

Sets & Relations

Basics of Sets



Relational Database

x	y	Likes(x,y)
Alice	Alice	TRUE
	Jabberwock	FALSE
	Flamingo	TRUE
Jabberwock	Alice	FALSE
	Jabberwock	TRUE
	Flamingo	FALSE
Flamingo	Alice	FALSE
	Jabberwock	FALSE
	Flamingo	TRUE

Relational DB Table

Likes	
x	y
Alice	Alice
Alice	Flamingo
Jabberwock	Jabberwock
Flamingo	Flamingo

- Queries to the DB are set/logical operations
 - ```

SELECT x
FROM Likes
WHERE y='Alice' OR y='Flamingo'

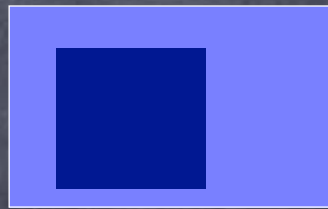
```
  - $\{ x \mid (x, Alice) \in Likes \} \cup \{ x \mid (x, Flamingo) \in Likes \}$

# Sets: Basics

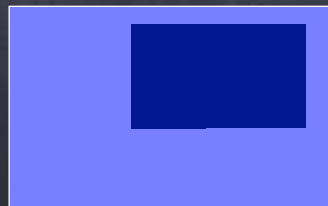
- Unordered collection of “elements”
  - e.g.:  $\mathbb{Z}$ ,  $\mathbb{R}$  (infinite sets),  $\emptyset$  (empty set),  $\{1, 2, 5\}$ , ...
- Will always be given an implicit or explicit universe (universal set) from which the elements come
  - (Aside: In developing the foundations of mathematics, often one starts from “scratch”, using only set theory to create the elements themselves)
- Set membership: e.g.  $0.5 \in \mathbb{R}$ ,  $0.5 \notin \mathbb{Z}$ ,  $\emptyset \notin \mathbb{Z}$
- Set inclusion: e.g.  $\mathbb{Z} \subseteq \mathbb{R}$ ,  $\emptyset \subseteq \mathbb{Z}$
- Set operations: complement, union, intersection, difference



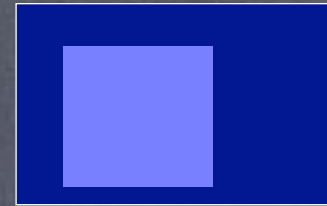
# Set Operations



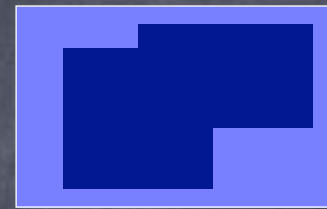
S



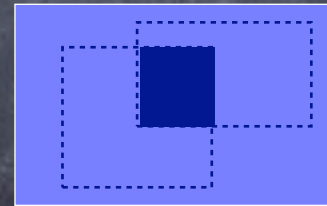
T



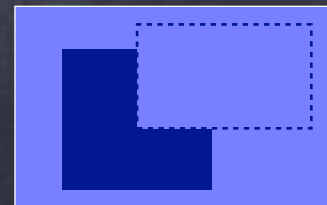
$\bar{S}$



$S \cup T$



$S \cap T$



$S - T$

# Sets as Predicates

| x          | Winged(x) | Flies(x) | Pink(x) | inClub(x) |
|------------|-----------|----------|---------|-----------|
| Alice      | FALSE     | FALSE    | FALSE   | TRUE      |
| Jabberwock | TRUE      | TRUE     | FALSE   | FALSE     |
| Flamingo   | TRUE      | TRUE     | TRUE    | TRUE      |

- Given a predicate can define the set of elements for which it holds
  - $\text{WingedSet} = \{ x \mid \text{Winged}(x) \} = \{\text{J'wock}, \text{Flamingo}\}$
  - $\text{FliesSet} = \{ x \mid \text{Flies}(x) \} = \{\text{J'wock}, \text{Flamingo}\}$
  - $\text{PinkSet} = \{ x \mid \text{Pink}(x) \} = \{\text{Flamingo}\}$
- Conversely, given a set, can define a **membership predicate** for it  
e.g. given set  $\text{Club} = \{\text{Alice}, \text{Flamingo}\}$ . Then, define predicate  $\text{inClub}(x)$  s.t.  $\text{inClub}(x) = \text{True}$  iff  $x \in \text{Club}$

