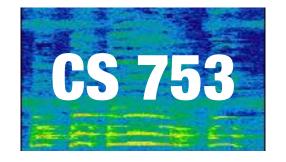
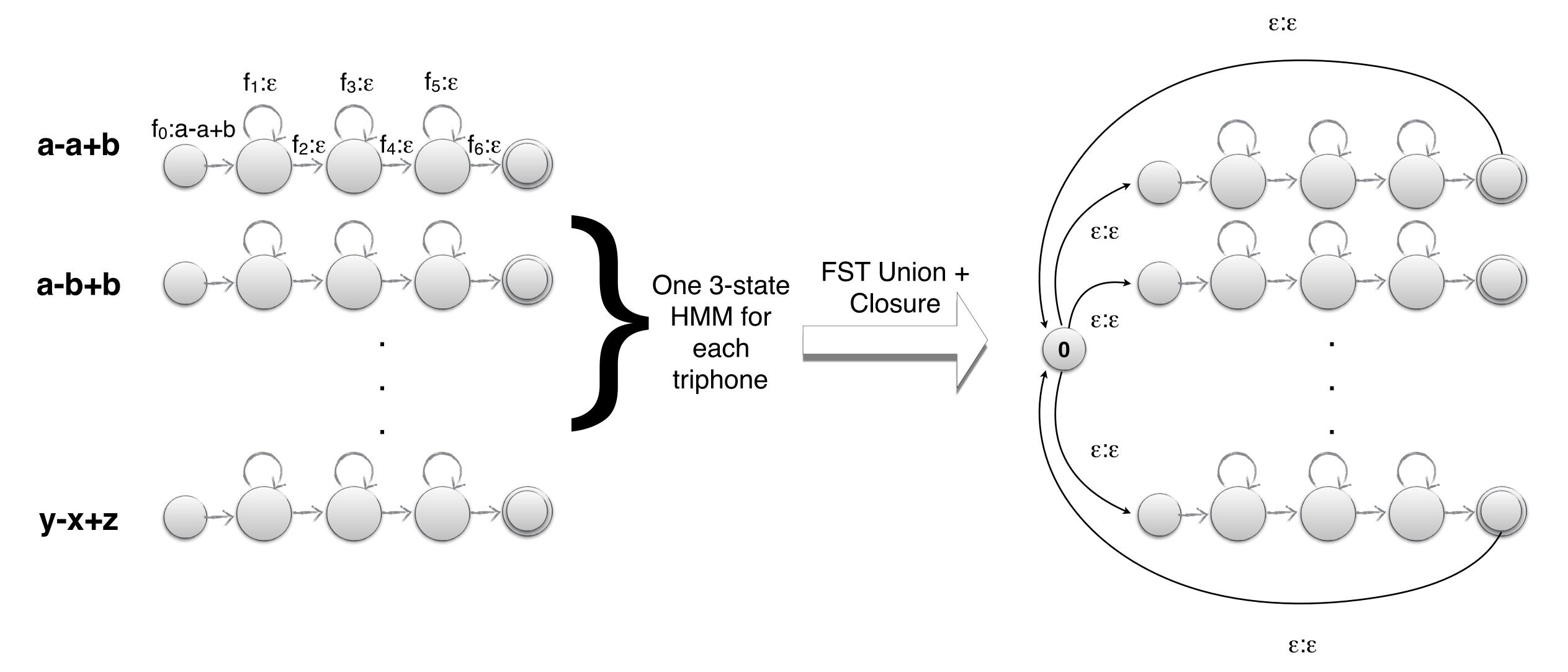
WFSTs continued

