Countability and the Uncountable

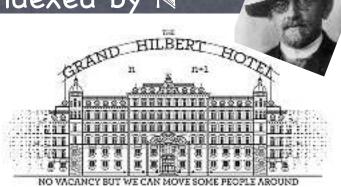


Hilbert's Hotel

The Grand Hilbert Hotel has infinite rooms indexed by ℕ

The natural numbers are all staying in this hotel, with number n occupying Room no. n

Suppose a new guest, -1 arrives



Can simply move everyone to the next room — n goes to Room no. n+1 — and make Room no. O available!

Suppose all the negative integers arrive

No problem! Move n to Room no. 2n, and make all the odd numbered rooms available. Can send -n to 2n-1 now.

But what if all the real numbers arrive?

The hotel can't accommodate them!

Image courtesy https://www.ias.edu/ideas/2016/pires-hilbert-hotel

How do you count infinity?

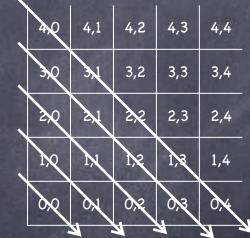
- How do you make precise the intuition that there are more real numbers than integers? Both are infinite...
- When do we say two infinite sets A & B have the same size?
- **Definition:** |A| = |B| if there is a bijection from A to B Definition
 - Ø |ℤ| = |2ℤ|. (2ℤ = evens). f : ℤ→2ℤ defined as f(x)=2x i finite sets too bijection

Countable

A set A is countably infinite if |A|=|N|
i.e., there is a bijection f : N → A
Note: |A|=|N| iff |A|=|Z|, |A|=|2Z| etc.
A set is countable if it is finite or countably infinite
Intuition: all "discrete" sets are countable

How do you count infinity? We defined: A is countably infinite if |A| = |N|, i.e., if there is a bijection between A and N.

(0,0), (1,0), (0,1), (2,0), (1,1), (0,2), ...
(i.e., f(0)=(0,0), f(1)=(1,0), f(2)=(0,1) ...)
Note: (0,0), (1,0), (2,0), (3,0) ... will not give a bijection



∑² is countable. f : ℤ²→ℕ defined as f(a,b) = h (g(a),g(b)), where
 g : ℤ→ℕ and h : ℕ²→ℕ are bijections, is a bijection

More generally, if A and B are countable, the A×B is countable (extended to any finite number of sets by induction)

But Things Get Messy...

- We saw bijection between ℤ² and ℕ. Enough to find a bijection between ℚ and ℤ².
- Not immediately clear: not all pairs (a,b) correspond to a distinct rational number a/b

a and b can have a common divisor; also, trouble with b=0
But easier to construct a <u>one-to-one</u> function f : Q→Z² as f(x) = (p,q) where x=p/q is the "canonical representation" of x (i.e., gcd(p,q)=1 and q > 0).

 ${}_{\textcircled{O}}$ Hence one-to-one function gof : ${}_{\textcircled{O}} \rightarrow {}_{\swarrow}$, where g : ${}_{\textcircled{Z}} \rightarrow {}_{\bigtriangledown}$ is a bijection

But Things Get Messy...

👩 Is 🔍 countable?

 \odot One-to-one function $f: \mathbb{Q} \rightarrow \mathbb{N}$

Intuitively, if a one-to-one function from A to B, |A|≤|B|

True for finite sets

Definition good for finite sets too

• **Definition:** $|A| \leq |B|$ if there is a one-to-one function from A to B

Also can construct a one-to-one function h : N→Q as h(a)=a.
 So |N|≤|Q|

Bijection from Two Injections

Cantor-Schröder-Bernstein

- Theorem [CSB]: There is a bijection from A to B if and only if there is a one-to-one function from A to B, and a one-to-one function from B to A
- Restated: |A|=|B| ⇔ |A| ≤ |B| and |B| ≤ |A| <
 </p>

Trivial for finite sets

- Proof idea: Let $f: A \rightarrow B$ and $g: B \rightarrow A$ (one-to-one).
- Consider infinite chains obtained by following the arrows
 - \odot One-to-one \Rightarrow Each node in a unique chain
 - Chain could start from an A node, start from a B node or has no starting node (doubly infinite or cyclic). Say, types A,B and C
 - Let h : A→B s.t. h(a)=f(a) if a's chain type A;
 else h(a)=g⁻¹(a).