# Set Operations

Unary operator

Binary operators

Associative

S complement  
Symbol:  $\bar{S}$

$$\text{in}\bar{S}(x) \equiv \neg \text{in}S(x)$$

S union T  
Symbol:  $S \cup T$

$$\begin{aligned} \text{in}S \cup T(x) \\ \equiv \text{in}S(x) \vee \text{in}T(x) \end{aligned}$$

S intersection T  
Symbol:  $S \cap T$

$$\begin{aligned} \text{in}S \cap T(x) \\ \equiv \text{in}S(x) \wedge \text{in}T(x) \end{aligned}$$

S difference T  
Symbol:  $S - T$   
(Alternately:  $S \setminus T$ )

$$\begin{aligned} \text{in}S - T(x) \\ \equiv \text{in}S(x) \wedge \neg \text{in}T(x) \\ \equiv \text{in}S(x) \nrightarrow \text{in}T(x) \end{aligned}$$

$$S - T = S \cap \bar{T}$$

S symmetric diff. T  
Symbol:  $S \Delta T$

$$\begin{aligned} \text{in}S \Delta T(x) \\ \equiv \text{in}S(x) \oplus \text{in}T(x) \end{aligned}$$

Note: Notation  $\text{in}S(x)$  used only to explicate the connection with predicate logic. Will always write  $x \in S$  later.

# De Morgan's Laws

$$\overline{S \cup T} = \bar{S} \cap \bar{T}$$

$$x \in \overline{S \cup T} \equiv \neg(x \in S \cup T)$$

$$\equiv \neg(x \in S \vee x \in T) \equiv \neg(x \in S) \wedge \neg(x \in T)$$

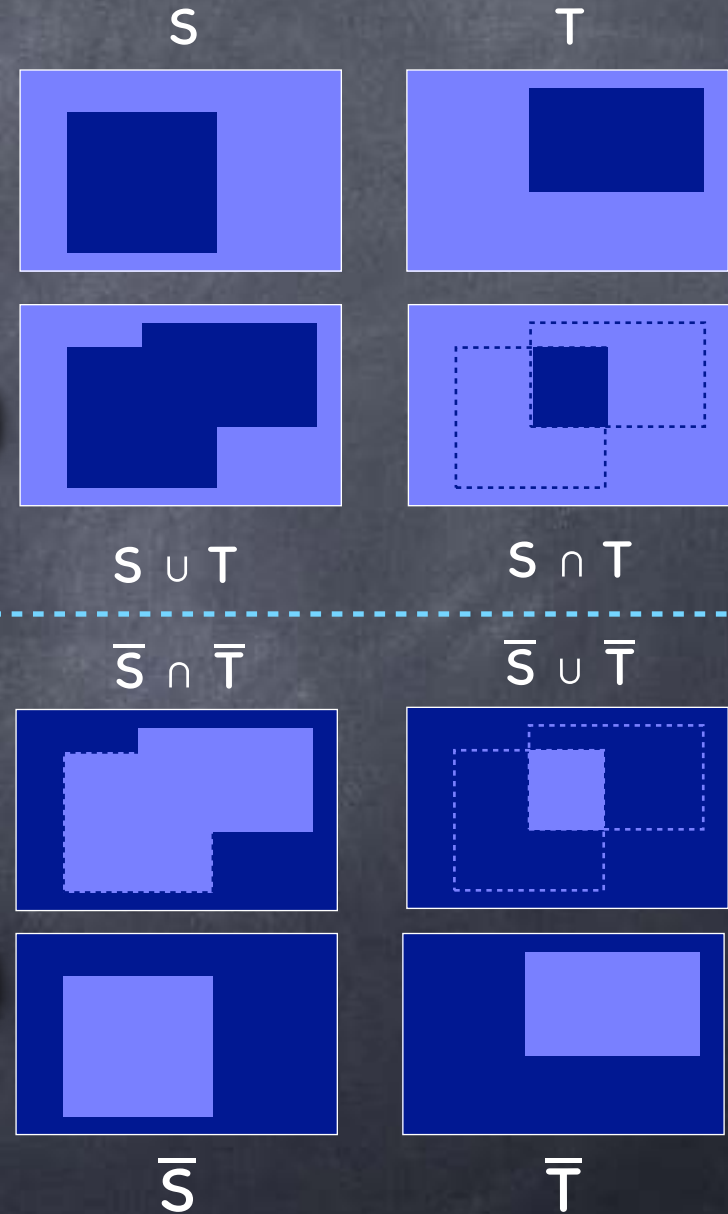
$$\equiv x \in \bar{S} \wedge x \in \bar{T} \equiv x \in \bar{S} \cap \bar{T}$$