Lecture 5

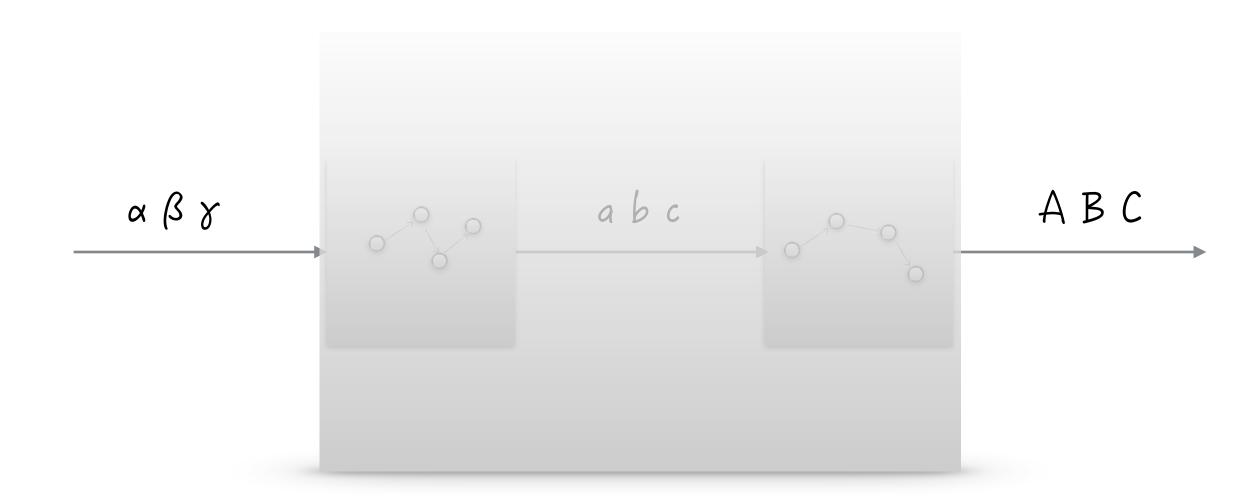


Instructor: Preethi Jyothi

Recap: Construct acoustic Model WFST using simple rational operations



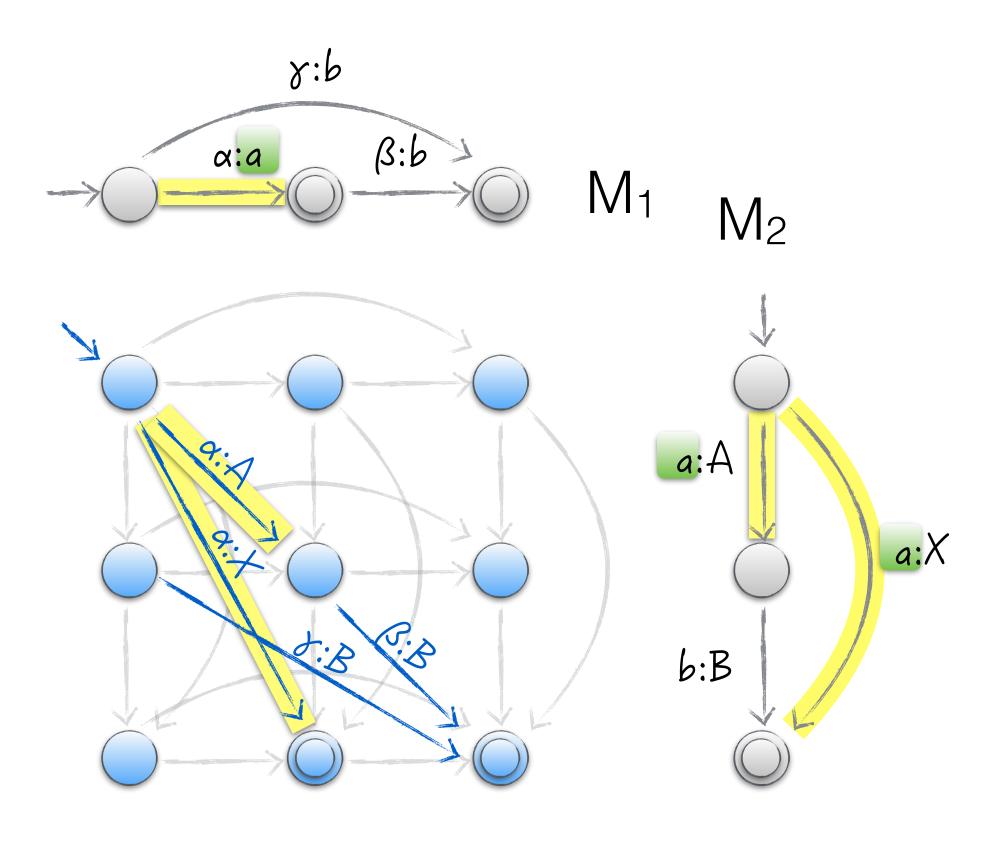
Recap: Composition



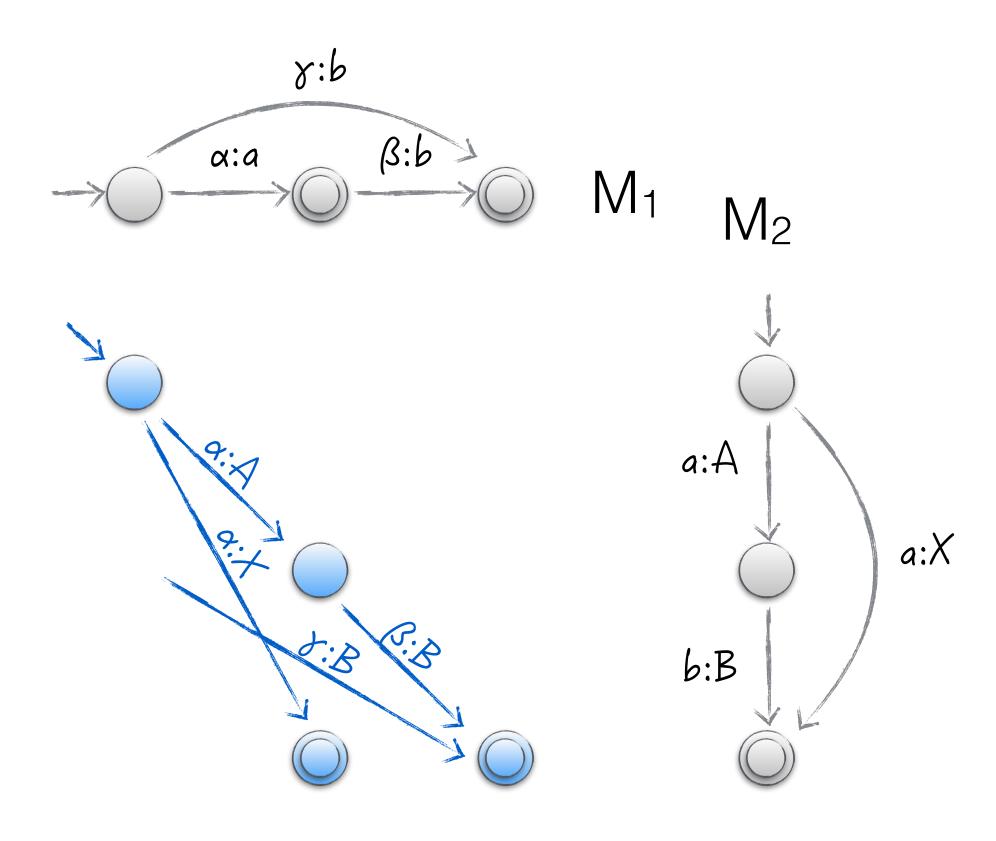
• If T_1 transduces x to z, and T_2 transduces z to y, then $T_1 \circ T_2$ transduces x to y

