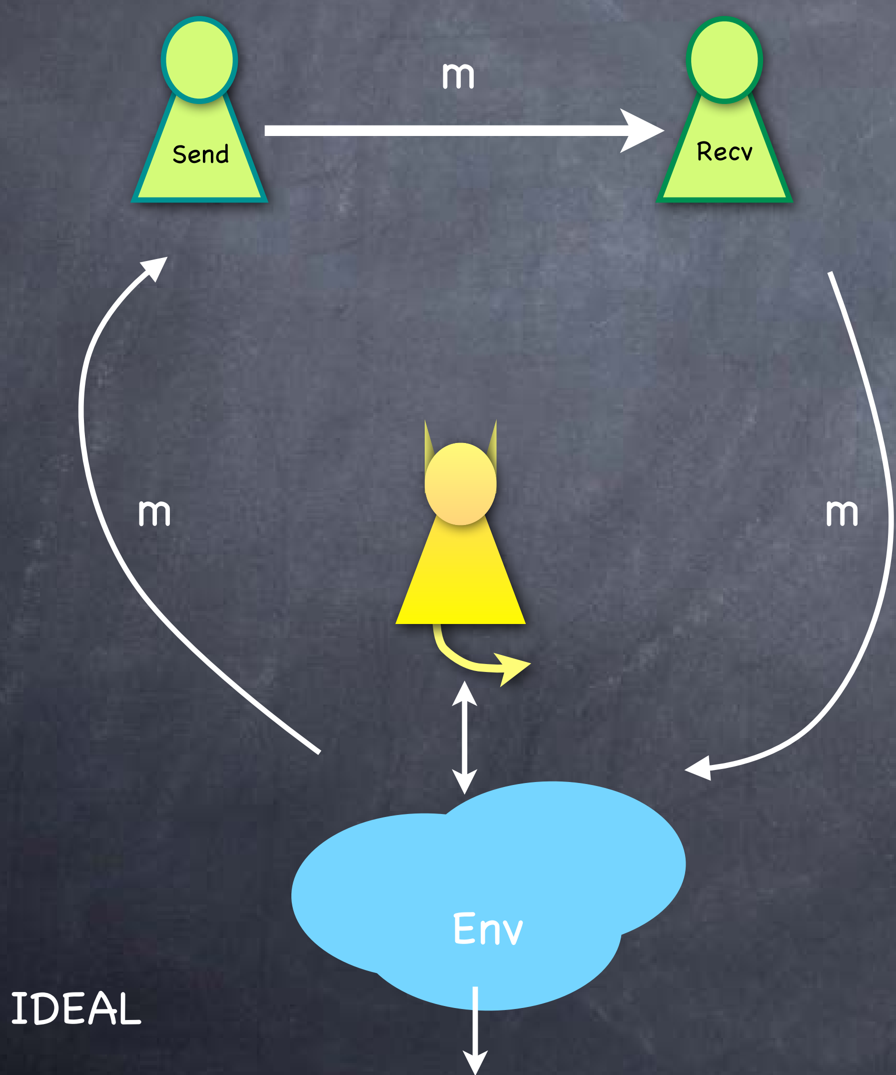
Shared/Symmetric-Key Encryption (a.k.a. private-key encryption)

- SKE:
  - Syntax
    - KeyGen outputs $K \leftarrow \mathcal{K}$
    - Enc: $\mathcal{M} \times \mathcal{K} \times \mathcal{R} \rightarrow \mathcal{C}$
    - Dec: $\mathcal{C} \times \mathcal{K} \rightarrow \mathcal{M}$
  - Correctness
    - $\forall K \in \text{Range(KeyGen)}$, $\text{Dec}(\text{Enc}(m,K), K) = m$
  - Security (SIM/IND-CPA)

a.k.a. asymmetric-key encryption

- PKE
  - Syntax
    - KeyGen outputs $(PK,SK) \leftarrow \mathcal{PK} \times \mathcal{SK}$
    - Enc: $\mathcal{M} \times \mathcal{PK} \times \mathcal{R} \rightarrow \mathcal{C}$
    - Dec: $\mathcal{C} \times \mathcal{SK} \rightarrow \mathcal{M}$
  - Correctness
    - $\forall (PK,SK) \in \text{Range(KeyGen)}$, $\text{Dec}(\text{Enc}(m,PK), SK) = m$
  - Security (SIM/IND-CPA, PKE version)

SIM-CPA (PKE Version)