Find a perfect matching

Bijection from Two Injections

Since |Q| ≤ |N| and |N| ≤ |Q|, by CSB-theorem |Q| = |N|

👩 🔍 is countable

Example: The set S of all finite-length strings made of [A-Z] is countably infinite

Interpret A to Z as the <u>non-zero</u> digits in base 27. Given s∈S, interpret it as a number. This mapping (S→N) is one-to-one (because no leading zeroes).

Map an integer n to Aⁿ (string with n As). This is one-to-one.

Bijection from Two Injections

Another example: T be the set of all infinitely long binary strings
 Claim: |T| = |R|
 Not true if we used binary representation instead of decimal representation, we'll have strings 011111.. and

This is a one-to-one mapping: a finite difference between the real numbers that two different strings map to

10000... map to the same real number

IR ≤ T: map each x ∈ R injectively to a real number in (0,1), say as (1+e^{-x})⁻¹ and then injectively to the string in its binary expansion

 \odot Corollary: $|\mathbb{R}^2| = |\mathbb{R}|$

Because $|T^2| = |T|$ (bijection by interleaving), and $|\mathbb{R}^2| = |T^2|$

Summary

Equivalently: there is an onto function from B to A (relying on the "Axiom of Choice")

Definition: |A| = |B| if there is a bijection from A to B • Definition: $|A| \leq |B|$ if there is a one-to-one function from A to B • Theorem [CSB]: $|A|=|B| \Leftrightarrow |A| \leq |B|$ and $|B| \leq |A|$ \odot A is countably infinite if |A| = |N| \circ e.q., $|\mathbf{Q}| = |\mathbf{N}|$ (saw one-to-one functions in both directions) A is <u>uncountable</u> if A is infinite but not countably infinite

 \odot Equivalently, if no function $f : A \rightarrow \mathbb{N}$ is one-to-one

 \odot Equivalently, if no function $f : \mathbb{N} \rightarrow A$ is onto

Uncountable Sets

- Related claims:
 - \odot Set T of all infinitely long binary strings is uncountable
 - Contrast with set of all finitely long binary strings, which is a countably infinite set
 - The power-set of \mathbb{N} , $\mathbb{P}(\mathbb{N})$ is uncountable

There is a bijection f: T → P(N) defined as $f(s) = \{i \mid s_i = 1\}$

How do we show something is not countable?!
Cantor's "diagonal slash"

e.g., set of even numbers corresponds to the string 101010...

Cantor's Diagonal Slash

- ${f o}$ Take any function ${f f}:{f N}{
 ightarrow}{f P}({f N})$
- Make a binary table with $T_{ij} = 1$ iff j∈f(i)
- Consider the set X ⊆ N
 corresponding to the "flipped diagonal"

X doesn't appear as a row in this table (why?)
 So f not onto

	0	1	0	0	1	1	1
f(0) -			0	1	$\mathbf{\hat{\mathbf{A}}}$	\bigcirc	0
f(0) =	1	0	0	1	0	0	0
f(1) =	0	0	1	0	1	0	0
f(2) =	1	1	1	1	1	1	1
f(3) =	1	1	0	1	0	1	0
f(4) =	1	1	0	0	0	0	1
f(5) =	0	1	0	1	1	0	1
f(6) =	0	1	0	1	0	1	0

Cantor's Diagonal Slash

- Take any function $f: \mathbb{N} \rightarrow \mathbb{P}(\mathbb{N})$
- Make a binary table with $T_{ij} = 1$ iff $j \in f(i)$
- Consider the set X ⊆ N
 corresponding to the "flipped diagonal"

X doesn't appear as a row in this
 table (why?)
 So f not onto

 $\frac{\text{Generalizes:}}{\text{No onto function } f: A \rightarrow \mathbb{P}(A)}$

May not have a <u>table</u> enumerating f (if A is uncountable)

for any set A

Let $X = \{ j \in A \mid j \notin f(j) \}$

Claim: $\nexists i \in A$ s.t. X = f(i)

Suppose not: i.e., $\exists i, X=f(i)$. $i \in X \leftrightarrow i \in f(i) \leftrightarrow i \notin X$ Contradiction!

Reals are Uncountable

We saw P(N) is uncountable, i.e., |P(N)| ≠ |N|
Recall T, the set of infinite binary strings
We saw |T| = |P(N)| (via an easy bijection) and |T| = |R| (via CSB)
Hence |R| ≠ |N|, i.e., R uncountable
|R| is a "higher infinity" than |N|
Denoted as N₀, N₁, N₂,...
|P(R)| even higher