•
$$(T_1 \circ T_2)(x, y) = \bigoplus_{z} T_1(x, z) \otimes T_2(z, y)$$

Composition: Construction

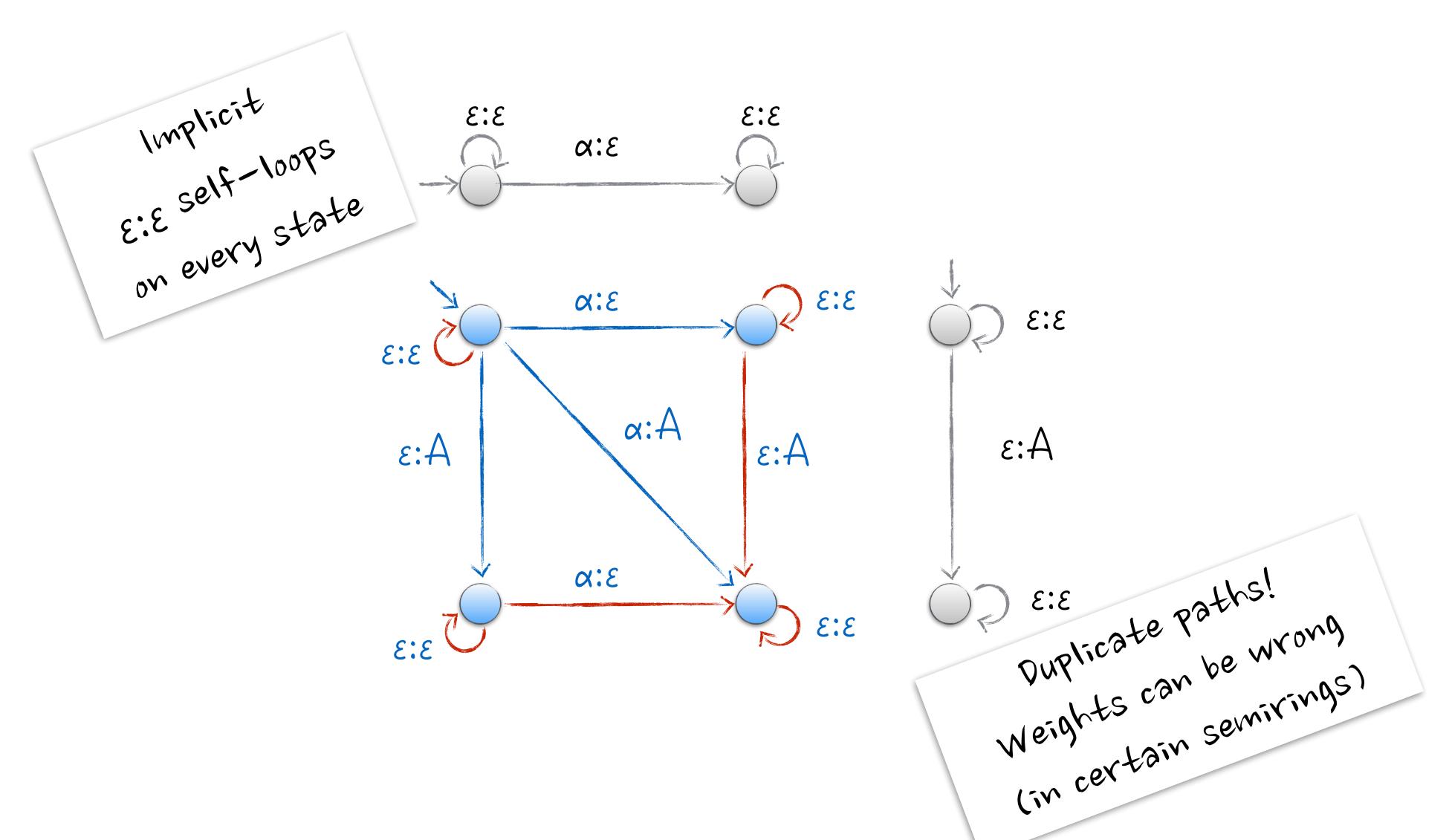


Composition

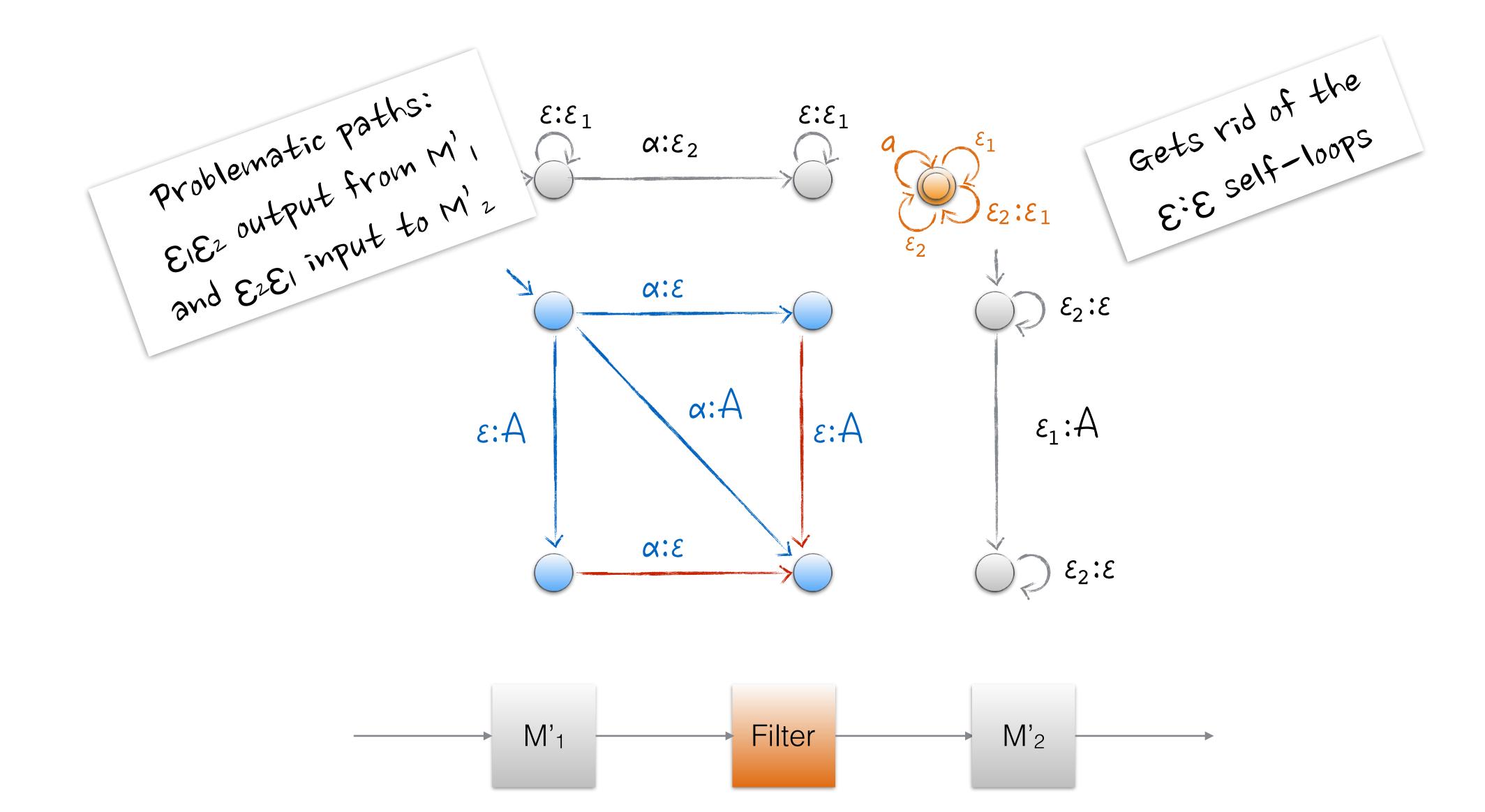


 $M_1 \circ M_2$