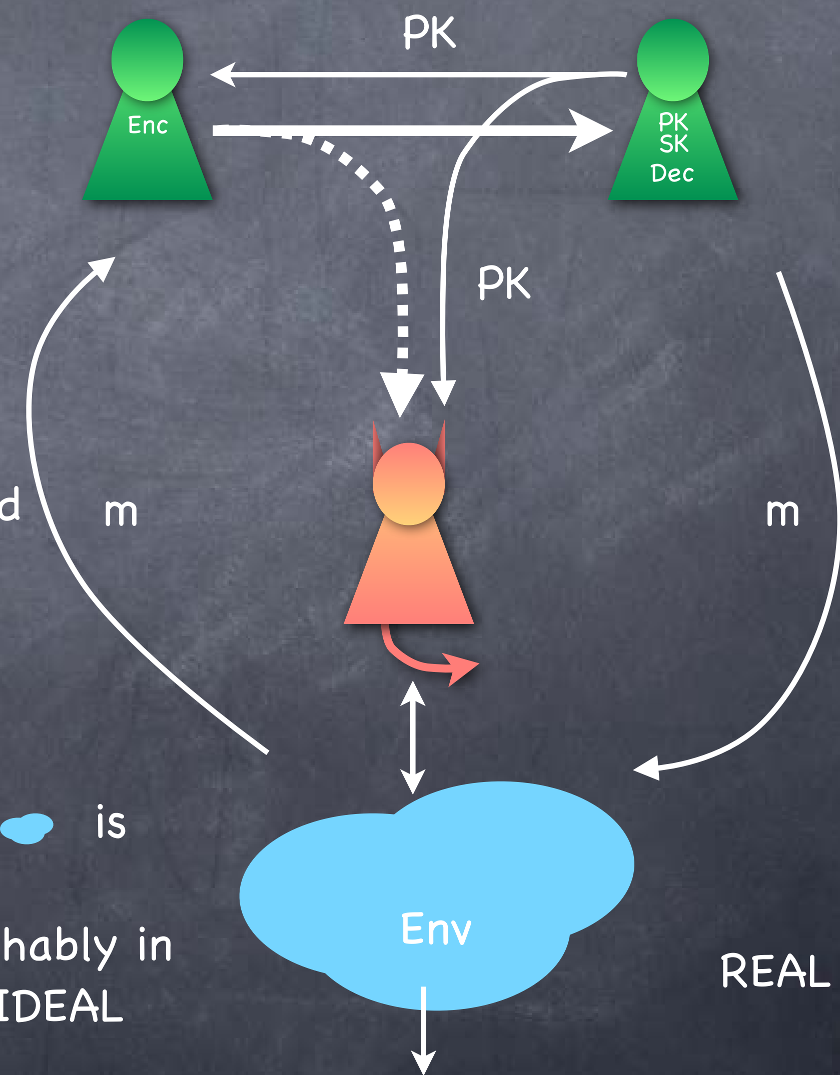
Secure (and correct) if:

∀ 

∃  s.t.