$$\overline{S \cap T} = \bar{S} \cup \bar{T}$$

$$x \in \overline{S \cap T} \equiv \neg(x \in S \cap T)$$

$$\equiv \neg(x \in S \wedge x \in T) \equiv \neg(x \in S) \vee \neg(x \in T)$$

$$\equiv x \in \bar{S} \vee x \in \bar{T} \equiv x \in \bar{S} \cup \bar{T}$$





# Distributivity

- $R \cap (S \cup T) = (R \cap S) \cup (R \cap T)$

- $x \in R \cap (S \cup T) \equiv$

- $\equiv x \in R \wedge (x \in S \vee x \in T) \equiv (x \in R \wedge x \in S) \vee (x \in R \wedge x \in T)$

- $\equiv x \in (R \cap S) \cup (R \cap T)$

- $R \cup (S \cap T) = (R \cup S) \cap (R \cup T)$

- $x \in R \cup (S \cap T) \equiv$

- $\equiv x \in R \vee (x \in S \wedge x \in T) \equiv (x \in R \vee x \in S) \wedge (x \in R \vee x \in T)$

- $\equiv x \in (R \cup S) \cap (R \cup T)$



# Set Inclusion

| x          | Winged(x) | Flies(x) | Pink(x) |
|------------|-----------|----------|---------|
| Alice      | FALSE     | FALSE    | FALSE   |
| Jabberwock | TRUE      | TRUE     | FALSE   |
| Flamingo   | TRUE      | TRUE     | TRUE    |

•  $\text{PinkSet} \subseteq \text{FliesSet} = \text{WingedSet}$

•  $S \subseteq T$  same as the proposition  $\forall x \ x \in S \rightarrow x \in T$

•  $S \supseteq T$  same as the proposition  $\forall x \ x \in S \leftarrow x \in T$

•  $S = T$  same as the proposition  $\forall x \ x \in S \leftrightarrow x \in T$

# Set Inclusion

•  $S \subseteq T$  same as the proposition  $\forall x \ x \in S \rightarrow x \in T$

• If  $S = \emptyset$ , and  $T$  any arbitrary set,  $S \subseteq T$

•  $\forall x$ , vacuously we have  $x \in S \rightarrow x \in T$

• **If  $S \subseteq T$  and  $T \subseteq R$ , then  $S \subseteq R$**

If no such  $x$ , already done

• Consider arbitrary  $x \in S$ . Since  $S \subseteq T$ ,  $x \in T$ . Then since  $T \subseteq R$ ,  $x \in R$ .

•  $S \subseteq T \iff \bar{T} \subseteq \bar{S}$

•  $\forall x \ \underline{x \in S \rightarrow x \in T} \equiv \forall x \ \underline{x \notin T \rightarrow x \notin S}$  (contrapositive)

$\equiv \forall x \ \underline{x \in \bar{T} \rightarrow x \in \bar{S}}$

# Proving Set Equality

• To prove  $S = T$ , show  $S \subseteq T$  and  $T \subseteq S$

• e.g.,  $L(a,b) = \{ x : \exists u,v \in \mathbb{Z} \ x=au+bv \}$

$$M(a,b) = \{ x : (\gcd(a,b) \mid x) \}$$

• [Recall] **Theorem:**  $L(a,b) = M(a,b)$

• Proof in two parts:

•  $L(a,b) \subseteq M(a,b)$  : i.e.,  $\forall x \in \mathbb{Z} \ x \in L(a,b) \rightarrow x \in M(a,b)$

•  $M(a,b) \subseteq L(a,b)$  : i.e.,  $\forall x \in \mathbb{Z} \ x \in M(a,b) \rightarrow x \in L(a,b)$