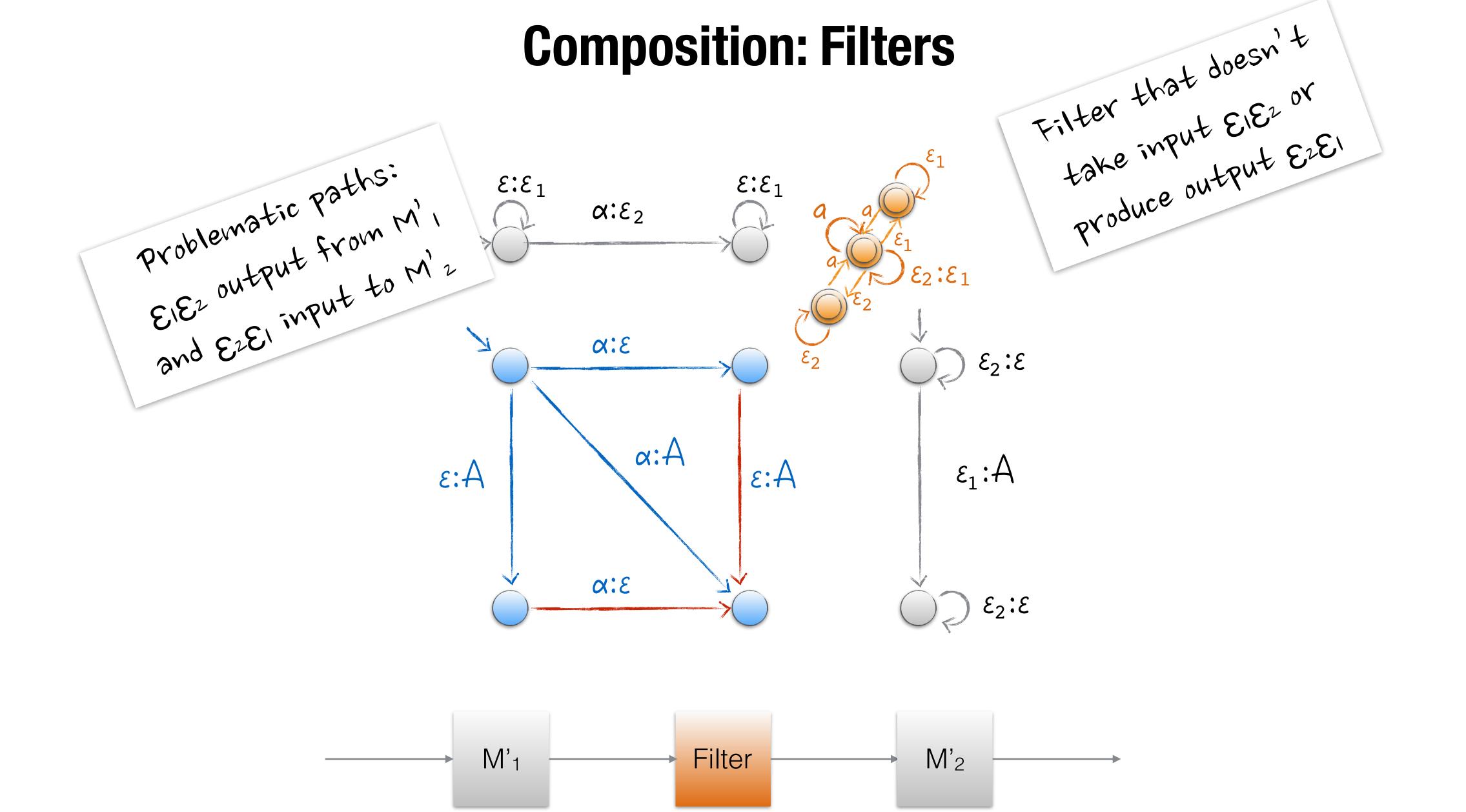
Composition: Handling epsilons

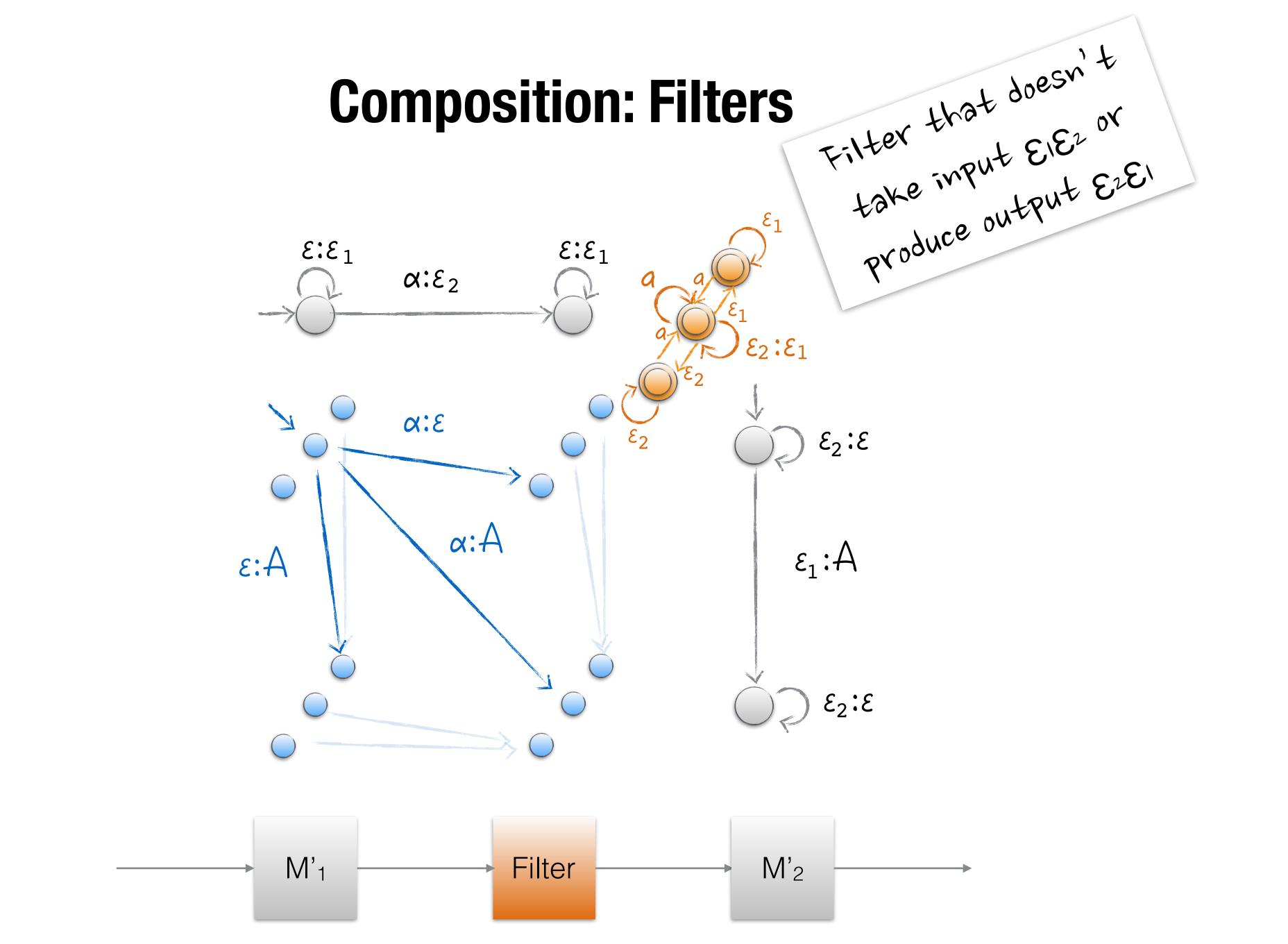


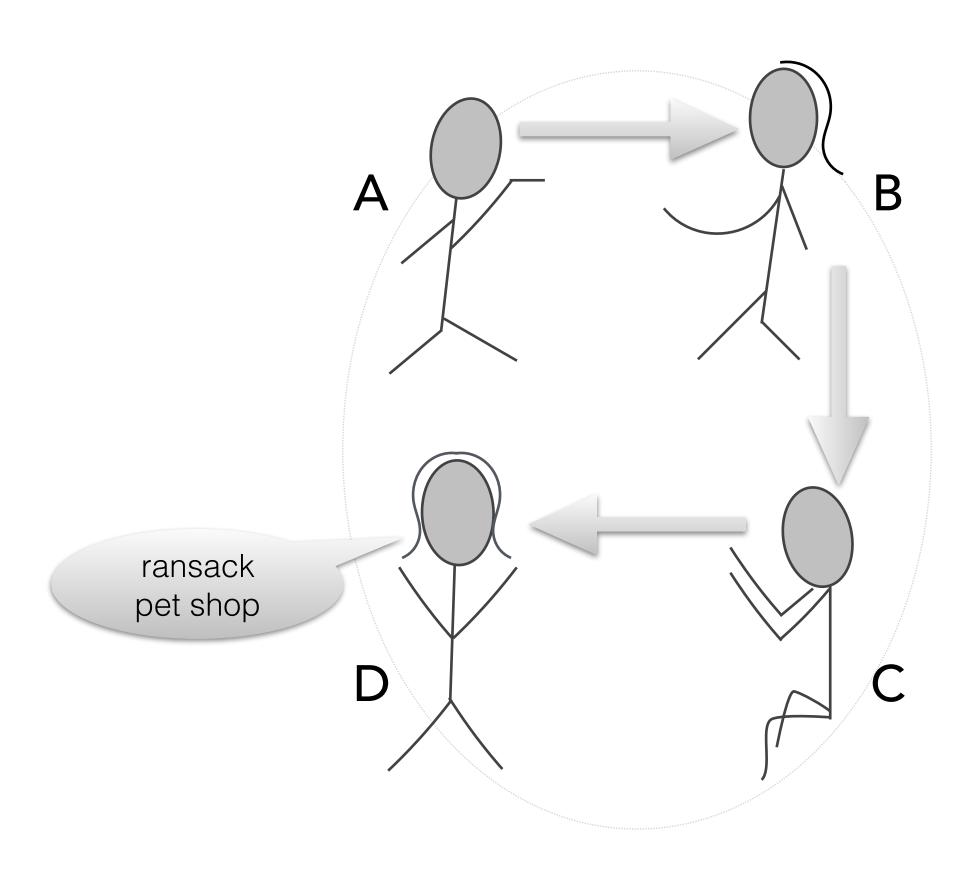
Composition: Filters



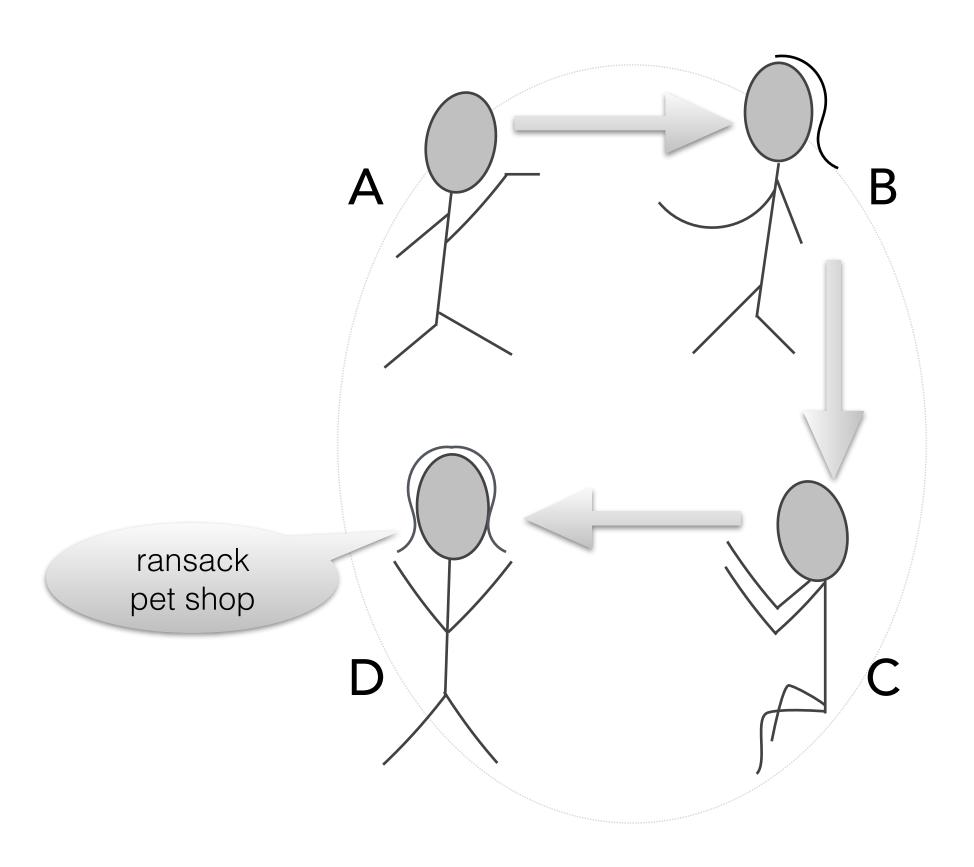