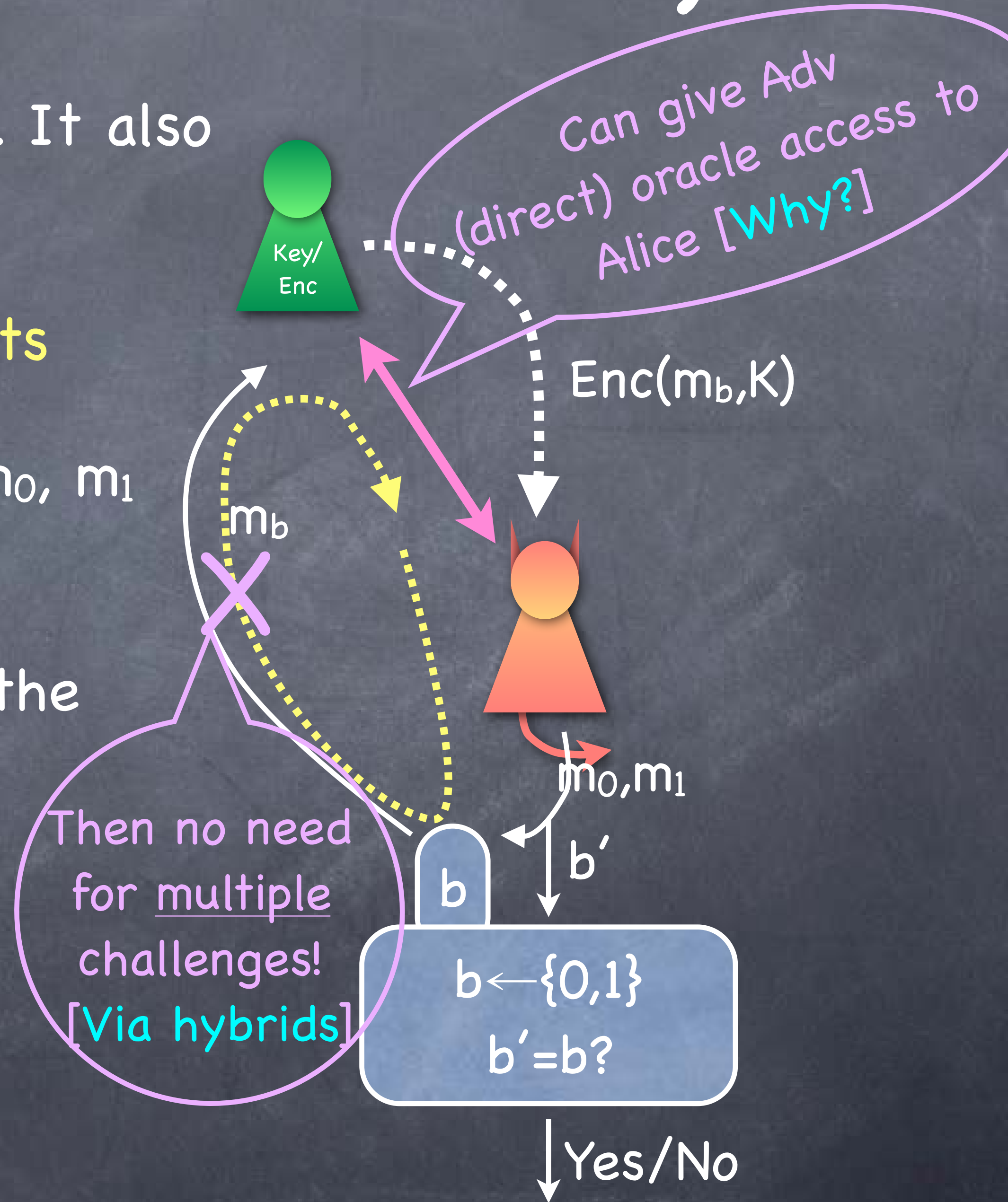
∀ 

output of is distributed indistinguishably in REAL and IDEAL



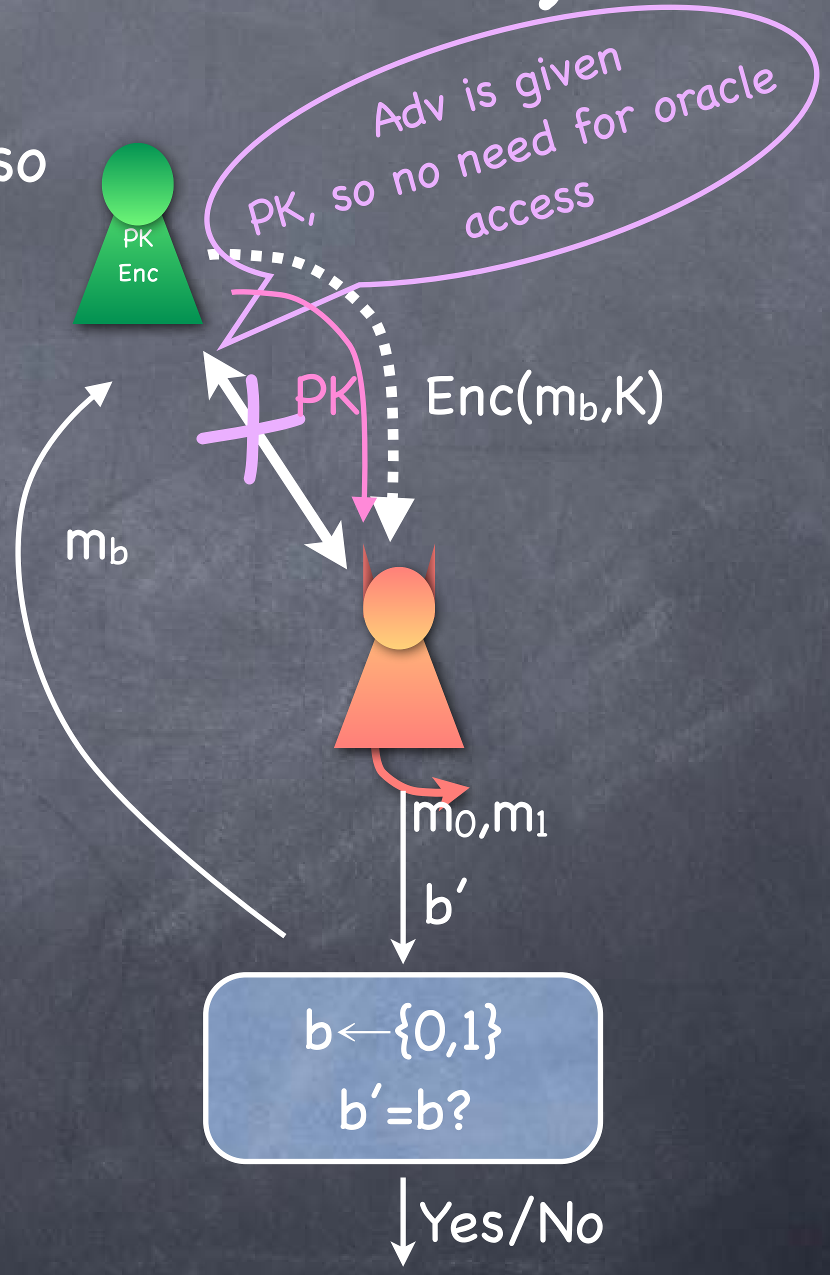
IND-CPA (SKE version)

