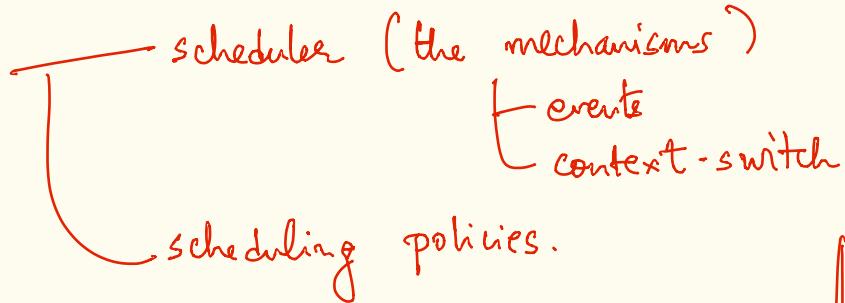


Lecture 12

CS 347/333

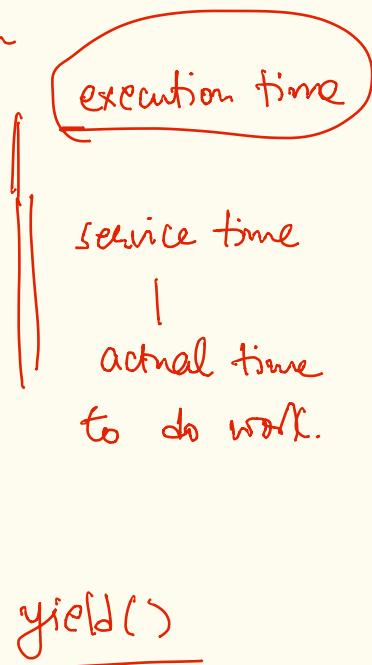
30.8.2022

(*) scheduling



(*) (i) metrics:

utilization	fairness (no starvation)
finish time	waiting time
response time	throughput
# deadlines met.	



(ii) categories:

(i) non-preemptive vs. preemptive

Schedule only when current process yields, exits or blocks.

Schedule a new process even if current process is not blocked/done (basically wants to continue running).

(ii) non-work conserving

if resource is available & work to be done, may not schedule work on resource.

vs work conserving.

if resource is available & work is to be done, schedule work, no resource idling.



(iii) policy criteria / challenges.

- interactive / bursty vs. CPU-intensive
- real time vs. best effort

Examples:

- ① FCFS: first-come-first-served
- add to end of ready queue
 - dequeue from head of queue. (on block/exit/yield)
- ※ work-conserving, no starvation, waiting time could be high!
(~non-preemptive)

② RR ~ Round Robin

~ preemptive at every time interval.

- dequeue from head of ready queue & execute for a specific max. time interval.
- on time-out or block (move to end of queue).

↳ dequeue next ready process

※ work-conserving, no starvation,

pre-emptive.

good for: fairness, response time.

Execution & finish time

can suffer (esp. at high load).

{ time interval $\rightarrow 0$ fair scheduler
 (context switch overheads increase)
 - time interval $>> 0$ high overheads
 (large) \Rightarrow FCFS.

③ Priority-based scheduler

e.g. priority levels

```

graph TD
    A(( )) --> B[gold]
    A --> C[silver]
    A --> D[bronze]
  
```

~ a separate ready queue per priority level.

~ each process has a tagged priority.

~ FCFS/RR used for each ready queue.

※ ~ always schedule process from higher priority first!

⊕ pre-emptive, priority sensitive, (higher priority processes get priority for CPU usage).

⊖ Starvation! - lower priority process may never get CPU.

④ Proportionate Scheduling

WRR

weighted round robin

- one ready queue per priority.
- time slice per queue (weight) is proportional to priority.
e.g.: 4:2:1 for 3 queues.

↳ 4 tasks B priority 1

2 tasks of —— 2

1 task B —— 1

and repeat (round robin).

⊕ priority sensitive
no starvation.

⊖ finish time depends
on weights & load.

Lottery Scheduling

- # tickets issued to each process.
- random number generated to select winner/ticket \Rightarrow process.
- # tickets \Rightarrow probability of scheduling.
tickets $> 0 \Rightarrow$ no starvation!

⊕ probabilistic & proportionate scheduling.

more policies:

- SJF: shortest/smallest job first.
- STCF: Smallest Time to Completion first.
Shortest

⑤ Real world

- arrival of processes (in ready) non-deterministic queue

- process execution time unknown. (and execution times are different for different processes)

- process behaviour (interactive vs CPU intensive) unknown.

⑥ real scheduler example

(i) # multi-level feedback queue based scheduler. (MLFQ).

- part of Linux for a long time

- multiple queues (one per priority)

- priority of each process is dynamic/changing.

the MLFQ Scheduling Policy —

1. if priority (process A) \geq priority (process B)
A gets CPU.
2. if priority (A) = priority (B), A & B run in RR fashion
using a time slice for the
given queue.
3. a process entering the system, Placed at highest priority
(READY for first time) (topmost queue).
4. Once process uses time slice/allotment at a given level
its priority is reduced (moved down one queue level).
5. After some time period, move all processes to topmost queue
or wrap around after last queue

⊕ note: each queue has a different time slice for RR.

⊕ no starvation, preemptive, priority sensitive & fair.

short jobs finish quickly, CPU-intensive make regular progress.
long jobs

(ii) CFS ————— current Linux scheduler
Completely Fair Scheduler. (default for non real time tasks).

(iii) multi CPU (multiprocessor) scheduling.

- multiple queues (one per CPU)
- which queue a process should be placed in? (load balancing).
- should process move across queues? (migration).