Composition: Filters



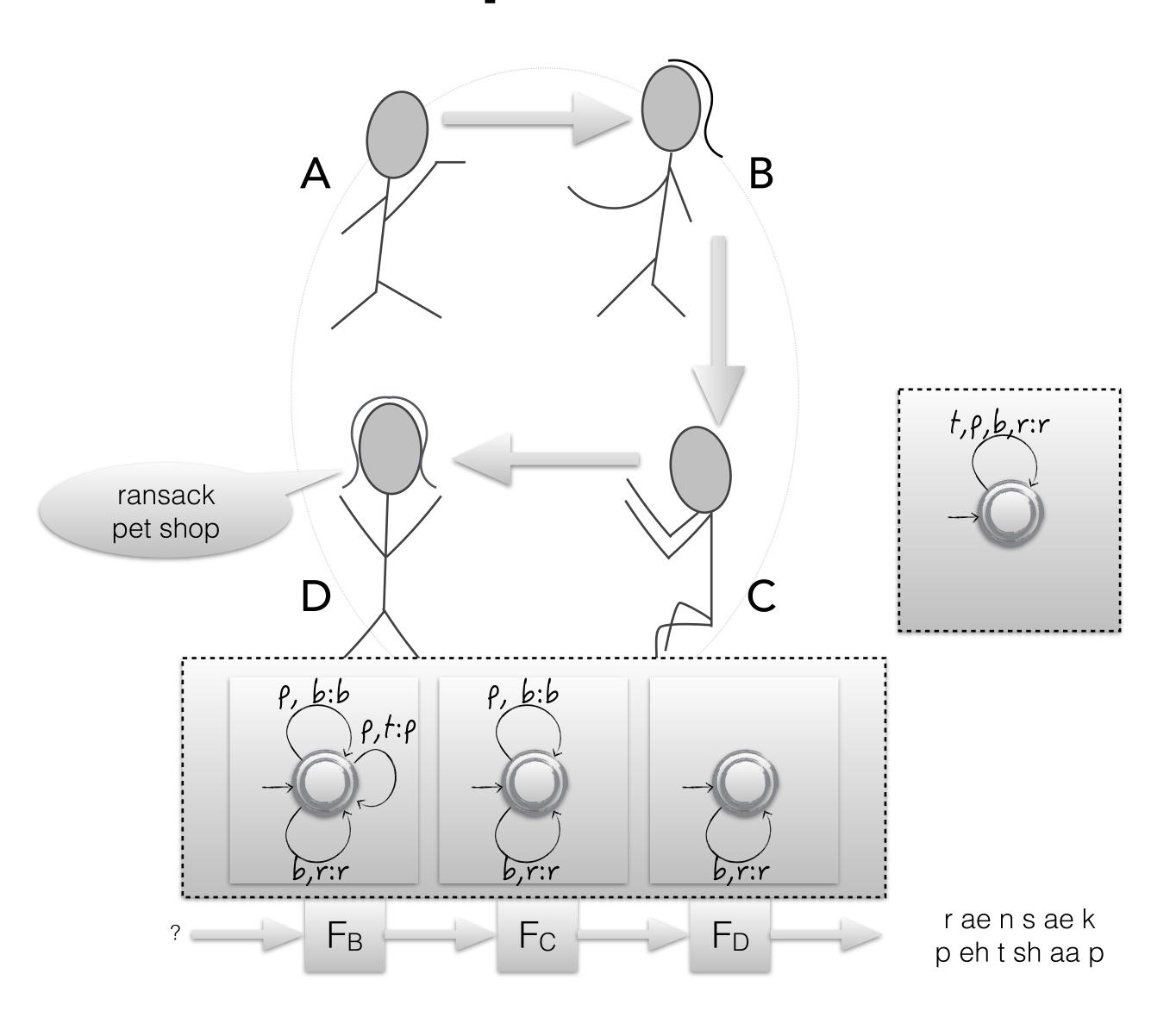


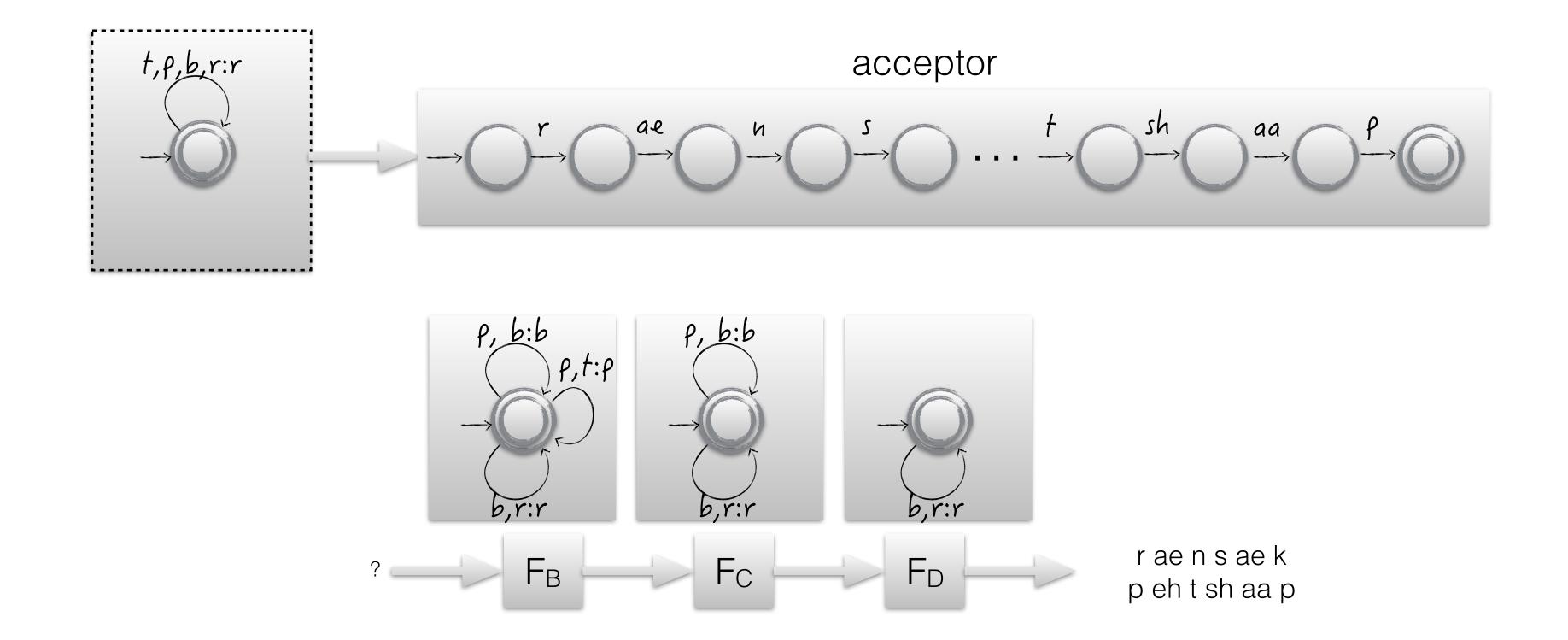


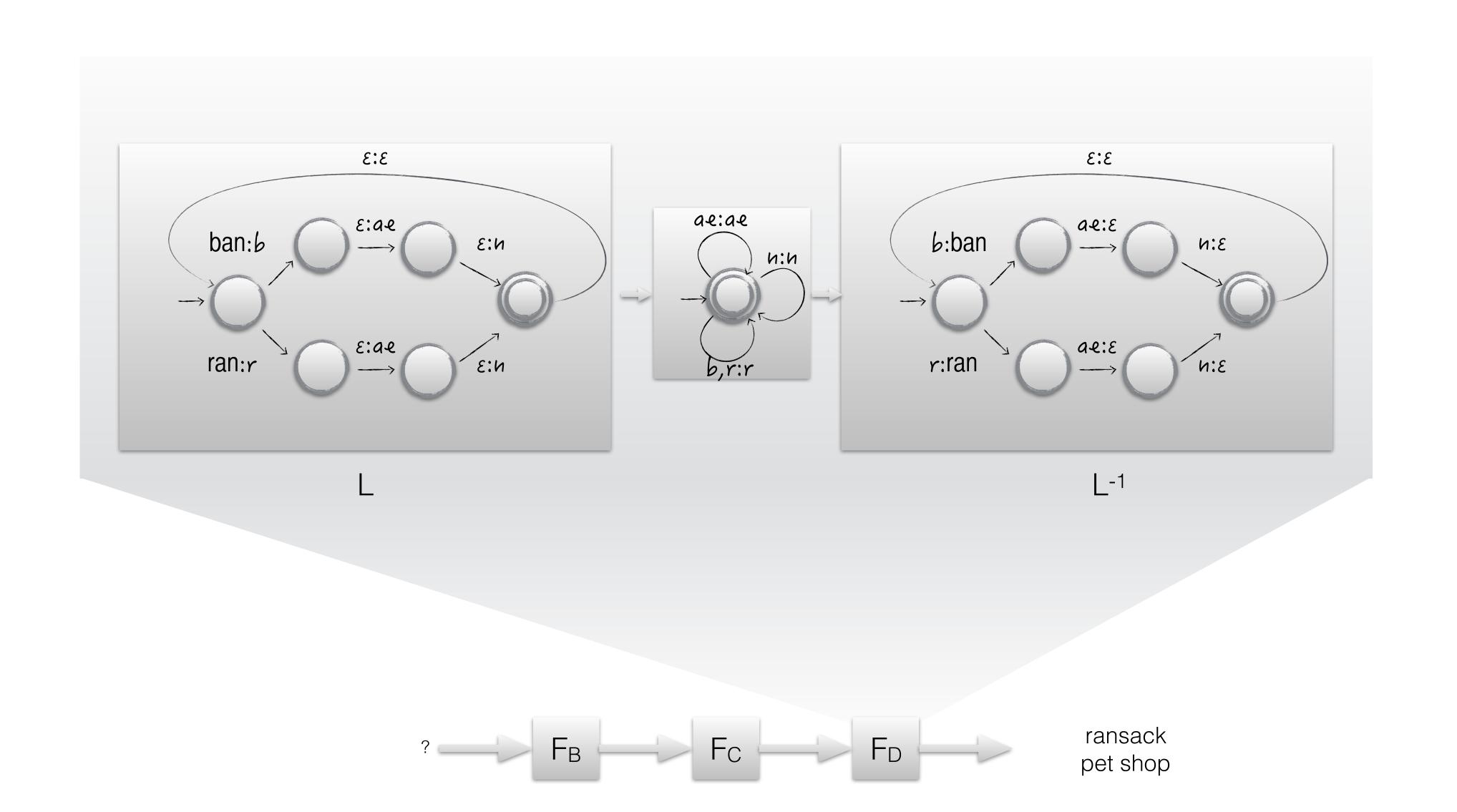
Given what D said, can we infer the message A started with?

