

## Recap:

- OS provides several services
  - process creation / deletion
  - memory allocation / separation
  - safe device access

via different abstractions

- files
- processes
- address space
- Network interfaces

- two main tasks performed by an OS
  - resource mgmt. / multiplexing
  - isolation & control.

- Key requirement to do above tasks is that the OS owns all resources.

⇒ all access to resources need OS mediation.

## Types of OSes.

- monolithic vs. microkernel vs. Unikernel
- server vs. desktop vs. handhelds vs. embedded

## \* two key building blocks for OSes

- (privileged) modes of execution
- interrupts / interrupt-driven ops.

to exercise control on (a) resources & their mgmt.

(to prevent bad things)

- memory isolation
- safe / fair / no-starvation scheduling

- (b) idea: all tasks related to resource control / mgmt. to be handled by OS, rest user programs can directly execute on CPU.
- scheduling, allocation etc.

user-mode — less privileges

kernel-mode — for OS actions

(c) how are the modes put in action?

- requires h/w support (in CPU)
- every instruction requires a minimum privilege level to execute.
- h/w regs. stores CPL, which CPU implicitly checks.

- (d) e.g.: to access memory of a process certain h/w regs. need to be setup.

| 32-bit virtual (x86)

— a process certain hw reg. need to be setup.

— e.g. regs. CR3 on x86 systems.

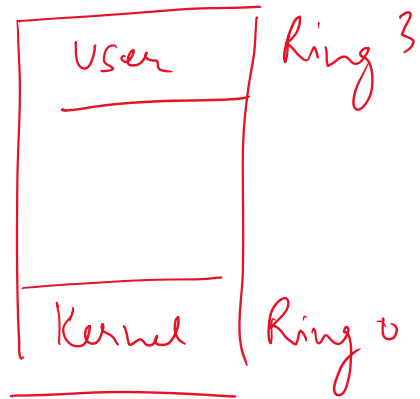
update of CR3 requires Kernel mode

writes from user mode Ring 0

have no effect or result in error.

— write to interval timer  
(same idea).

logical view (x86)



(ii) interrupts

— an interrupt, interrupts current execution on CPU so that other work can be done.

— interrupts "handled" by interrupt handlers, interrupt service routines (ISRs) / functions.

two building blocks  
(i) (privileged) modes of execution

(ii) interrupts

# Why interrupts important?

(i) the world is non-deterministic. "anything can happen any time"

— pkt. arrivals, key presses, disk reads.

— how does a CPU / OS know when to process these events.

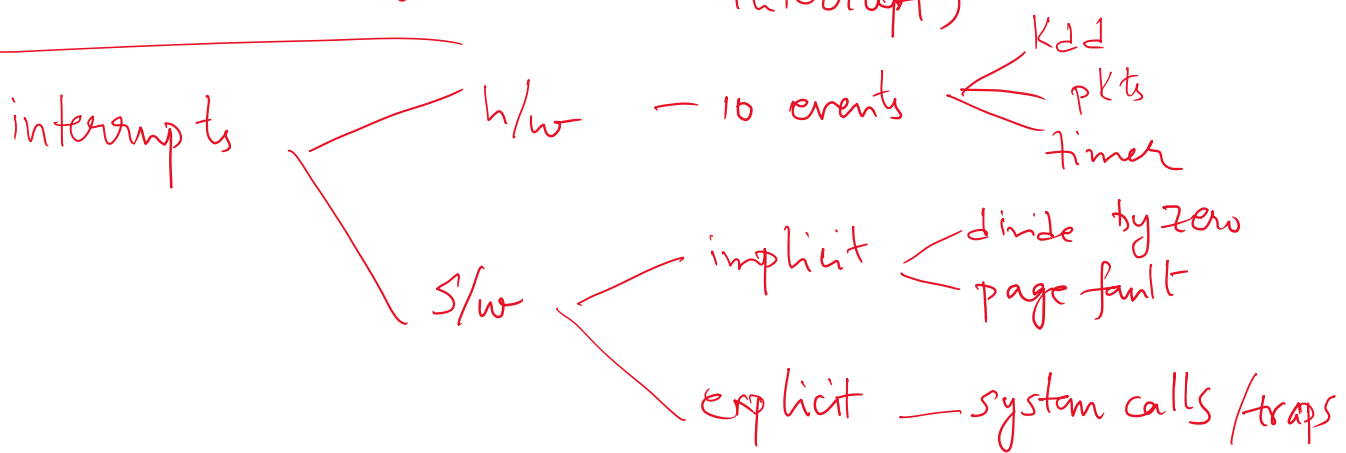
— interrupts, polling.

(ii) basis mechanism to invoke OS (for services).

— use of interrupt to switch to kernel mode for OS service

— implicit ————— access unallocated memory (page fault)

vs.  
explicit — system call (issues a software interrupt)



Q.

— What is special about the timer interrupt?

- how does CPU know which ISR to execute?
- constraints / conditions on ISR code?

Lab 2 Heads up. Processes.

Process creation, mgmt from user space.

- fork & c