- Experiment picks a random bit b. It also runs KeyGen to get a key K
  - For as long as Adversary wants
    - Adv sends two messages m_0, m_1 to the experiment
    - Expt returns Enc(m_b,K) to the adversary
  - Adversary returns a guess b′
  - Experiment outputs 1 iff b′=b
- IND-CPA secure if for all PPT adversaries Pr[b′=b] – 1/2 ≤ ν(k)



IND-CPA (~~SKE~~ PKE version)

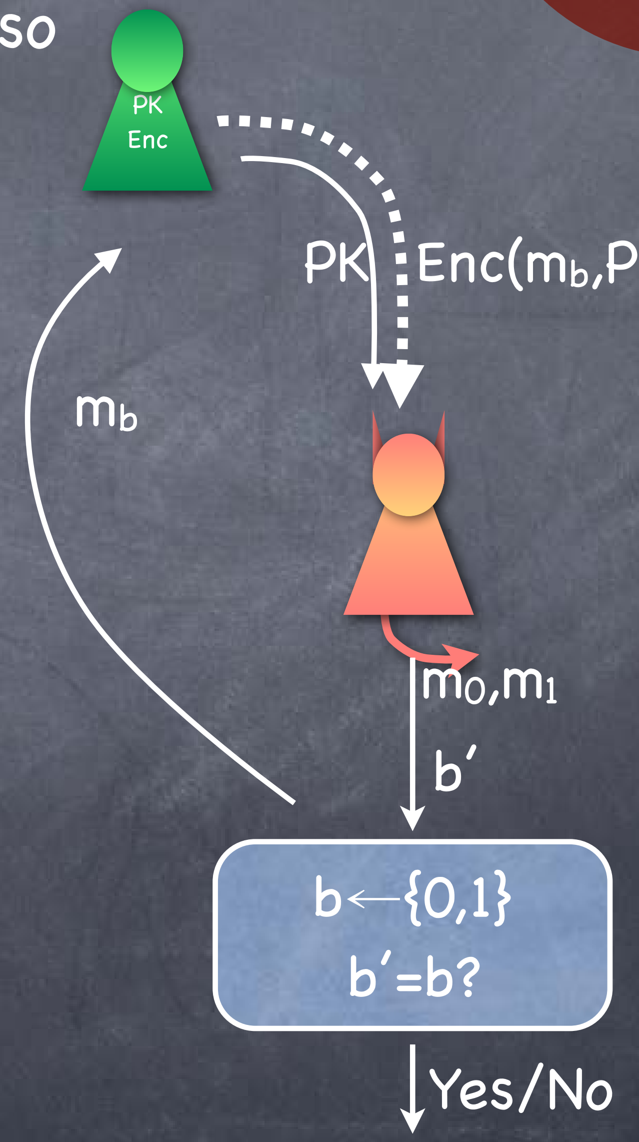
- Experiment picks a random bit b. It also runs KeyGen to get a key (PK,SK). Adv given PK
  - Adv sends two messages m_0, m_1 to the experiment
  - Expt returns Enc(m_b,K) to the adversary
  - Adversary returns a guess b′
  - Experiment outputs 1 iff b′=b
- IND-CPA secure if for all PPT adversaries Pr[b′=b] – 1/2 ≤ ν(k)