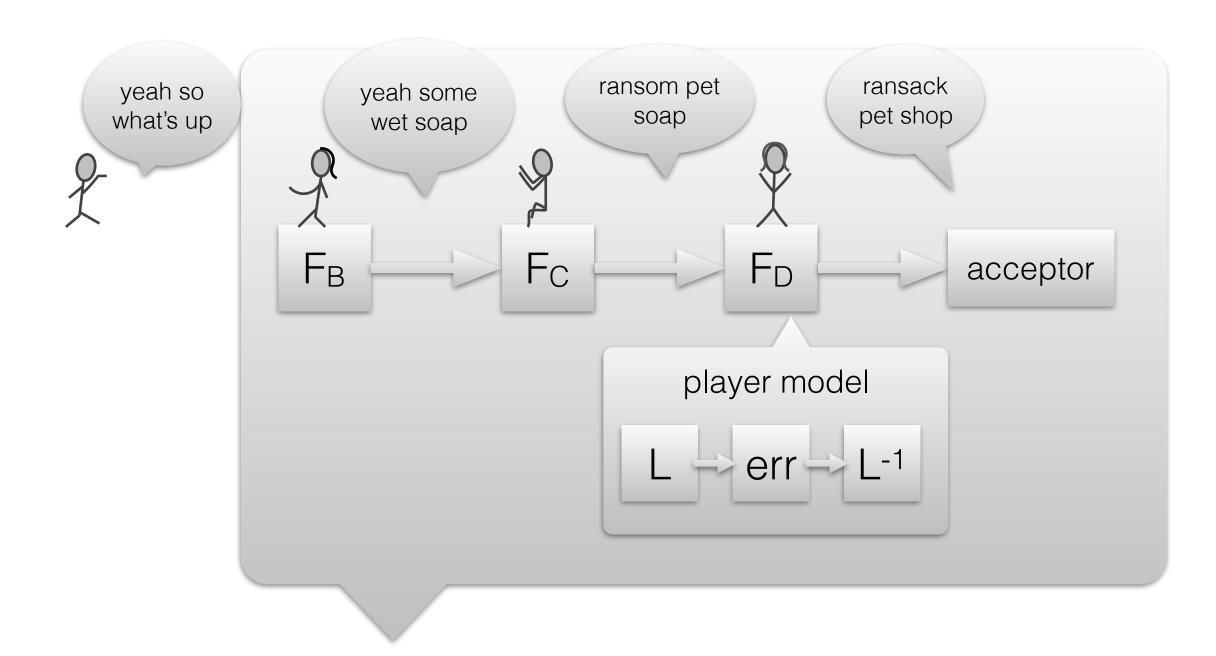


Model the errors made by each player using an FST









Find the best path in this FST

Read off the input words on the arcs

Can also find the best combination of paths in each player FST

Composition: Recap

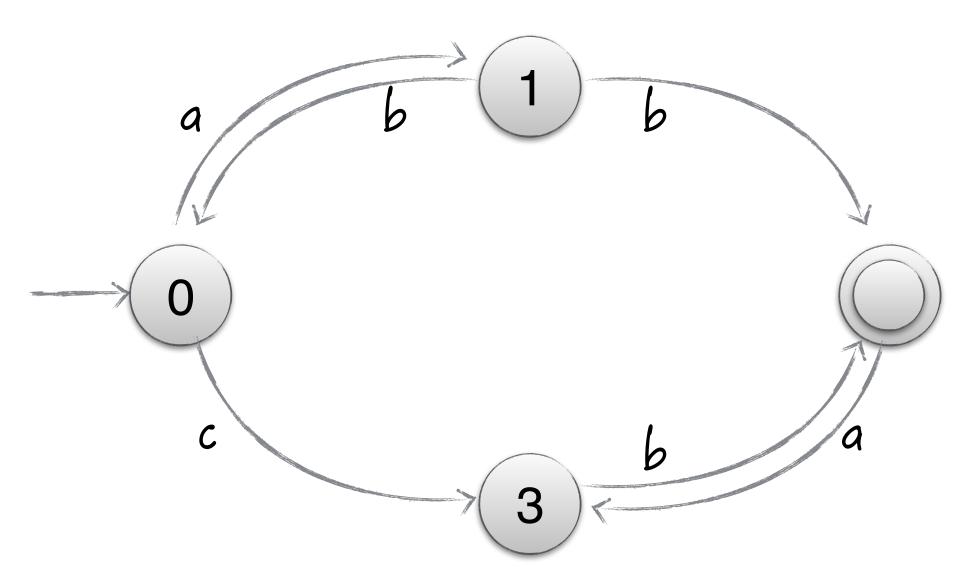
- If T_1 transduces x to z, and T_2 transduces z to y, then $T_1 \circ T_2$ transduces x to y
- Note: output alphabet of T₁ ⊆ input alphabet of T₂
- E.g. If T_1 removes punctuation symbols from a string, and T_2 changes uppercase letters to lowercase letters, then $T_1 \circ T_2$ brings about both changes

Determinization and Minimization

- WFSTs constructed using various operations (or designed by hand) may have several redundancies
 - Affects the efficiency of subsequent operations
- Determinization and minimization seek to remove redundancies
 - Determinization can expand a WFST, but makes it faster to process an input string
 - Minimization results in the smallest number of states
- Will discuss WFSAs here. Extends to WFSTs.

Deterministic FSAs

- An FSA is deterministic if:
 - Unique start state
 - No two transitions from a state share the same label
 - No epsilon labels



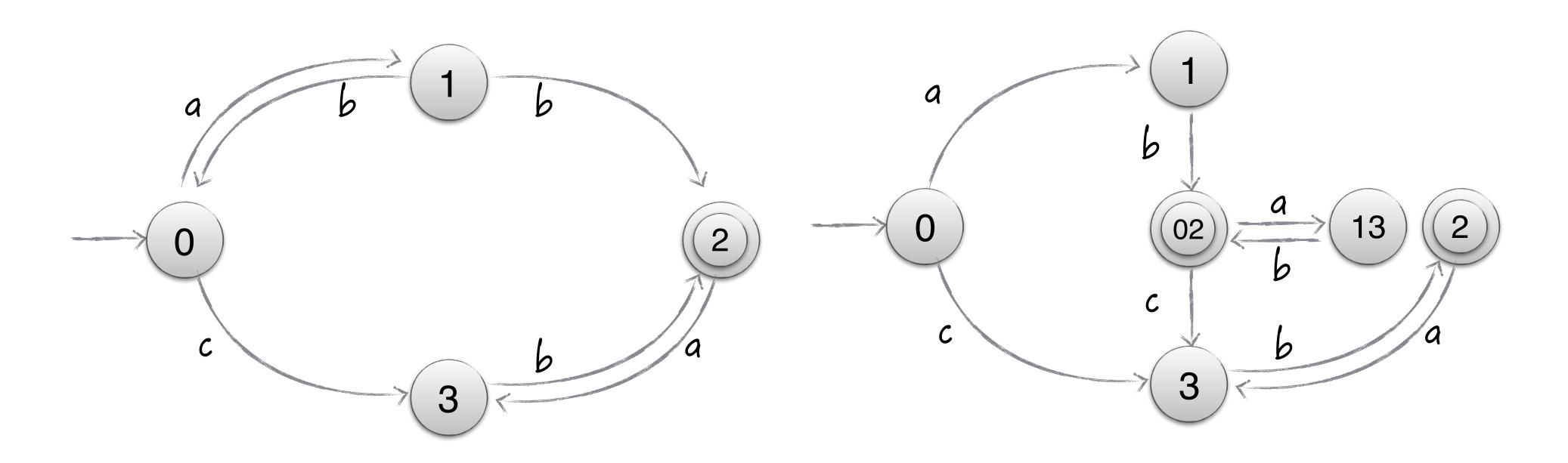
Deterministic or non-deterministic?

Any input sequence yields a unique path (if at all)

Determinization

Construct an equivalent deterministic FSA

States correspond to subsets of states in the original FSA

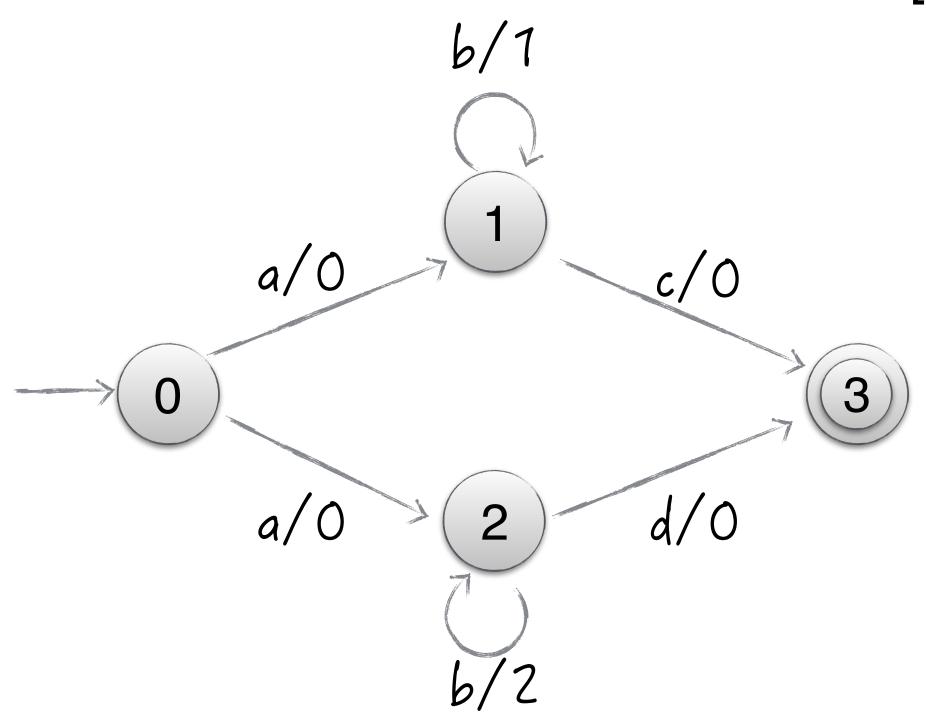


non-deterministic FSA

equivalent deterministic FSA

Determinization: Weighted FSA

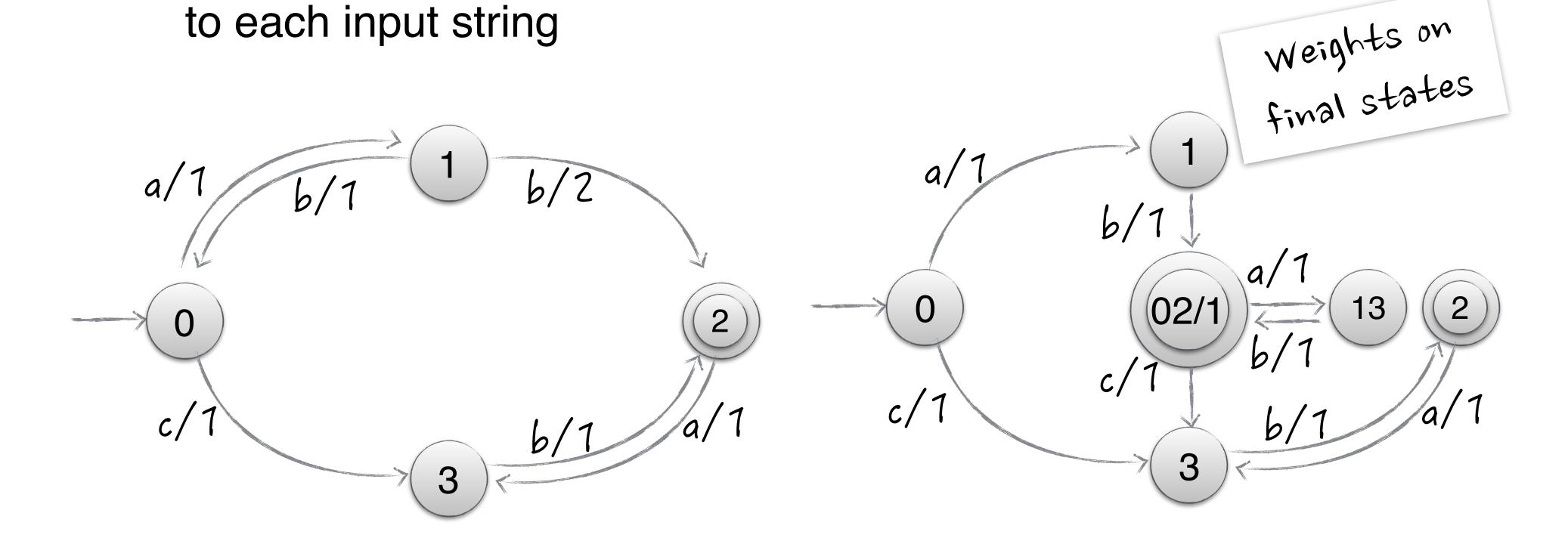
Some Weighted-FSAs are not determinizable! [M97]



