

lecture 43

05 November 2018 09:39

① in-memory FS state:

- inode cache
- dentry cache
- page cache
- FS logic + other state

on-disk

- super block
- metadata blocks (inode)
- data blocks.

② dentry cache ~ directory entries.

/home /data /file.txt } logical pathname

no explicit
(data) description
of hierarchy
setup.

dentry cache — < dentry, inode >

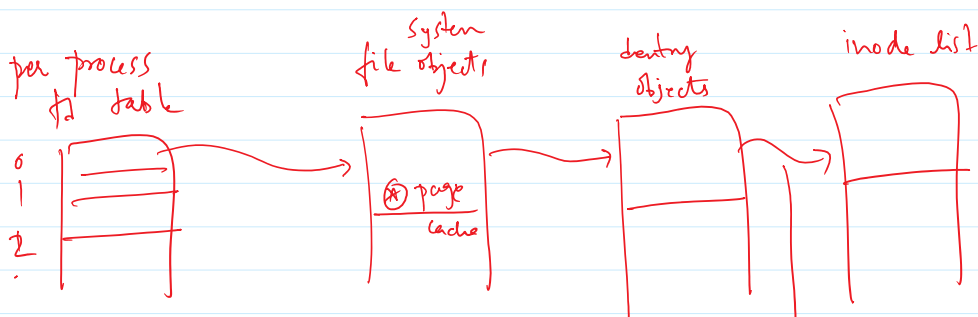
① lookup path component in dentry cache

- if hit
get/use inode

- else
create dentry and lookup inode from disk.

② - use inode to read contents (file) of directory.

- check for next path component
 - name, permissions etc.
 - goto ①
 - else, error!



→ /home /data1 /file1.txt
 → /home /data2 /file2.txt
 ? ?

< / → inode1 >

< /home → inode2 >

↑
path component

directory - is a file

directory - is a file

data blocks of this file : $\langle \text{filename, inode \#} \rangle$

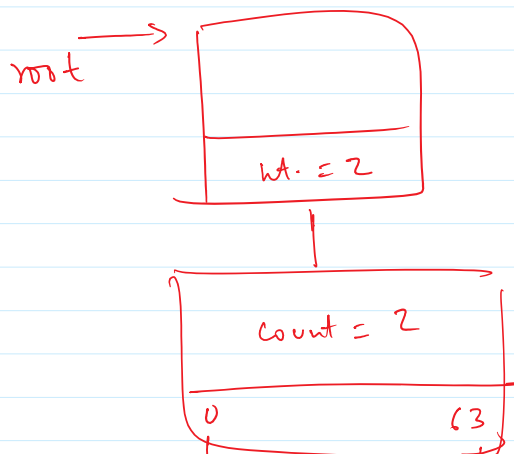
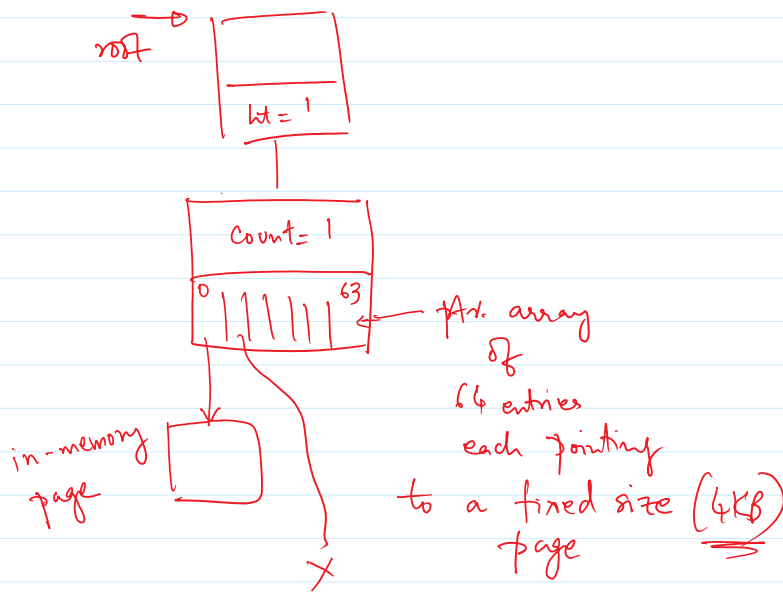
(3) page cache : 'data' cache.

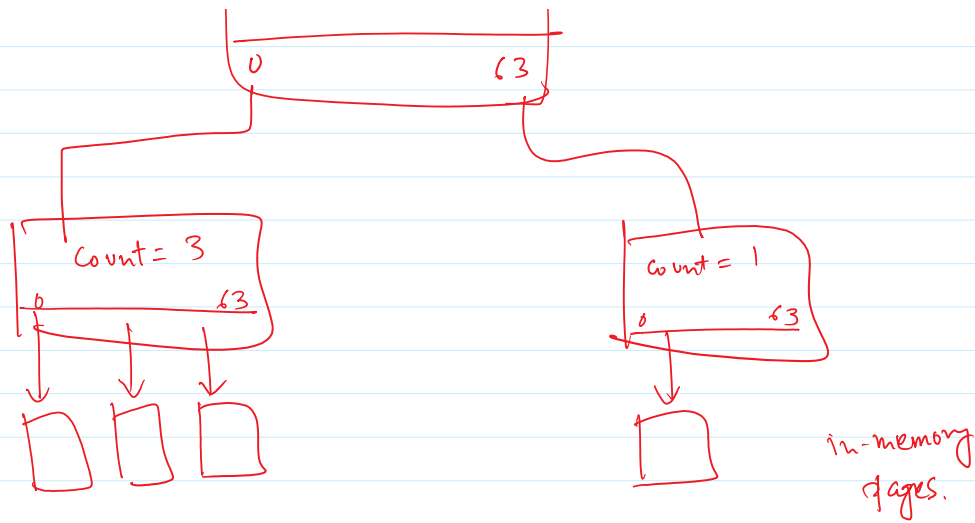
~ file/disk
page cache — $\langle \text{file, offset} \rangle$ ②

buffer cache — $\langle \text{disk, block} \# \rangle$

~ challenge ~ how to get to a cached object quickly?
& store/build cache efficiently.

linux radix tree base page cache.





② all reads/writes go through page cache
 ↳ one-mode of operation.

— cache consisting of ~~for~~ techniques

— write-through — sync. writes to disk

— write-back — async. writes to disk

↳ periodic flush to disk (Kernel threads).

③ file system consistency. ~ JFS journaling log file system.

↳ meta-data not in sync with data

— super-block not in sync with meta data.

e.g: tool 'fsck' file system check