

recap: process abstraction

- PCB  $\left\{ \begin{array}{l} \text{pid, ppid, state} \\ \text{memory info.} \end{array} \right.$

- OS game plan  $\longrightarrow$  handcraft the first user process & start its execution

- first user process can consume the OS interface/services.

- fork

↑ duplicates a process

& exec

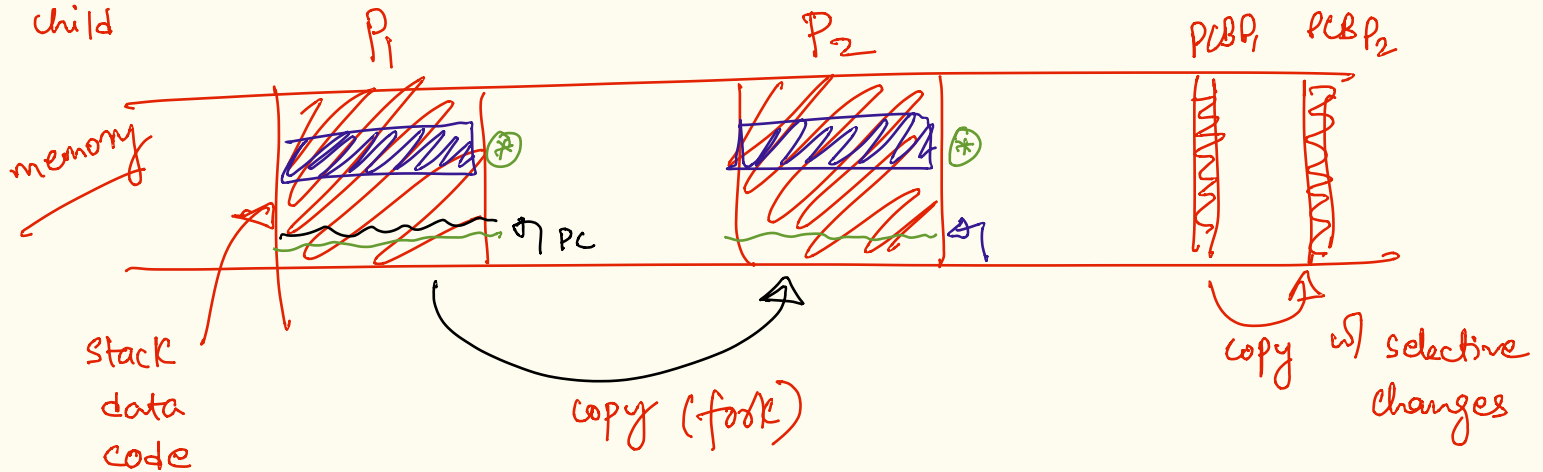
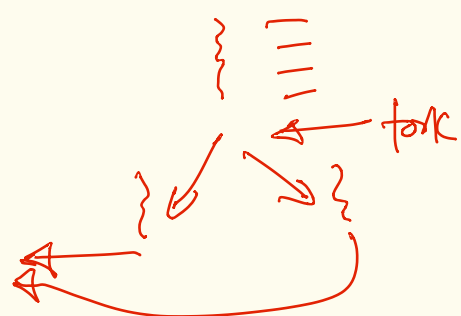
↑ loads a new program (into a process)

```
int a = 23;
```

returns pid of child in parent

```
fork();
a++;
Print(a);
```

returns 0 in child



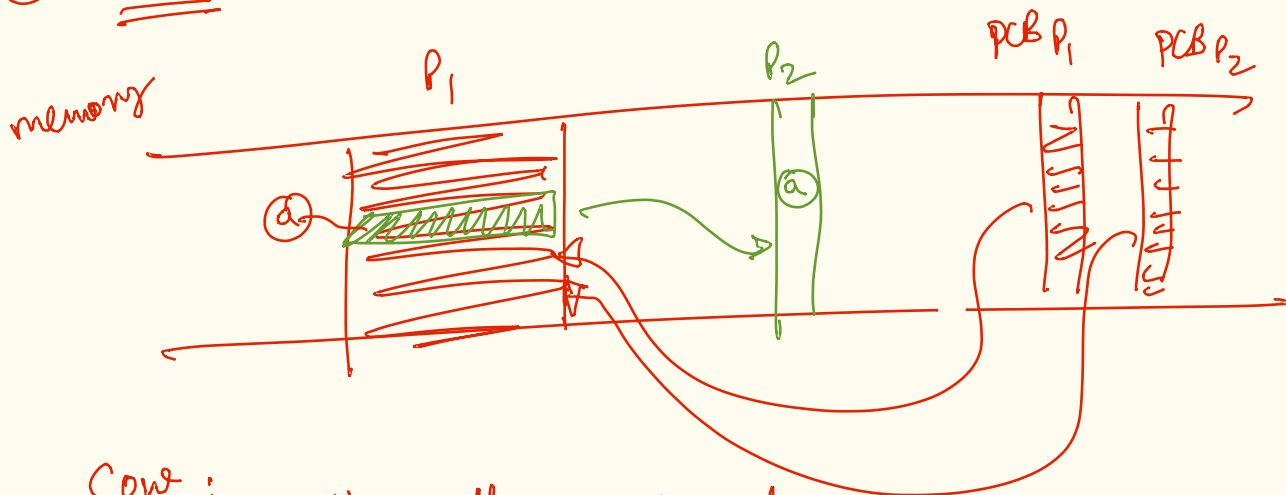
## \* why fork?

- duplicate & setup a process with custom configurations.
  - default all state is copied (open files etc.)
  - after fork, can set the subset of open files, cwd, etc. before
- multi-process parallel work - eg. webserver
- is create a process framework for new programs!

## \* exec

- loads a new program (in a process)
  - cleans up context (CPU regs)
  - sets PC to first instruction
- from disk into process memory
- code
  - data
- stack & heap

## # Cow - copy-on-write



Cow: - share all memory regions || mark all regions "read-only"

- only make copy per process on a write!

## (ii) signals

- OS mechanism for inter-process communication (of events)
- process-level intercept/signaling mechanisms

- usage

OS-support.

+ <sup>pending</sup> per process signal state.  
in the PCB

+ before a process is scheduled on  
the CPU, pending signals are  
processed!