Weight of string $ab^nc = n$ and weight of $ab^nd = 2n$ After seeing ab^n an FSA can't remember n

Determinization: Weighted FSA

Two WFSAs are equivalent if they associate the same weight to each input string



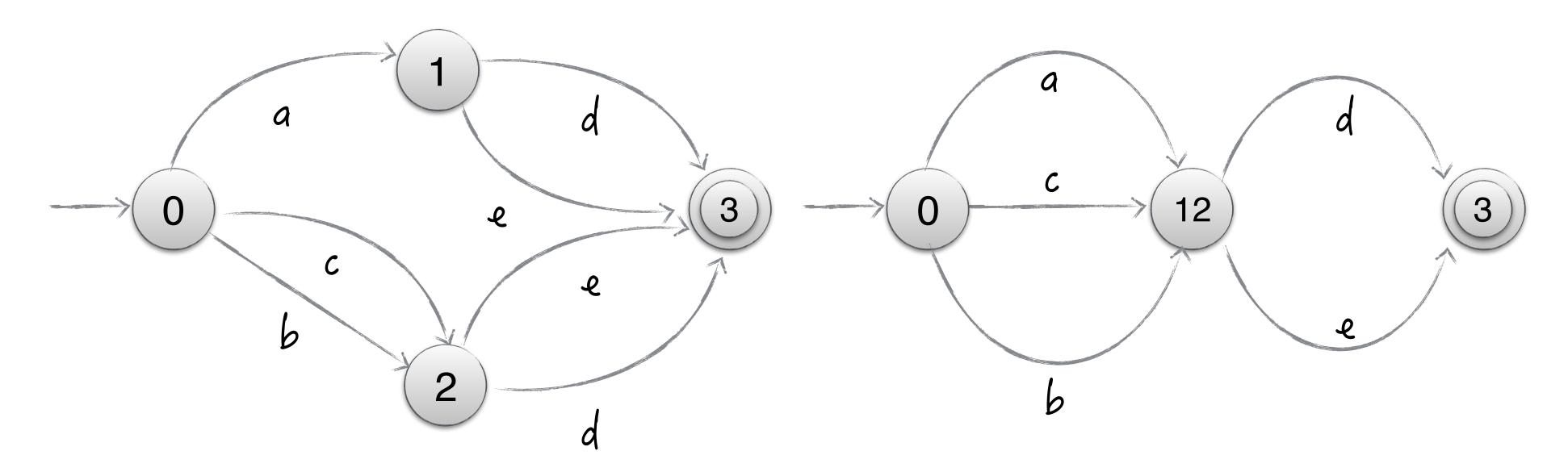
non-deterministic WFSA

equivalent deterministic WFSA

Minimization

Minimization: find an equivalent deterministic FSA with the least number of states (and transitions)

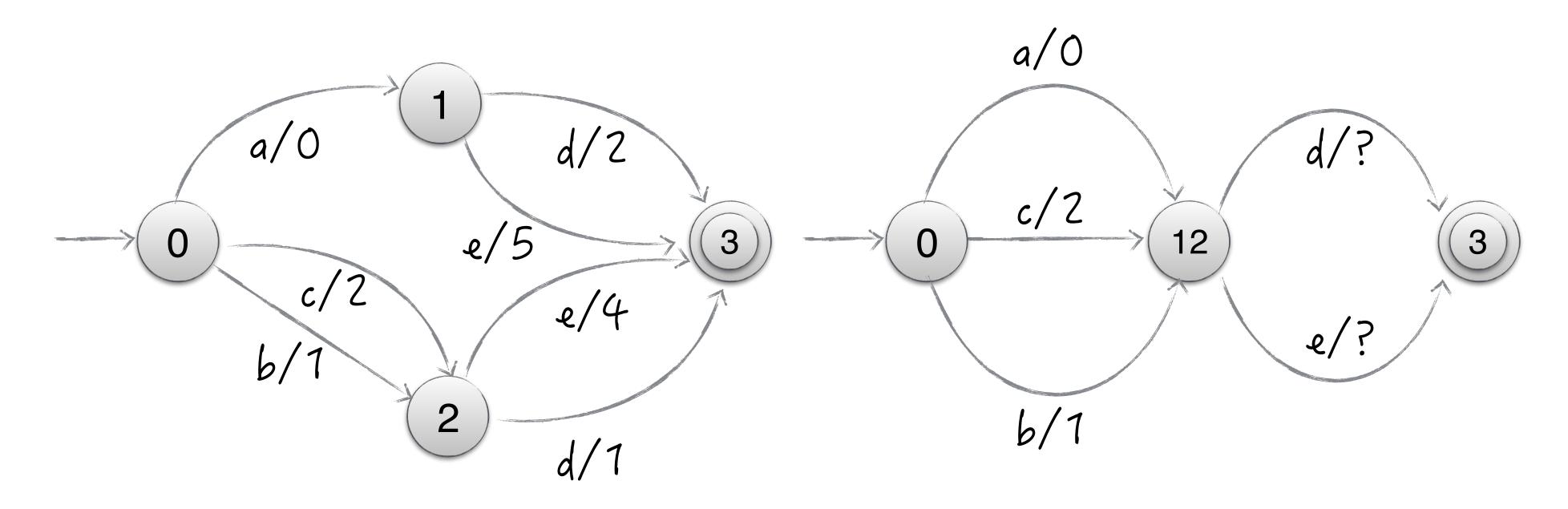
Unweighted FSAs have a unique minimal FSA [Aho74]



Obtained by identifying and merging equivalent states

Minimization: Weighted FSA

Two states are equivalent only if for every input string, the outcome — weight assigned to the string, if accepted — starting from the two states are the same