IND-CPA (PKE version)

IND-CPA + ~correctness equivalent to SIM-CPA

- Experiment picks a random bit b. It also runs KeyGen to get a key (PK,SK). Adv given PK
  - Adv sends two messages m_0, m_1 to the experiment
  - Expt returns Enc(m_b,K) to the adversary
  - Adversary returns a guess b′
  - Experiment outputs 1 iff b′=b
- IND-CPA secure if for all PPT adversaries Pr[b′=b] – 1/2 ≤ ν(k)



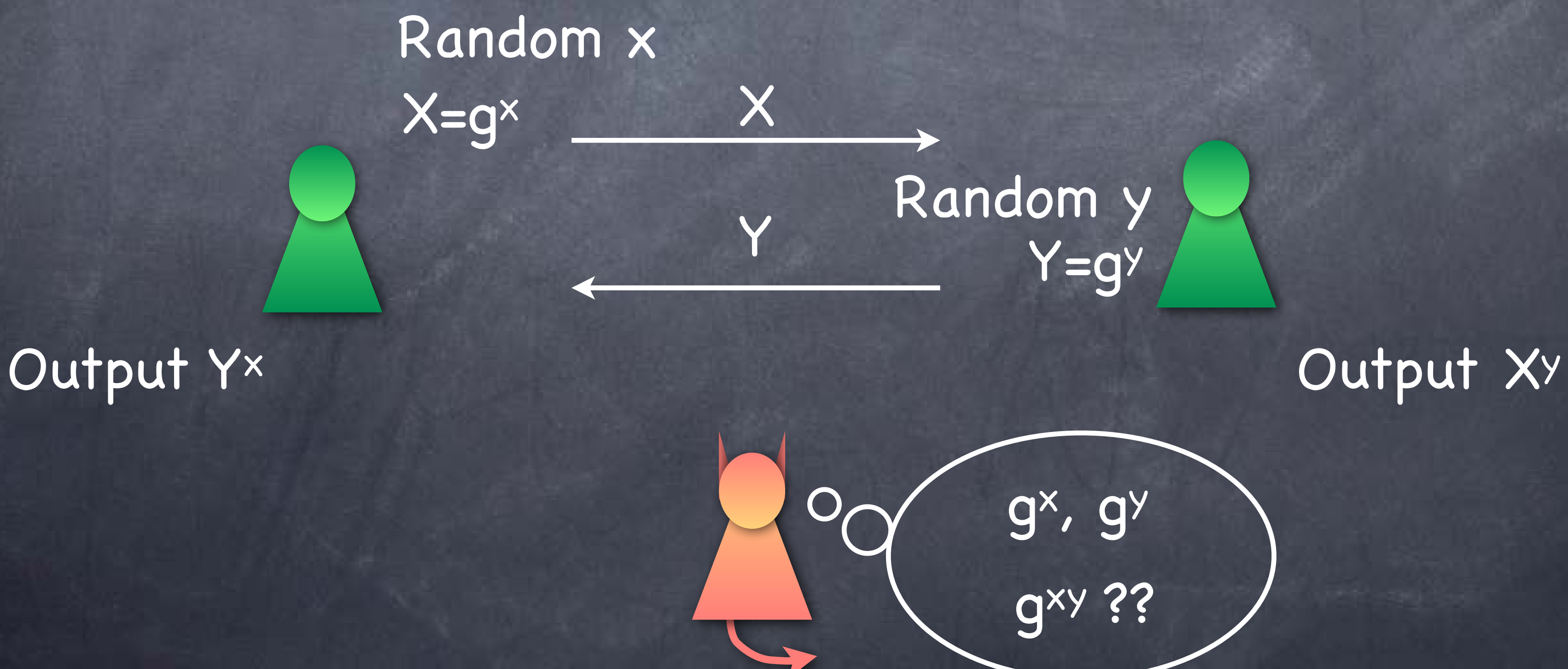
Perfect Secrecy?