First show that

$$g \in L(a,b)$$

(as the smallest +ve  
element in  $L(a,b)$ )

Let  $x=ng$ . But

$$g=au+bv \Rightarrow x=au'+bv'$$

Let  $x=au+bv$ .

$$g \mid a, g \mid b \Rightarrow g \mid x$$

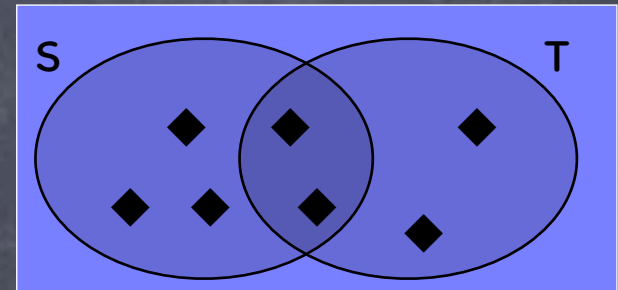


# Inclusion-Exclusion

- $|S| + |T|$  counts every element that is in  $S$  or in  $T$ 
  - But it double counts the number of elements that are in both:  
i.e., elements in  $S \cap T$

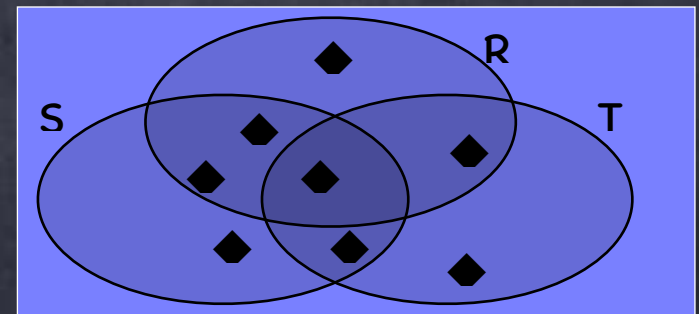
- So,  $|S| + |T| = |S \cup T| + |S \cap T|$

- Or,  $|S \cup T| = |S| + |T| - |S \cap T|$



- $|R \cup S \cup T| = |R| + |S| + |T| - |R \cap S| - |S \cap T| - |T \cap R| + |R \cap S \cap T|$

- $$\begin{aligned} |R \cup S \cup T| &= |R| + |S \cup T| - |R \cap (S \cup T)| \\ &= |R| + |S \cup T| - |(R \cap S) \cup (R \cap T)| \\ &= |R| + |S| + |T| - |S \cap T| \\ &\quad - (|R \cap S| + |R \cap T| - |R \cap S \cap T|) \end{aligned}$$

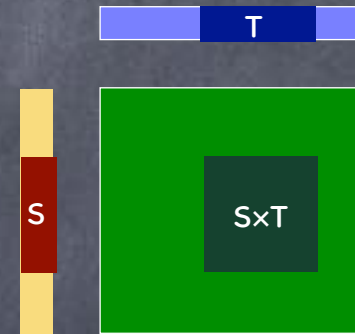


# Cartesian Product

- $S \times T = \{ (s,t) \mid s \in S \text{ and } t \in T \}$

- $(S = \emptyset \vee T = \emptyset) \leftrightarrow S \times T = \emptyset$

- $|S \times T| = |S| \cdot |T|$



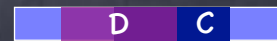
- $R \times S \times T = \{ (r,s,t) \mid r \in R, s \in S, t \in T \}$

$\{ ((r,s),t) \mid r \in R, s \in S, t \in T \}$

- Not the same as  $(R \times S) \times T$  (but "essentially" the same)

- $(A \cup B) \times C = (A \times C) \cup (B \times C)$ . Also,  $(A \cap B) \times C = (A \times C) \cap (B \times C)$

- $(A \cup B) \times (C \cup D) = (A \times (C \cup D)) \cup (B \times (C \cup D))$   
 $= (A \times C) \cup (A \times D) \cup (B \times C) \cup (B \times D)$



- Complement:  $\overline{S \times T} = ?$



- $(\overline{S} \times \overline{T}) \cup (\overline{S} \times T) \cup (S \times \overline{T})$