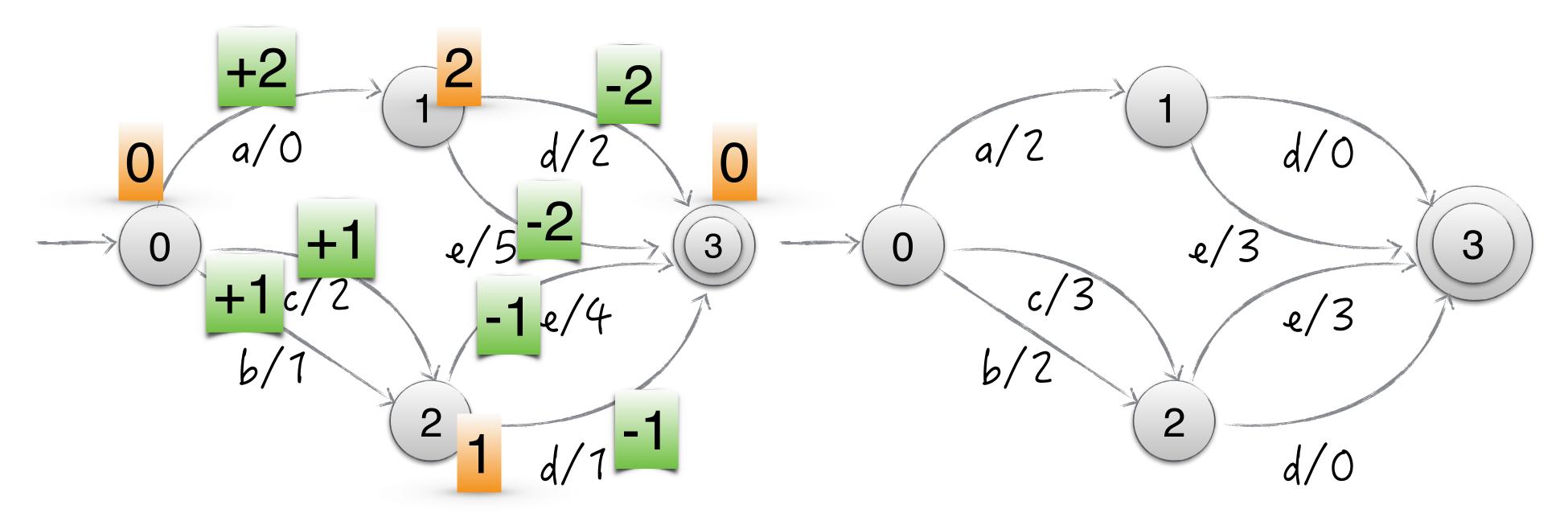


Redistribute weights before identifying equivalent states

Minimization: Weighted FSA

Reweighting OK as long as resulting WFSA is equivalent

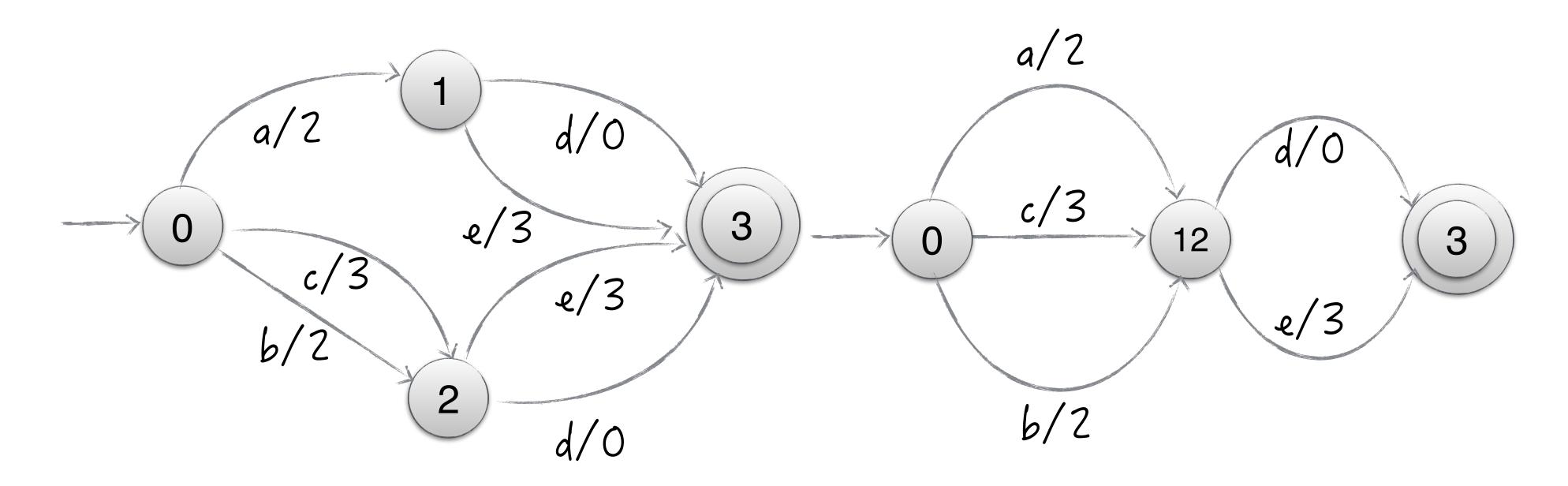
Can reweight using a "potential function" on states



"Weight pushing": Reweighting using a potential function that optimally moves weights towards the start state

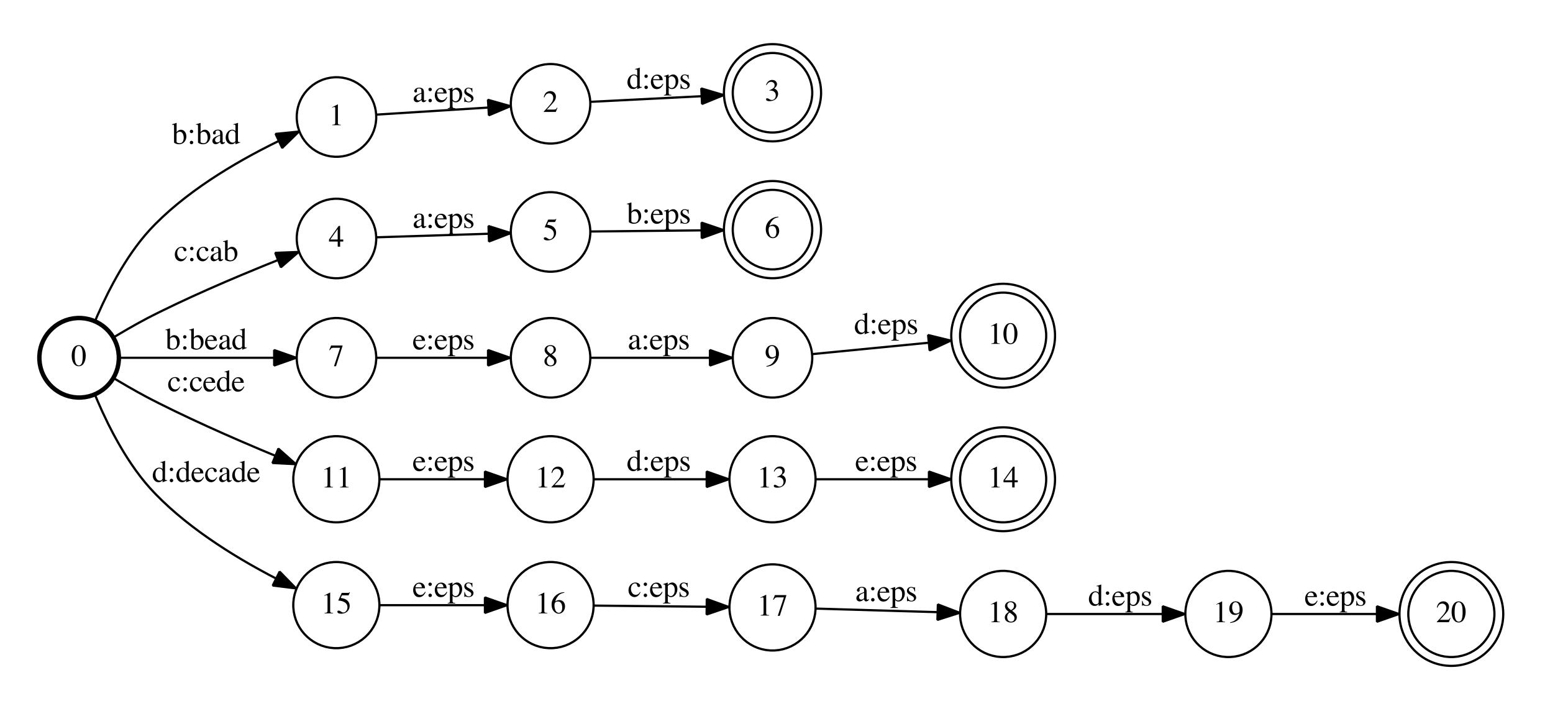
Minimization: Weighted FSA

After weight-pushing, can simply apply unweighted FSA minimization (treating label/weight as label)

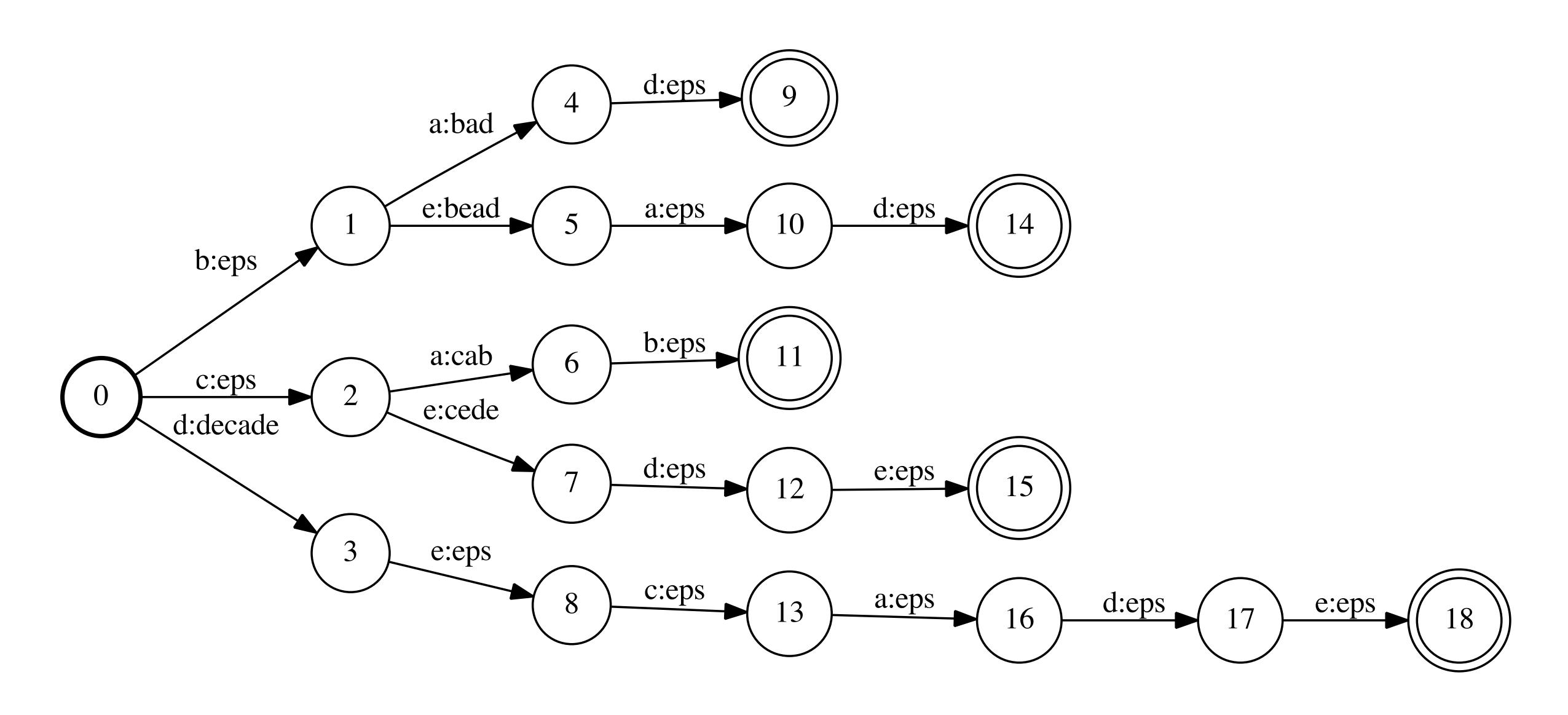


Guaranteed to yield a minimal WFSA (under some technical conditions required for weight-pushing)

Example: Dictionary WFST



Determinized Dictionary WFST



Minimized Dictionary WFST