- No perfectly secret and correct PKE (even for one-time encryption)
  - Public-key and ciphertext (the total shared information between Alice and Bob at the end) should together have entire information about the message
    - Intuition: If Eve thinks Bob could decrypt it as two messages based on different SKs, Alice should be concerned too
    - i.e., Alice conveys same information to Bob and Eve
    - [Exercise]
- PKE only with computational security

Unless assumptions of imperfect eavesdropping

Diffie-Hellman Key-exchange

- A candidate for how Alice and Bob could generate a shared key, which is "hidden" from Eve



Why DH-Key-exchange could be secure

- Given g^x, g^y for random x, y, g^{xy} should be "hidden"
  - i.e., could still be used as a pseudorandom element
  - i.e., $(g^x, g^y, g^{xy}) \approx (g^x, g^y, R)$
- Is that reasonable to expect?
  - Depends on the "group"

Groups, by examples

- A group (G, $*$) specified by a set G (for us finite, unless otherwise specified) and a "group operation" $*$ that is associative, has an identity, is invertible, and (for us) commutative

Abelian

- Examples: $\mathbb{Z}$ = (integers, +) (this is an infinite group),

  $\mathbb{Z}_N$ = (integers modulo N, + mod N),

  G^n = (Cartesian product of a group G, coordinate-wise operation)

Direct Product

- Order of a group G: |G| = number of elements in G

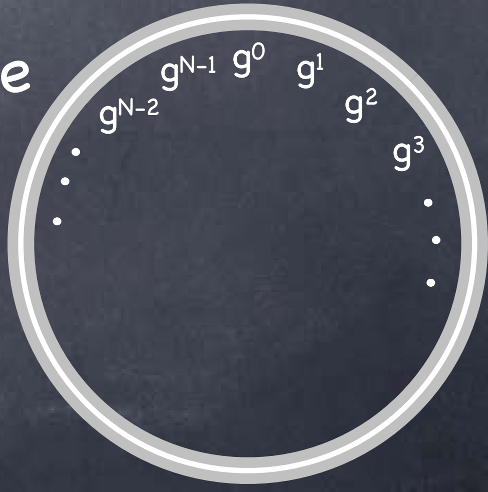
Lagrange's theorem

- For any a∈G, $a^{|G|} = a * a * \ldots * a$ (|G| times) = identity

- Finite Cyclic group (in multiplicative notation): there is one element g such that G = $\{g^0, g^1, g^2, \ldots g^{|G|-1}\}$

  - Prototype: $\mathbb{Z}_N$ (additive group), with g=1

    - or any g s.t. gcd(g,N) = 1



Groups, by examples

- $\mathbb{Z}_N^*$ = (generators of $\mathbb{Z}_N$, multiplication mod N)
  - Numbers in {1,..,N-1} which have a multiplicative inverse mod N
  - **Fact**: If N is prime, $\mathbb{Z}_N^*$ is a cyclic group, of order N-1
    - e.g. $\mathbb{Z}_5^*$ = {1,2,3,4} is generated by 2 (as 1,2,4,3), and by 3 (as 1,3,4,2). But 1 and 4 are not generators.
    - (Also cyclic for certain other values of N)

Generators are called Primitive Roots of N

Discrete Log Assumption

- Discrete Log (w.r.t g) in a (multiplicative) cyclic group G generated by g: $DL_g(X)$:= unique x such that $X = g^x$ ($x \in \{0,1,\ldots,|G|-1\}$)
- In a (computationally efficient) group, given integer x and the standard representation of a group element g, can efficiently find the standard representation of $X=g^x$ (How?)
  - But given X and g, may not be easy to find x (depending on G)
  - DLA: Every PPT Adv has negligible success probability in the DL Expt: $(G,g) \leftarrow$ GroupGen; $X \leftarrow G$; $Adv(G,g,X) \rightarrow z$; $g^z=X$?
- If DLA broken, then Diffie-Hellman key-exchange broken
  - Eve gets x, y from g^x, g^y (sometimes) and can compute g^{xy} herself
    - A "key-recovery" attack
  - Note: could potentially break pseudorandomness without breaking DLA too

Repeated squaring

OWF collection:
Raise(x;G,g) = $(g^x;G,g)$

Decisional Diffie-Hellman (DDH) Assumption

- $\{(g^x, g^y, g^{xy})\}_{(G,g)\leftarrow\textbf{GroupGen};\ x,y\leftarrow[|G|]} \approx \{(g^x, g^y, g^r)\}_{(G,g)\leftarrow\textbf{GroupGen};\ x,y,r\leftarrow[|G|]}$
- At least as strong as DLA
  - If DDH assumption holds, then DLA holds [Why?]
- But possible that DLA holds and DDH assumption doesn't
  - e.g.: DLA is widely assumed to hold in $\mathbb{Z}_p^*$ (p prime), but DDH assumption doesn't hold there!
    - Next time